TMS320C24x DSP Controllers Reference Set

Volume 1: CPU, System, and Instruction Set

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Preface

Read This First

About This Manual

This manual (volume 1 of a 2-volume set) describes the architecture, central processing unit (CPU), system hardware, assembly language instructions, and general operation of the TMS320C24x digital signal processor (DSP) controllers. In this document, the TMS320C24x is also referred to as the ’C24x.

The *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices* (literature number SPRU161) describes the peripherals available in the ’C24x family and their operation. Also described are specific device configurations of the ’C24x family.

For a summary of updates in this book, see Appendix E, *Summary of Updates in This Document.*

How to Use This Manual

The following table summarizes the ’C24x information contained in this manual:

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Addressing modes (for addressing data memory)	Chapter 7, *Addressing Modes*
Assembly language instructions	Chapter 8, *Assembly Language Instructions*
Comparison of assembly language instructions for TMS320C1x, ’C2x, ’C24x, and ’C5x	Appendix A, *TMS320C1x/C2x/C24x/C5x Instruction Set Comparisons*
CPU	Chapter 3, *Central Processing Unit*
Custom ROM from TI	Appendix B, *Submitting ROM Codes to TI*
Emulator	Appendix C, *Design Considerations for Using XDS510 Emulator*
Features	Chapter 1, *Introduction* Chapter 2, *Architectural Overview*

If you are looking for information about	Turn to
Input/output ports	Chapter 4, *Memory and I/O Spaces*
Interrupts	Chapter 6, *System Functions*
Memory configuration	Chapter 4, *Memory and I/O Spaces*
Peripheral interface	Chapter 6, *System Functions*
Pipeline	Chapter 5, *Program Control*
Power-down modes	Chapter 6, *System Functions*
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Reset	Chapter 6, *System Functions*
Stack	Chapter 5, *Program Control*
Status registers	Chapter 6, *System Functions*
Summary of Changes in This Document	Appendix E, *Summary of Updates in This Document*

Notational Conventions

This document uses the following conventions:

- Program listings and program examples are shown in a `special typeface`.

  Here is a segment of a program listing:

```
OUTPUT  LDP     #6          ;select data page 6
        BLDD    #300, 20h   ;move data at address 300h to 320h
        RET
```

- In syntax descriptions, the instruction is in a **bold typeface** and parameters are in an *italic typeface*. Portions of a syntax in **bold** must be entered as shown; portions of a syntax in *italics* describe the type of information that you specify. Here is an example of an instruction syntax:

  BLDD *source, destination*

  BLDD is the instruction and has two parameters, *source* and *destination*. When you use **BLDD**, the first parameter must be an actual data memory source address and the second parameter must be a destination address. A comma and a space (optional) must separate the two addresses.

- Square brackets, [ ], identify an optional parameter. If you use an optional parameter, specify the information within the brackets; do not type the brackets themselves. When you specify more than one optional parameter from a list, you separate them with a comma and a space. Here is a sample syntax:

BLDD *source*, *destination* [, **AR***n*]

BLDD is the instruction. The two required operands are *source* and *destination*, and the optional operand is **AR***n*. **AR** is bold and *n* is italic; if you choose to use **AR***n*, you must type the letters A and R and then supply a chosen value for *n* (in this case, a value from 0 to 7). Here is an example:

```
BLDD *, #310h, AR3
```

Information About Cautions

This book contains cautions.

This is an example of a caution statement.

A caution statement describes a situation that could potentially damage your software or equipment.

CAUTION

The information in a caution is provided for your protection. Please read each caution carefully.

Related Documentation From Texas Instruments

The following books describe the ’C24x and related support tools. To obtain a copy of any of these TI documents, call the Texas Instruments Literature Response Center at (800) 477–8924. When ordering, please identify the book by its title and literature number.

TMS320C24x DSP Controllers Reference Set Volume 2: Peripheral Library and Specific Devices (literature number SPRU161) describes the peripherals available on the TMS320C24x digital signal processor controllers and their operation. Also described are specific device configurations of the ’C24x family.

TMS320C240, TMS320F240 DSP Controllers (literature number SPRS042) data sheet contains the electrical and timing specifications for these devices, as well as signal descriptions and pinouts for all of the available packages.

TMS320C1x/C2x/C2xx/C5x Code Generation Tools Getting Started Guide (literature number SPRU121) describes how to install the TMS320C1x, TMS320C2x, TMS320C2xx, and TMS320C5x assembly language tools and the C compiler for the ’C1x, ’C2x, ’C2xx, and ’C5x devices. The installation for MS-DOS™, OS/2™, SunOS™, and Solaris™ systems is covered.

TMS320C1x/C2x/C2xx/C5x Assembly Language Tools User's Guide (literature number SPRU018) describes the assembly language tools (assembler, linker, and other tools used to develop assembly language code), assembler directives, macros, common object file format, and symbolic debugging directives for the 'C1x, 'C2x, 'C2xx, and 'C5x generations of devices.

TMS320C2x/C2xx/C5x Optimizing C Compiler User's Guide (literature number SPRU024) describes the 'C2x/C2xx/C5x C compiler. This C compiler accepts ANSI standard C source code and produces TMS320 assembly language source code for the 'C2x, 'C2xx, and 'C5x generations of devices.

TMS320C2xx C Source Debugger User's Guide (literature number SPRU151) tells you how to invoke the 'C2xx emulator and simulator versions of the C source debugger interface. This book discusses various aspects of the debugger interface, including window management, command entry, code execution, data management, and breakpoints. It also includes a tutorial that introduces basic debugger functionality.

TMS320C2xx Simulator Getting Started (literature number SPRU137) describes how to install the TMS320C2xx simulator and the C source debugger for the 'C2xx. The installation for MS-DOS™, PC-DOS™, SunOS™, Solaris™, and HP-UX™ systems is covered.

TMS320C2xx Emulator Getting Started Guide (literature number SPRU209) tells you how to install the Windows™ 3.1 and Windows™ 95 versions of the 'C2xx emulator and C source debugger interface.

XDS51x Emulator Installation Guide (literature number SPNU070) describes the installation of the XDS510™, XDS510PP™, and XDS510WS™ emulator controllers. The installation of the XDS511™ emulator is also described.

XDS522/XDS522A Emulation System Installation Guide (literature number SPRU171) describes the installation of the emulation system. Instructions include how to install the hardware and software for the XDS522™ and XDS522A™.

XDS522A Emulation System User's Guide (literature number SPRU169) tells you how to use the XDS522A™ emulation system. This book describes the operation of the breakpoint, tracing, and timing functionality in the XDS522A emulation system. This book also discusses BTT software interface and includes a tutorial that uses step-by-step instructions to demonstrate how to use the XDS522A emulation system.

XDS522A Emulation System Online Help (literature number SPRC002) is an online help file that provides descriptions of the BTT software user interface, menus, and dialog boxes.

JTAG/MPSD Emulation Technical Reference (literature number SPDU079) provides the design requirements of the XDS510™ emulator controller, discusses JTAG designs (based on the IEEE 1149.1 standard), and modular port scan device (MPSD) designs.

TMS320 DSP Development Support Reference Guide (literature number SPRU011) describes the TMS320 family of digital signal processors and the tools that support these devices. Included are code-generation tools (compilers, assemblers, linkers, etc.) and system integration and debug tools (simulators, emulators, evaluation modules, etc.). Also covered are available documentation, seminars, the university program, and factory repair and exchange.

Digital Signal Processing Applications with the TMS320 Family, Volumes 1, 2, and 3 (literature numbers SPRA012, SPRA016, SPRA017) Volumes 1 and 2 cover applications using the ’C10 and ’C20 families of fixed-point processors. Volume 3 documents applications using both fixed-point processors, as well as the ’C30 floating-point processor.

TMS320 DSP Designer’s Notebook: Volume 1 (literature number SPRT125) presents solutions to common design problems using ’C2x, ’C3x, ’C4x, ’C5x, and other TI DSPs.

TMS320 Third-Party Support Reference Guide (literature number SPRU052) alphabetically lists over 100 third parties that provide various products that serve the family of TMS320 digital signal processors. A myriad of products and applications are offered—software and hardware development tools, speech recognition, image processing, noise cancellation, modems, etc.

Related Technical Articles

The following technical articles contain beneficial information regarding designs, operations, and applications for signal-processing systems; all of the documents provide additional references.

“A Greener World Through DSP Controllers”, Panos Papamichalis, *DSP & Multimedia Technology*, September 1994.

“A Single-Chip Multiprocessor DSP for Image Processing—TMS320C80”, Dr. Ing. Dung Tu, *Industrie Elektronik*, Germany, March 1995.

“Application Guide with DSP Leading-Edge Technology”, Y. Nishikori, M. Hattori, T. Fukuhara, R.Tanaka, M. Shimoda, I. Kudo, A.Yanagitani, H. Miyaguchi, et al., *Electronics Engineering*, November 1995.

"Approaching the No-Power Barrier", Jon Bradley and Gene Frantz, *Electronic Design*, January 9, 1995.

"Beware of BAT: DSPs Add Brilliance to New Weapons Systems", Panos Papamichalis, *DSP & Multimedia Technology*, October 1994.

"Choose DSPs for PC Signal Processing", Panos Papamichalis, *DSP & Multimedia Technology*, January/February 1995.

"Developing Nations Take Shine to Wireless", Russell MacDonald, Kara Schmidt and Kim Higden, *EE Times*, October 2, 1995.

"Digital Signal Processing Solutions Target Vertical Application Markets", Ron Wages, *ECN*, September 1995.

"Digital Signal Processors Boost Drive Performance", Tim Adcock, *Data Storage*, September/October 1995.

"DSP and Speech Recognition, An Origin of the Species", Panos Papamichalis, *DSP & Multimedia Technology*, July 1994.

"DSP Design Takes Top-Down Approach", Andy Fritsch and Kim Asal, *DSP Series Part III*, *EE Times*, July 17, 1995.

"DSPs Advance Low-Cost 'Green' Control", Gregg Bennett, *DSP Series Part II*, *EE Times*, April 17, 1995.

"DSPs Do Best on Multimedia Applications", Doug Rasor, *Asian Computer World*, October 9–16, 1995.

"DSPs: Speech Recognition Technology Enablers", Gene Frantz and Gregg Bennett, *I&CS*, May 1995.

"Easing JTAG Testing of Parallel-Processor Projects", Tony Coomes, Andy Fritsch, and Reid Tatge, *Asian Electronics Engineer*, Manila, Philippines, November 1995.

"Fixed or Floating? A Pointed Question in DSPs", Jim Larimer and Daniel Chen, *EDN*, August 3, 1995.

"Function-Focused Chipsets: Up the DSP Integration Core", Panos Papamichalis, *DSP & Multimedia Technology*, March/April 1995.

"GSM: Standard, Strategien und Systemchips", Edgar Auslander, *Elektronik Praxis*, Germany, October 6, 1995.

"High Tech Copiers to Improve Images and Reduce Paperwork", Karl Guttag, *Document Management*, July/August 1995.

"Host-Enabled Multimedia: Brought to You by DSP Solutions", Panos Papamichalis, *DSP & Multimedia Technology*, September/October 1995.

"Integration Shrinks Digital Cellular Telephone Designs", Fred Cohen and Mike McMahan, *Wireless System Design*, November 1994.

"On-Chip Multiprocessing Melds DSPs", Karl Guttag and Doug Deao, *DSP Series Part III*, *EE Times*, July 18, 1994.

"Real-Time Control", Gregg Bennett, *Appliance Manufacturer*, May 1995.

"Speech Recognition", P.K. Rajasekaran and Mike McMahan, *Wireless Design & Development*, May 1995.

"Telecom Future Driven by Reduced Milliwatts per DSP Function", Panos Papamichalis, *DSP & Multimedia Technology*, May/June 1995.

"The Digital Signal Processor Development Environment", Greg Peake, *Embedded System Engineering*, United Kingdom, February 1995.

"The Growing Spectrum of Custom DSPs", Gene Frantz and Kun Lin, *DSP Series Part II*, *EE Times*, April 18, 1994.

"The Wide World of DSPs, " Jim Larimer, *Design News*, June 27, 1994.

"Third-Party Support Drives DSP Development for Uninitiated and Experts Alike", Panos Papamichalis, *DSP & Multimedia Technology*, December 1994/January 1995.

"Toward an Era of Economical DSPs", John Cooper, *DSP Series Part I*, *EE Times*, Jan. 23, 1995.

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❑ World-Wide Web Sites

TI Online	http://www.ti.com
Semiconductor Product Information Center (PIC)	http://www.ti.com/sc/docs/pic/home.htm
DSP Solutions	http://www.ti.com/dsps
320 Hotline On-line™	http://www.ti.com/sc/docs/dsps/support.htm

❑ North America, South America, Central America

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TI Literature Response Center U.S.A.	(800) 477-8924		
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DSP Hotline	+03-3769-8735 or (INTL) 813-3769-8735	Fax: +03-3457-7071 or (INTL) 813-3457-7071
DSP BBS via Nifty-Serve	Type “Go TIASP”	

❑ Documentation

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Mail: Texas Instruments Incorporated
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Houston, Texas 77251-1443

Email: comments@books.sc.ti.com

Note: When calling a Literature Response Center to order documentation, please specify the literature number of the book.

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Chapter 1

Introduction

The TMS320C24x ('C24x) series is a member of the TMS320 family of digital signal processors (DSPs). The 'C24x series is designed to meet a wide range of digital motor control (DMC) applications. This chapter provides an overview of the current TMS320 family, describes the background and benefits of the 'C24x DSP controller products, and introduces the 'C240 device.

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1.1 TMS320 Family Overview

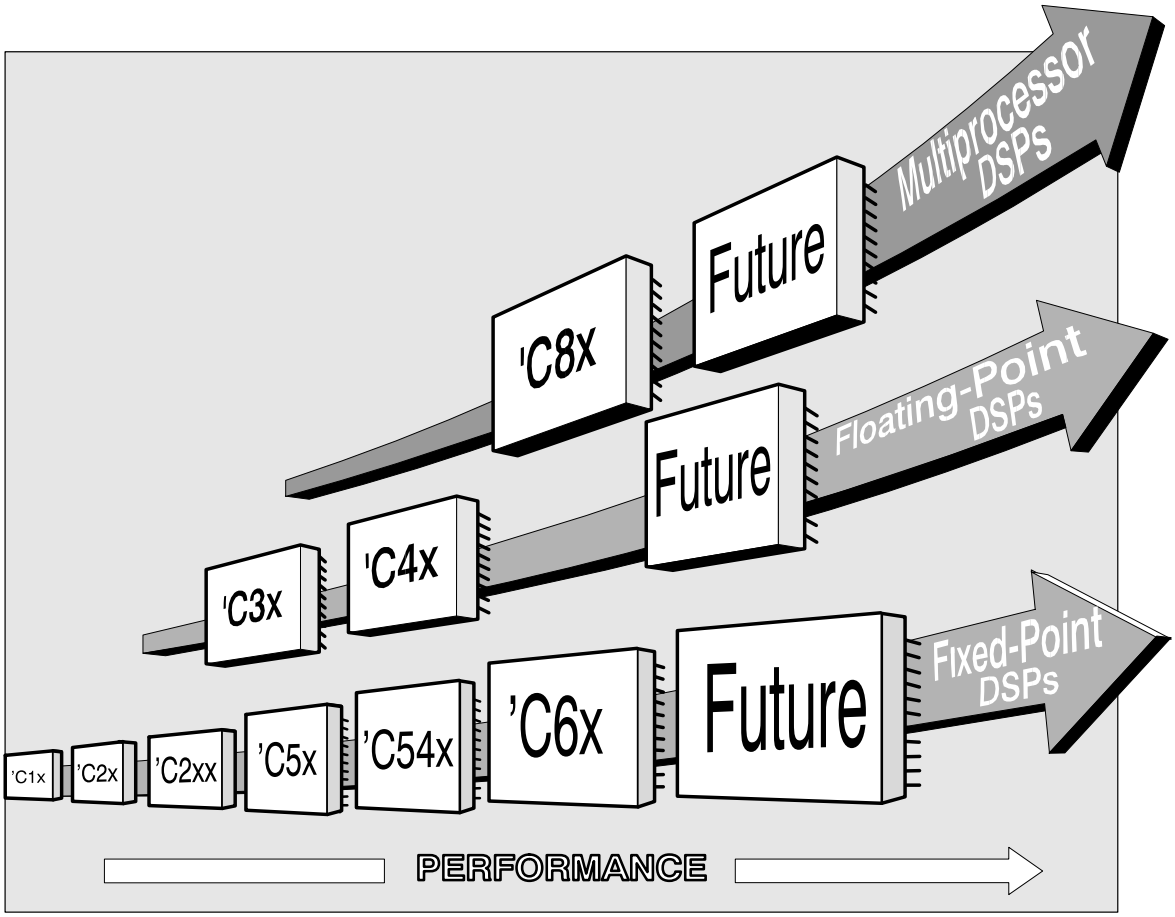
The TMS320 family consists of fixed-point, floating-point, multiprocessor digital signal processors (DSPs), and fixed-point DSP controllers. TMS320 DSPs have an architecture designed specifically for real-time signal processing. The ’C24x series of DSP controllers combines this real-time processing capability with controller peripherals to create an ideal solution for control system applications. The following characteristics make the TMS320 family the right choice for a wide range of processing applications:

- Very flexible instruction set
- Inherent operational flexibility
- High-speed performance
- Innovative parallel architecture
- Cost-effectiveness

In 1982, Texas Instruments introduced the TMS32010, the first fixed-point DSP in the TMS320 family. Before the end of the year, *Electronic Products* magazine awarded the TMS32010 the title “Product of the Year”. Today, the TMS320 family consists of these generations (Figure 1–1): ’C1x, ’C2x, ’C2xx, ’C5x, ’C54x, and ’C6x fixed-point DSPs; ’C3x and ’C4x floating-point DSPs; and ’C8x multiprocessor DSPs. The ’C24x is considered part of the ’C2xx family of fixed-point DSPs.

Devices within a generation of the TMS320 family have the same CPU structure but different on-chip memory and peripheral configurations. Spin-off devices use new combinations of on-chip memory and peripherals to satisfy a wide range of needs in the worldwide electronics market. By integrating memory and peripherals onto a single chip, TMS320 devices reduce system costs and save circuit board space.

Figure 1–1. TMS320 Family



1.2 TMS320C24x Series of DSP Controllers

Designers are recognizing the opportunity to redesign existing DMC systems to use advanced algorithms, yielding better performance and reducing system component count. DSPs are enabling:

- Design of robust controllers for a new generation of inexpensive motors, such as AC induction, DC permanent magnet, and switched-reluctance motors
- Full variable-speed control of brushless motor types that have lower manufacturing cost and higher reliability
- Energy savings through variable-speed control, saving up to 25% of the energy used by fixed-speed controllers
- Increased fuel economy, improved performance, and elimination of hydraulic fluid in automotive electronic power steering (EPS) systems
- Reduced manufacturing and maintenance costs by eliminating hydraulic fluids in automotive electronic braking systems
- More efficient and quieter operation due to less generation of torque ripple, resulting in less loss of power, lower vibration, and longer life
- Elimination or reduction of memory lookup tables through real-time polynomial calculation, thereby reducing system cost
- Use of advanced algorithms that can reduce the number of sensors required in a system
- Control of power switching inverters along with control algorithm processing
- Single-processor control of multimotor systems

The ’C24x DSP controllers are designed to meet the needs of control-based applications. By integrating the high performance of a DSP core and the on-chip peripherals of a microcontroller into a single-chip solution, the ’C24x series yields a device that is an affordable alternative to traditional microcontroller units (MCUs) and expensive multichip designs. At 20 million instructions per second (MIPS), the ’C24x DSP controllers offer significant performance over traditional 16-bit microcontrollers and microprocessors.

The 16-bit, fixed-point DSP core of the ’C24x devices provides analog designers a digital solution that does not sacrifice the precision and performance of their systems. In fact, system performance can be enhanced through the use

of advanced control algorithms for techniques such as adaptive control, Kalman filtering, and state control. The ’C24x DSP controllers offer reliability and programmability. Analog control systems, on the other hand, are hard-wired solutions and can experience performance degradation due to aging, component tolerance, and drift.

The high-speed central processing unit (CPU) allows the digital designer to process algorithms in real time rather than approximate results with look-up tables. The instruction set of these DSP controllers, which incorporates both signal processing instructions and general-purpose control functions, coupled with the extensive development support available for the ’C24x devices, reduces development time and provides the same ease of use as traditional 8- and 16-bit microcontrollers. The instruction set also allows you to retain your software investment when moving from other general-purpose TMS320 fixed-point DSPs. It is source- and object-code compatible with the other members of the ’C2xx generation, source code compatible with the ’C2x generation, and upward source code compatible with the ’C5x generation of DSPs from Texas Instruments.

The ’C24x architecture is also well-suited for processing control signals. A 16-bit word length is used along with 32-bit registers for storing intermediate results, and two hardware shifters are available to scale numbers independently of the CPU. This combination minimizes quantization and truncation errors, and increases processing power for additional functions. Such functions might include a notch filter that could cancel mechanical resonances in a system or an estimation technique that could eliminate state sensors in a system.

The ’C24x DSP controllers take advantage of an existing set of peripheral functions that allow Texas Instruments to quickly configure various series members for different price/performance points or for application optimization. This library of both digital and mixed-signal peripherals includes:

- Timers
- Serial communications ports (SCI, SPI)
- Analog-to-digital converters (ADC)
- Event manager
- System protection, such as low-voltage detection and watchdog timers

The DSP controller peripheral library is continually growing and changing to suit the needs of tomorrow’s embedded control marketplace.

1.3 TMS320C240 Overview

The TMS320C240 is the first standard device introduced in the ’C24x series of DSP controllers. It sets the standard for a single-chip digital motor controller. The ’C240 can execute 20 MIPS. Almost all instructions are executed in a single cycle of 50 ns. This high performance allows real-time execution of very complex control algorithms, such as adaptive control and Kalman filters. Very high sampling rates can also be used to minimize loop delays.

The ’C240 has the architectural features necessary for high-speed signal processing and digital control functions, and it has the peripherals needed to provide a single-chip solution for motor control applications. The ’C240 is manufactured using submicron CMOS technology, achieving a low power dissipation rating. Also included are several power-down modes for further power savings.

Applications that benefit from the advanced processing power of the ’C240 include:

- Industrial motor drives
- Power inverters and controllers
- Automotive systems, such as electronic power steering, anti-lock brakes, and climate control
- Appliance and HVAC blower/compressor motor controls
- Printers, copiers, and other office products
- Tape drives, magnetic optical drives, and other mass storage products
- Robotics and CNC milling machines

To function as a system manager, DSPs must have robust on-chip I/O and other peripherals. The event manager of the ’C240 is unlike any other available on a DSP. This application-optimized peripheral unit, coupled with the high-performance DSP core, enables the use of advanced control techniques for high-precision and high-efficiency full variable-speed control of all motor types. Included in the event manager are special pulse-width modulation (PWM) generation functions, such as a programmable dead-band function and a space vector PWM state machine for three-phase motors that provides state-of-the-art maximum efficiency in the switching of power transistors. Three independent up/down timers, each with it’s own compare register, support the generation of asymmetric (noncentered) as well as symmetric (centered) PWM waveforms. Two of the four capture inputs are direct connections for quadrature encoder pulse signals from an optical encoder.

Here is a summary of ’C240 features:

- TMS320C2xx core CPU:
  - 32-bit central arithmetic logic unit (CALU)
  - 32-bit accumulator
  - 16-bit × 16-bit parallel multiplier with a 32-bit product capability
  - Three scaling shifters
  - Eight 16-bit auxiliary registers with a dedicated arithmetic unit for indirect addressing of data memory
- Memory:
  - 544 words × 16 bits of on-chip data/program dual-access RAM
  - 16K words × 16 bits of on-chip program ROM or flash EEPROM
  - 224K words × 16 bits of maximum addressable memory space (64K words of program space, 64K words of data space, 64K words of I/O space, and 32K words of global space)
  - External Memory Interface Module with a software wait-state generator, a 16-bit address bus, and a 16-bit data bus
  - Support of hardware wait-states
- Program control:
  - Four-level pipeline operation
  - Eight-level hardware stack
  - Six external interrupts: power-drive protection interrupt, reset, NMI, and three maskable interrupts
- Instruction set:
  - Source code compatibility with ’C2x, ’C2xx, and ’C5x fixed-point generations of the TMS320 family
  - Single-instruction repeat operation
  - Single-cycle multiply/accumulate instructions
  - Memory block move instructions for program/data management
  - Indexed-addressing capability
  - Bit-reversed indexed-addressing capability for radix-2 fast Fourier transforms (FFTs)
- Power:
  - Static CMOS technology
  - Four power-down modes to reduce power consumption

- Emulation: IEEE Standard 1149.1 test access port interface to on-chip scan-based emulation logic
- Speed: 50-ns (20 MIPS) instruction cycle time, with most instructions single-cycle
- Event manager:
  - 12 compare/pulse-width modulation (PWM) channels (9 independent)
  - Three 16-bit general-purpose timers with six modes, including continuous up counting and continuous up/down counting
  - Three 16-bit full compare units with dead band capability
  - Three 16-bit simple compare units
  - Four capture units, two of which have quadrature encoder-pulse interface capability
- Dual 10-bit analog-to-digital converter
- 28 individually programmable, multiplexed I/O pins
- Phase-locked loop (PLL)-based clock module
- Watchdog timer module with real-time interrupt
- Serial communication interface (SCI)
- Serial peripheral interface (SPI)

Chapter 2

Architectural Overview

This chapter provides an overview of the architectural structure and components of the ’C24x. The ’C24x uses an advanced, modified Harvard architecture that maximizes processing power by maintaining separate bus structures for program memory and data memory.

Topic **Page**

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2.2	Memory	2-5
2.3	Central Processing Unit	2-8
2.4	Program Control	2-10
2.5	On-Chip Peripherals	2-11
2.6	Serial-Scan Emulation	2-11

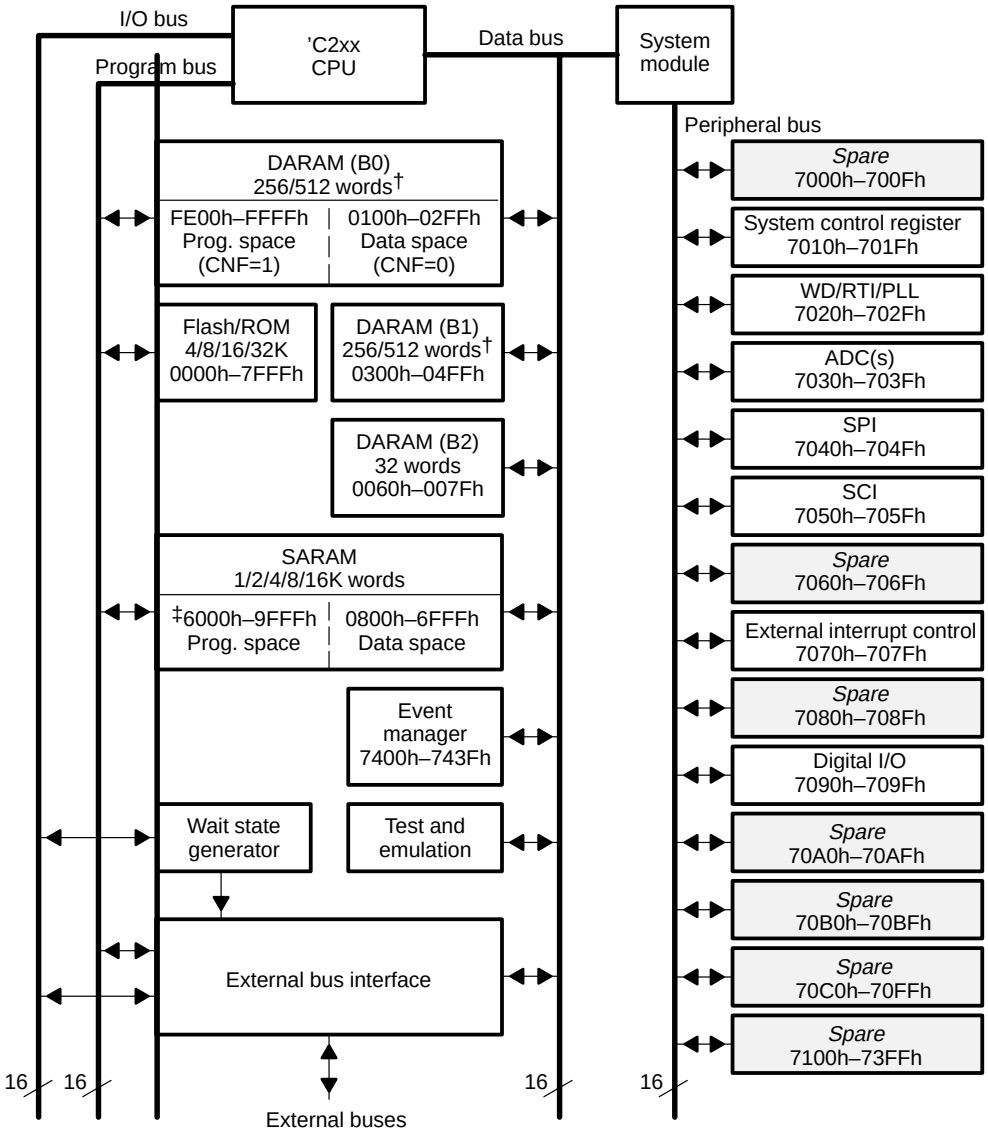
2.1 Architecture Summary

A high-level block diagram of the ’C24x architecture is shown in Figure 2–1. The ’C24x architecture is based on the modified Harvard architecture, which supports separate bus structures for program space and data space. A third space, the input/output (I/O) space, is also available and is accessible through the *external bus interface* (shown at the bottom of the figure). To support a large selection of peripherals, a peripheral bus is used. The peripheral bus is mapped to the data space and interfaced to the data bus through the *system module*. Thus, all the instructions that operate on the data space also operate on all the peripheral registers.

Separate program and data spaces allow simultaneous access to program instructions and data. For example, while data is multiplied, a previous product can be added to the accumulator, and, at the same time, a new address can be generated. Such parallelism supports a set of arithmetic, logic, and bit-manipulation operations that can all be performed in a single machine cycle. The ’C24x also includes control mechanisms to manage interrupts, repeated operations, and function/subroutine calls.

The bus structure shown in Figure 2–1 forms the basis of the entire ’C24x generation of devices. In addition, the CPU is identical for all ’C24x devices. However, each different device configuration has a unique combination of memory and peripheral modules. In the figure, the address ranges given for the memory modules are for the maximum allowable memory sizes. Typically, specific ’C24x devices have subsets of these ranges. The peripheral locations shown in the figure are true for all ’C24x devices. If more than one instance of a certain peripheral is on a device, each additional instance occupies one of the slots labeled *Spare* in the figure. The exact memory and peripheral configurations for a specific ’C24x device are defined in the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices* and in the device data sheet.

Figure 2–1. ’C24x High-Level Block Diagram



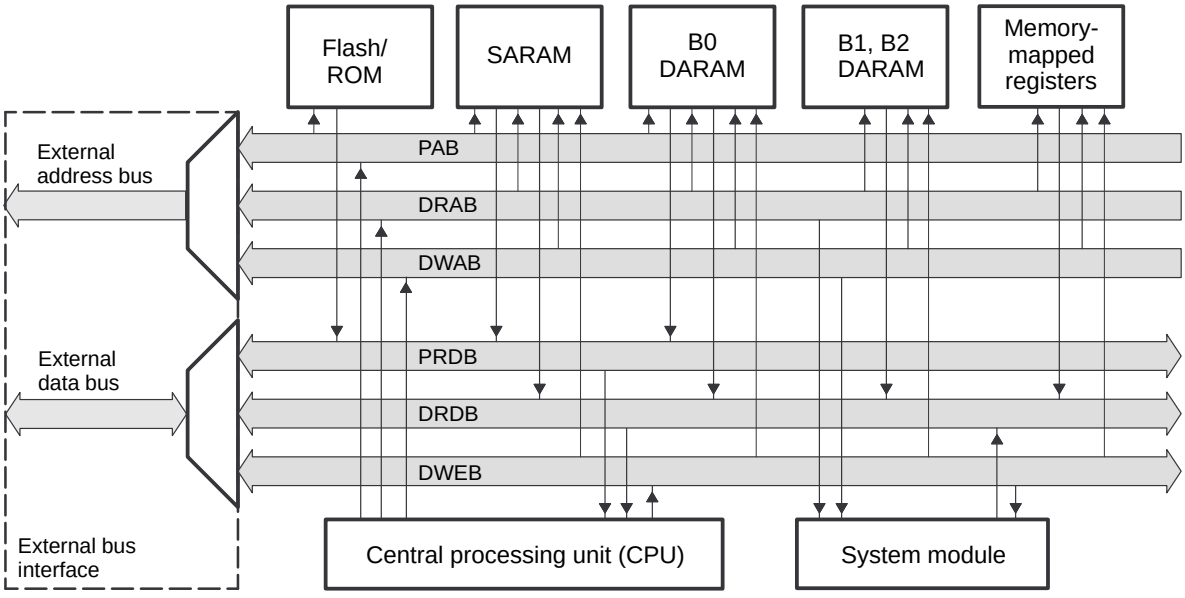
† Size of DARAM depends on device
‡ Start address of SARAM depends on end address of ROM/Flash

The internal data and program bus structure is further divided into six 16-bit buses (see Figure 2–2):

- **PAB.** The *program address bus* provides addresses for both reads from and writes to program memory.
- **DRAB.** The *data-read address bus* provides addresses for reads from data memory.
- **DWAB.** The *data-write address bus* provides addresses for writes to data memory.
- **PRDB.** The *program read bus* carries instruction code and immediate operands, as well as table information, from program memory to the CPU.
- **DRDB.** The *data-read bus* carries data from data memory to the central arithmetic logic unit (CALU) and the auxiliary register arithmetic unit (ARAU).
- **DWEB.** The *data-write bus* carries data to both program memory and data memory.

Having separate address buses for data reads (DRAB) and data writes (DWAB) allows the CPU to read and write in the same machine cycle.

Figure 2–2. ’C24x Address and Data Bus Structure



2.2 Memory

The ’C24x can contain the following kinds of on-chip memory:

- Dual-access RAM (DARAM)
- Single-access RAM (SARAM)
- Flash EEPROM or ROM (masked)

The ’C24x memory is organized into four individually-selectable spaces:

- Program (64K words)
- Local data (64K words)
- Global data (32K words)
- Input/Output (64K words)

These spaces form an address range of 224K words. For a detailed description of the ’C24x memory and I/O spaces, see Chapter 4, *Memory and I/O Spaces*.

2.2.1 On-Chip Dual-Access RAM (DARAM)

The ’C24x devices can have up to a maximum of 1056 words of on-chip DARAM, which can be accessed twice per machine cycle. This memory is primarily intended to hold data but, when needed, can also be used to hold programs. The memory can be configured in one of two ways, depending on the state of the CNF bit of status register ST1.

If total DARAM is 1056 words:

- When CNF = 0, all 1056 words are configured as data memory.
- When CNF = 1, 544 words are configured as data memory and 512 words are configured as program memory.

If total DARAM is 544 words:

- When CNF = 0, all 544 words are configured as data memory.
- When CNF = 1, 288 words are configured as data memory and 256 words are configured as program memory.

For the DARAM configurations of a particular ’C24x device, see the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices* and the device data sheet.

Because DARAM can be accessed twice per cycle, it improves the speed of the CPU. The CPU operates within a 4-cycle pipeline. In this pipeline, the CPU reads data on the third cycle and writes data on the fourth cycle. However, DARAM allows the CPU to write and read in one cycle; the CPU writes to DARAM on the master phase of the cycle and reads from DARAM on the slave phase. For example, suppose two instructions, A and B, store the accumulator value to DARAM and load the accumulator with a new value from DARAM. Instruction A stores the accumulator value during the master phase of the CPU cycle, and instruction B loads the new value to the accumulator during the slave phase. Because part of the dual-access operation is a write, it only applies to RAM.

2.2.2 On-Chip Program/Data Single-Access RAM (SARAM)

The ’C24x can have up to 16K 16-bit words of single-access RAM (SARAM), starting at address 800h in data space and the top of ROM/Flash in program space. These addresses can be used for both data memory and program memory. For example, in Figure 2–1, the SARAM block is double mapped to both program and data space. Code can be booted from off-chip ROM and then executed at full speed once it is loaded into the on-chip SARAM.

SARAM is accessed only once per CPU cycle. When the CPU requests multiple accesses, the SARAM schedules the accesses by providing a not-ready condition to the CPU and then executing the accesses, one per cycle. For example, if the instruction sequence involves storing the accumulator value and then loading a value to the accumulator, it would take two cycles to complete in SARAM, compared to one cycle in DARAM.

The SARAM block allows for more flexible address mapping than the DARAM block because SARAM can be mapped to program and data memory at the same time. Because of this, however, an instruction fetch and a data fetch that could be performed in one cycle using DARAM may take two cycles with SARAM.

2.2.3 Flash EEPROM

The ’C24x family supports from 4K words to 64K words of on-chip flash EEPROM. The flash memory supports single cycle/single access in read mode. In write mode (programming), the ’C24x requires the regular 5-V supply on pin V_{CCP} . The higher programming voltage is generated by on-chip charge pumps.

Detailed descriptions of the control register and bit functions used to program the flash block are given in the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices*.

2.2.4 Factory-Masked ROM

For large-volume applications in which the software is stable and free of bugs, low-cost, masked ROM is available. ROM sizes from 4K words to 32K words are supported. If you want a custom ROM, you can provide the code or data to be programmed into the ROM in object-file format, and Texas Instruments will generate the appropriate process mask to program the ROM. See Appendix B, *Submitting ROM Codes to TI*, for details.

2.2.5 External Memory Interface Module

In addition to full, on-chip memory support, the ’C24x can provide access to external memory by way of the *external memory interface* module. This interface provides 16 external address lines, 16 external data lines, and relevant control signals to select data, program, and I/O spaces. A wait-state generator allows interfacing with slower off-chip memory and peripherals. For a detailed description of the external memory interface, see the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices*.

2.3 Central Processing Unit

The ’C2xx CPU is on all the ’C24x devices. It contains:

- A 32-bit central arithmetic logic unit (CALU)
- A 32-bit accumulator
- Input and output data-scaling shifters for the CALU
- A 16-bit × 16-bit multiplier
- A product-scaling shifter
- Data-address generation logic, which includes eight auxiliary registers and an auxiliary register arithmetic unit (ARAU)
- Program-address generation logic

2.3.1 Central Arithmetic Logic Unit (CALU) and Accumulator

The ’C24x performs 2s-complement arithmetic using the 32-bit CALU. The CALU uses 16-bit words taken from data memory or derived from an immediate instruction, or it uses the 32-bit result from the multiplier. In addition to arithmetic operations, the CALU can perform Boolean operations.

The accumulator stores the output from the CALU; it can also provide a second input to the CALU. The accumulator is 32 bits wide and is divided into a high-order word (bits 31 through 16) and a low-order word (bits 15 through 0). Assembly language instructions are provided for storing the high- and low-order accumulator words to data memory.

2.3.2 Scaling Shifters

The ’C24x has three 32-bit shifters that allow for scaling, bit extraction, extended arithmetic, and overflow-prevention operations:

- **Input data-scaling shifter (input shifter).** This shifter left shifts 16-bit input data by 0 to 16 bits to align the data to the 32-bit input of the CALU.
- **Output data-scaling shifter (output shifter).** This shifter can left shift output from the accumulator by 0 to 7 bits before the output is stored to data memory. The content of the accumulator remains unchanged.
- **Product-scaling shifter (product shifter)**. The product register (PREG) receives the output of the multiplier. The product shifter shifts the output of the PREG before that output is sent to the input of the CALU. The product shifter has four product shift modes (no shift, left shift by one bit, left shift by four bits, and right shift by six bits), which are useful for performing multiply/accumulate operations, performing fractional arithmetic, or justifying fractional products.

2.3.3 Multiplier

The on-chip multiplier performs 16-bit × 16-bit 2s-complement multiplication with a 32-bit result. In conjunction with the multiplier, the ’C24x uses the 16-bit temporary register (TREG) and the 32-bit product register (PREG). The TREG always supplies one of the values to be multiplied. The PREG receives the result of each multiplication.

Using the multiplier, TREG, and PREG, the ’C24x efficiently performs fundamental DSP operations such as convolution, correlation, and filtering. The effective execution time of each multiplication instruction can be as short as one CPU cycle.

2.3.4 Auxiliary Register Arithmetic Unit (ARAU) and Auxiliary Registers

The ARAU generates data memory addresses when an instruction uses indirect addressing (see Chapter 7, *Addressing Modes*) to access data memory. The ARAU is supported by eight auxiliary registers (AR0 through AR7), each of which can be loaded with a 16-bit value from data memory or directly from an instruction word. Each auxiliary register value can also be stored to data memory. The auxiliary registers are referenced by a 3-bit auxiliary register pointer (ARP) embedded in status register ST0.

2.4 Program Control

Several hardware and software mechanisms provide program control:

- Program control logic decodes instructions, manages the 4-level pipeline, stores the status of operations, and decodes conditional operations. Hardware elements included in the program control logic are the program counter, the status registers, the stack, and the address-generation logic.
- Software mechanisms used for program control include branches, calls, conditional instructions, a repeat instruction, reset, interrupts, and power-down modes.

Table 2–1 shows where you can find detailed information about these program control features.

Table 2–1. Where to Find Information About Program Control Features

For information about	See
Address-generation logic	Chapter 5, *Program Control*
Branches, calls, and returns	Chapter 5, *Program Control*
Conditional operations	Chapter 5, *Program Control*
Interrupts	Chapter 6, *System Functions*
Pipeline	Chapter 5, *Program Control*
Power-down modes	Chapter 6, *System Functions*
Program counter	Chapter 5, *Program Control*
Repeat instruction	Chapter 5, *Program Control*
Reset	Chapter 6, *System Functions*
Stack	Chapter 5, *Program Control*
Status registers	Chapter 3, *Central Processing Unit*

2.5 On-Chip Peripherals

As shown in Figure 2–1 (page 2-3), the ’C24x bus structure supports access to numerous peripherals. Two types of bus interfaces are used for the on-chip peripherals. Most of the peripherals are accessed using the peripheral bus. This bus is mapped to the data space through the control of the system module. Each access to one of these peripherals requires more than one cycle. However, the event manager fits directly onto the data bus and takes advantage of the full speed of the DSP central processing unit (CPU). An access to the event manager is made with zero wait states; a read takes one cycle and a write takes two cycles.

Each individual ’C24x device has a unique combination of peripheral modules. However, the address locations for the peripherals are fixed and are the same for all ’C24x devices. If more than one instance of a particular peripheral is on the device, each additional instance occupies one of the locations labeled *Spare* in Figure 2–1.

For detailed descriptions of the peripherals, refer to the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices*.

2.6 Serial-Scan Emulation

The ’C24x has seven pins dedicated to the serial scan emulation port (JTAG port). This port allows for nonintrusive emulation of the ’C24x devices and is supported by Texas Instruments emulation tools and by many third party debugger tools. For documentation on these emulation and debugger tools, see *Related Documentation From Texas Instruments* on page v of the preface and Appendix C, *Design Considerations for Using XDS510 Emulator*.

Chapter 3

Central Processing Unit

This chapter describes the ’C24x central processing unit (CPU) operations. The CPU can perform high-speed arithmetic operations within one instruction cycle because of its parallel architectural design.

First, this chapter describes three fundamental sections of the CPU (see Figure 3–1). The chapter then describes the auxiliary register arithmetic unit (ARAU), which performs arithmetic operations independently of the central arithmetic logic section. The chapter concludes with a description of status registers ST0 and ST1, which contain bits for determining processor modes, addressing pointer values, and indicating various processor conditions and arithmetic logic results.

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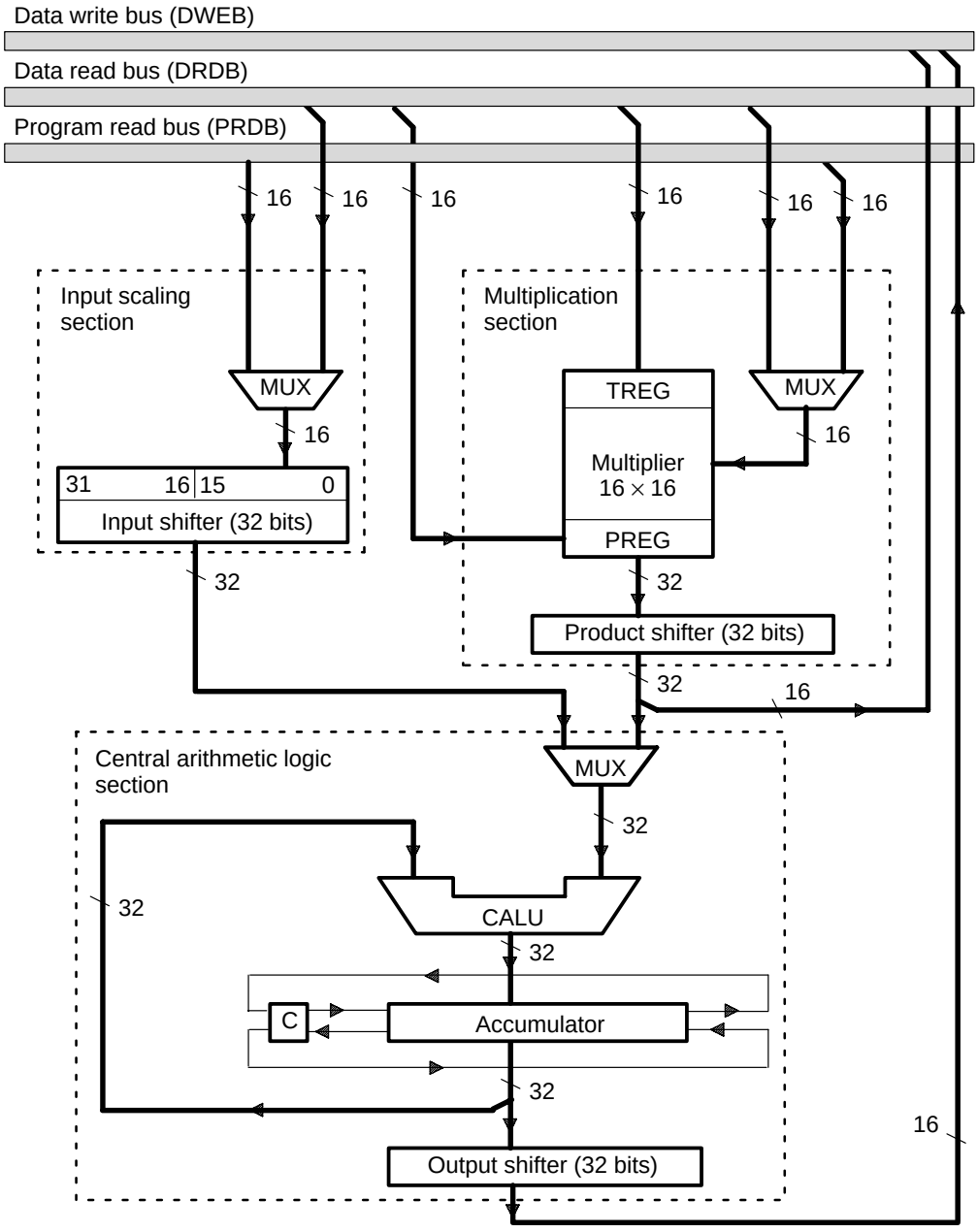
3.1 Input Scaling Section 3-3
3.2 Multiplication Section 3-5
3.3 Central Arithmetic Logic Section 3-8
3.4 Auxiliary Register Arithmetic Unit (ARAU) 3-12
3.5 Status Registers ST0 and ST1 3-15

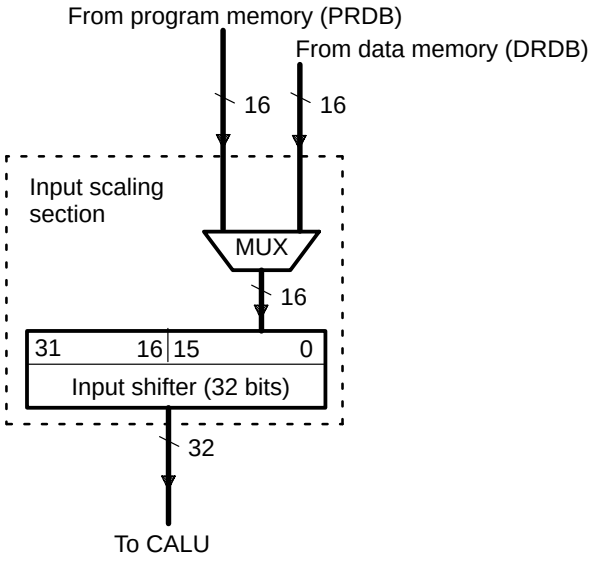
Figure 3–1. Block Diagram of the Input Scaling, Central Arithmetic Logic, and Multiplication Sections of the CPU



3.1 Input Scaling Section

A 32-bit input data-scaling shifter (input shifter) aligns a 16-bit value coming from memory to the 32-bit central arithmetic logic unit (CALU). This data alignment is necessary for data-scaling arithmetic, as well as aligning masks for logical operations. The input shifter operates as part of the data path between program or data space and the CALU and, thus, requires no cycle overhead. Described directly below are the input, the output, and the shift count of the input shifter. Throughout the discussion, refer to Figure 3–2.

Figure 3–2. Block Diagram of the Input Scaling Section



Input. Bits 15 through 0 of the input shifter accept a 16-bit input from either of two sources (see Figure 3–2):

- ❑ *The data read bus (DRDB).* This input is a value from a data memory location referenced in an instruction operand.
- ❑ *The program read bus (PRDB).* This input is a constant value given as an instruction operand.

Output. After a value has been accepted into bits 15 through 0, the input shifter aligns the16-bit value to the 32-bit bus of the CALU as shown in Figure 3–2. The shifter shifts the value left 0 to 16 bits and then sends the 32-bit result to the CALU.

During the left shift, unused LSBs in the shifter are filled with 0s, and unused MSBs in the shifter are either filled with 0s or sign extended, depending on the value of the sign-extension mode bit (SXM) of status register ST1.

Shift count. The shifter can left shift a 16-bit value by 0 to 16 bits. The size of the shift (or the shift count) is obtained from one of two sources:

- *A constant embedded in the instruction word.* Putting the shift count in the instruction word allows you to use specific data-scaling or alignment operations customized for your program code.
- *The four LSBs of the temporary register (TREG).* The TREG-based shift allows the data-scaling factor to be determined dynamically so that it can be adapted to the system's performance.

Sign-extension mode bit. For many but not all instructions, the sign-extension mode bit (SXM), bit 10 of status register ST1, determines whether the CALU uses sign extension during its calculations. If SXM = 0, sign extension is suppressed. If SXM = 1, the output of the input shifter is sign extended. Figure 3–3 shows an example of an input value shifted left by eight bits for SXM = 0. The MSBs of the value passed to the CALU are zero filled. Figure 3–4 shows the same shift but with SXM = 1. The value is sign extended during the shift.

Figure 3–3. Operation of the Input Shifter for SXM = 0

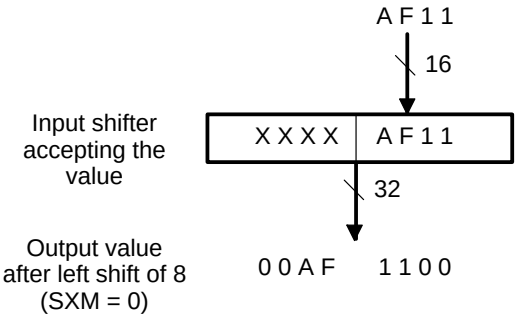
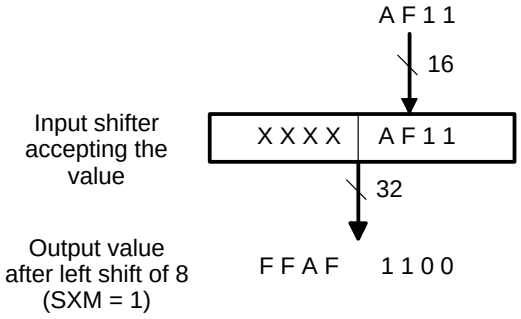


Figure 3–4. Operation of the Input Shifter for SXM = 1

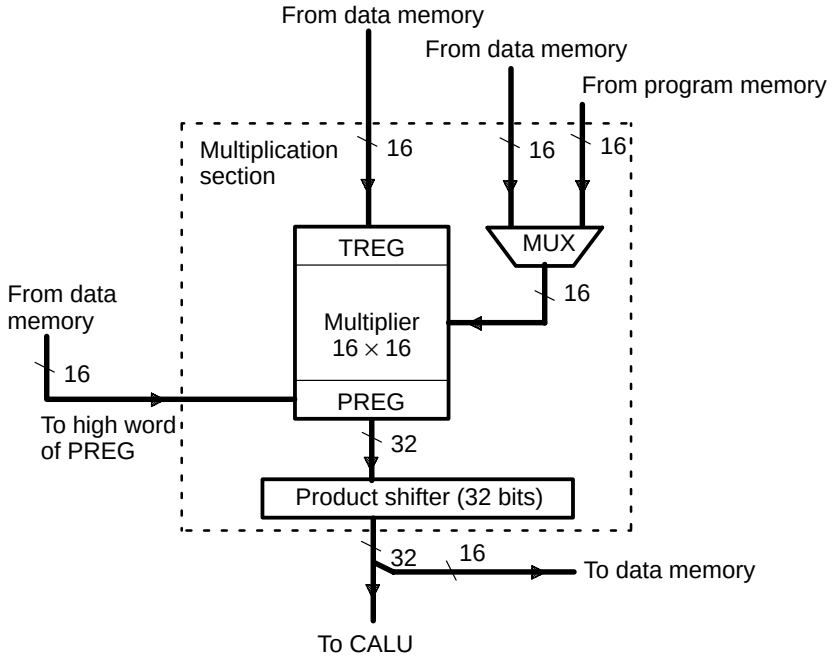


3.2 Multiplication Section

The ’C24x uses a 16-bit × 16-bit hardware multiplier that can produce a signed or unsigned 32-bit product in a single machine cycle. As shown in Figure 3–5, the multiplication section consists of:

- The 16-bit temporary register (TREG), which holds one of the multiplicands
- The multiplier, which multiplies the TREG value by a second value from data memory or program memory
- The 32-bit product register (PREG), which receives the result of the multiplication
- The product shifter, which scales the PREG value before passing it to the CALU

Figure 3–5. Block Diagram of the Multiplication Section



3.2.1 Multiplier

The 16-bit × 16-bit hardware multiplier can produce a signed or unsigned 32-bit product in a single machine cycle. The two numbers being multiplied are treated as 2s-complement numbers, except during unsigned multiplication (MPYU instruction). Descriptions of the inputs and output of the multiplier follow.

Inputs. The multiplier accepts two 16-bit inputs:

- One input is always from the 16-bit temporary register (TREG). The TREG is loaded before the multiplication with a data-value from the data read bus (DRDB).
- The other input is one of the following:
  - A data-memory value from the data read bus (DRDB)
  - A program memory value from the program read bus (PRDB)

Output. After the two 16-bit inputs are multiplied, the 32-bit result is stored in the product register (PREG). The output of the PREG is connected to the 32-bit product-scaling shifter. Through this shifter, the product may be transferred from the PREG to the CALU or to data memory (by the SPH and SPL instructions).

3.2.2 Product-Scaling Shifter

The product-scaling shifter (product shifter) facilitates scaling of the product register (PREG) value. The shifter has a 32-bit input connected to the output of the PREG and a 32-bit output connected to the input of the CALU.

Input. The shifter has a 32-bit input connected to the output of the PREG.

Output. After the shifter completes the shift, all 32 bits of the result can be passed to the CALU, or 16 bits of the result can be stored to data memory.

Shift Modes. This shifter uses one of four product shift modes, summarized in Table 3–1. As shown in the table, these modes are determined by the product shift mode (PM) bits of status register ST1. In the first shift mode (PM = 00), the shifter does not shift the product at all before giving it to the CALU or to data memory. The next two modes cause left shifts (of one or four), which are useful for implementing fractional arithmetic or justifying products. The right-shift mode shifts the product by six bits, enabling the execution of up to 128 consecutive multiply-and-accumulate operations without causing the accumulator to overflow. Note that the content of the PREG remains unchanged; the value is copied to the product shifter and shifted there.

Note:

The right shift in the product shifter is always sign extended, regardless of the value of the sign-extension mode bit (SXM) of status register ST1.

Table 3–1. Product Shift Modes for the Product-Scaling Shifter

PM	Shift	Comments†
00	No shift	Product sent to CALU or data write bus (DWEB) with no shift
01	Left 1	Removes the extra sign bit generated in a 2s-complement multiply to produce a Q31 product
10c	Left 4	Removes the extra four sign bits generated in a 16-bit × 13-bit 2s-complement multiply to produce a Q31 product when multiplying by a 13-bit constant
11	Right 6	Scales the product to allow up to 128 product accumulations without overflowing the accumulator. The right shift is always sign extended, regardless of the value of the sign-extension mode bit (SXM) of status register ST1.

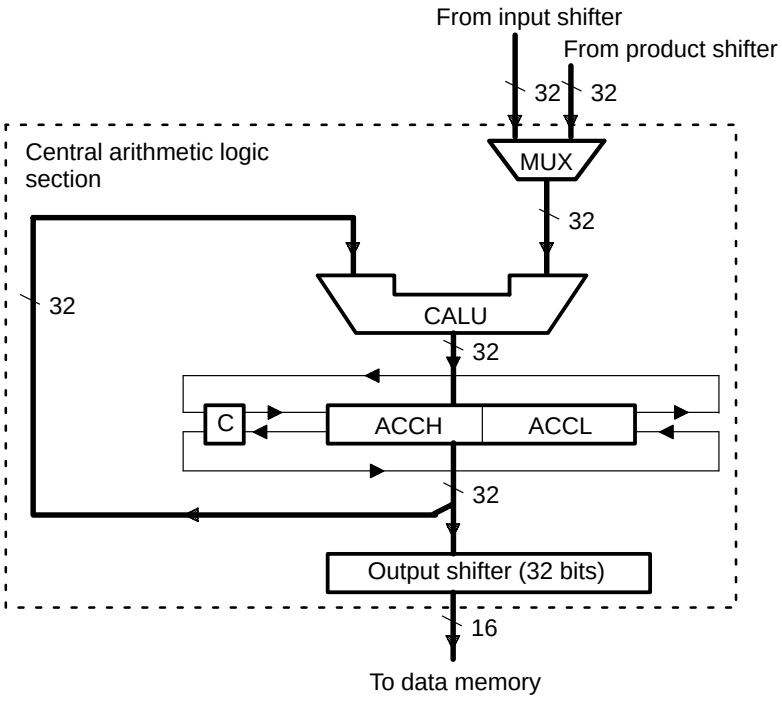
† A Q31 number is a binary fraction in which there are 31 digits to the right of the binary point (the base 2 equivalent of the base 10 decimal point).

3.3 Central Arithmetic Logic Section

Figure 3–6 shows the main components of the central arithmetic logic section, which are:

- The central arithmetic logic unit (CALU), which implements a wide range of arithmetic and logic functions
- The 32-bit accumulator (ACC), which receives the output of the CALU and is capable of performing bit shifts on its contents with the help of the carry bit (C). Figure 3–6 shows the accumulator's high word (ACCH) and low word (ACCL).
- The output shifter, which can shift a copy of either the high word or low word of the accumulator before sending it to data memory for storage

Figure 3–6. Block Diagram of the Central Arithmetic Logic Section



3.3.1 Central Arithmetic Logic Unit (CALU)

The CALU implements a wide range of arithmetic and logic functions, most of which execute in a single clock cycle. These functions can be grouped into four categories:

- 16-bit addition
- 16-bit subtraction
- Boolean logic operations
- Bit testing, shifting, and rotating

Because the CALU can perform Boolean operations, you can perform bit manipulation. For bit shifting and rotating, the CALU uses the accumulator. The CALU is referred to as central because there is an independent arithmetic unit, the auxiliary register arithmetic unit (ARAU), which is described in Section 3.4. A description of the inputs, the output, and an associated status bit of the CALU follows.

Inputs. The CALU has two inputs (see again Figure 3–6):

- One input is always provided by the 32-bit accumulator.
- The other input is provided by one of the following:
  - The product-scaling shifter (see subsection 3.2.2)
  - The input data-scaling shifter (see Section 3.1)

Output. Once the CALU performs an operation, it transfers the result to the 32-bit accumulator, which is capable of performing bit shifts of its contents. The output of the accumulator is connected to the 32-bit output data-scaling shifter. Through the output shifter, the accumulator's upper and lower 16-bit words can be individually shifted and stored to data memory.

Sign-extension mode bit. For many but not all instructions, the sign-extension mode bit (SXM), bit 10 of status register ST1, determines whether the CALU uses sign extension during its calculations. If SXM = 0, sign extension is suppressed. If SXM = 1, sign extension is enabled.

3.3.2 Accumulator

Once the CALU performs an operation, it transfers the result to the 32-bit accumulator, which can then perform single-bit shifts or rotations on its contents. Each of the accumulator's upper and lower 16-bit words can be passed to the output data-scaling shifter, where it can be shifted and then stored in data memory. The following describes the status bits and branch instructions associated with the accumulator.

Status bits. Four status bits are associated with the accumulator:

- *Carry bit (C).* C (bit 9 of status register ST1) is affected during:
  - Additions to and subtractions from the accumulator:

    C = 0 When the result of a subtraction generates a borrow

    When the result of an addition does not generate a carry (Exception: When the ADD instruction is used with a shift of 16 and no carry is generated, the ADD instruction has no effect on C.)

    C = 1 When the result of an addition generates a carry

    When the result of a subtraction does not generate a borrow (Exception: When the SUB instruction is used with a shift of 16 and no borrow is generated, the SUB instruction has no effect on C.)

  - Single-bit shifts and rotations of the accumulator value. During a left shift or rotation, the MSB of the accumulator is passed to C; during a right shift or rotation, the LSB is passed to C.

- *Overflow mode bit (OVM).* OVM (bit 11 of status register ST0) determines how the accumulator reflects arithmetic overflows. When the processor is in overflow mode (OVM = 1) and an overflow occurs, the accumulator is filled with one of two specific values:
  - If the overflow is in the positive direction, the accumulator is filled with its most positive value (7FFF FFFFh).
  - If the overflow is in the negative direction, the accumulator is filled with its most negative value (8000 0000h).

- *Overflow flag bit (OV).* OV is bit 12 of status register ST0. When no accumulator overflow is detected, OV is latched at 0. When overflow (positive or negative) occurs, OV is set to 1 and latched.

- *Test/control flag bit (TC).* TC (bit 11 of status register ST1) is set to 0 or 1 depending on the value of a tested bit. In the case of the NORM instruction, if the exclusive-OR of the two MSBs of the accumulator is true, TC is set to 1.

A number of branch instructions are implemented, based on the status of bits C, OV, and TC, and on the value in the accumulator (as compared to 0). For more information about these instructions, see Section 5.4, *Conditional Branches, Calls, and Returns*, on page 5-10.

3.3.3 Output Data-Scaling Shifter

The output data-scaling shifter (output shifter) has a 32-bit input connected to the 32-bit output of the accumulator and a 16-bit output connected to the data bus. The shifter copies all 32 bits of the accumulator and then performs a left shift on its content; it can be shifted from zero to seven bits, as specified in the corresponding store instruction. The upper word (SACH instruction) or lower word (SACL instruction) of the shifter is then stored to data memory. The content of the accumulator remains unchanged.

When the output shifter performs the shift, the MSBs are lost and the LSBs are zero filled. Figure 3–7 shows an example in which the accumulator value is shifted left by four bits and the shifted high word is stored to data memory. Figure 3–8 shows the same accumulator value shifted left by six bits and then the shifted low word stored.

Figure 3–7. Shifting and Storing the High Word of the Accumulator

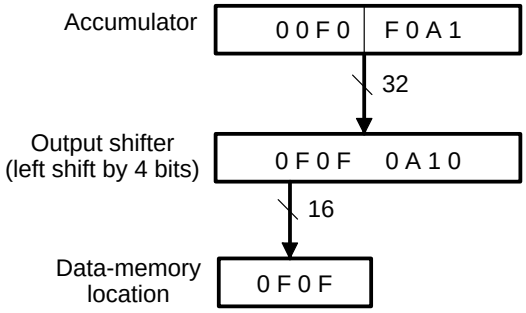
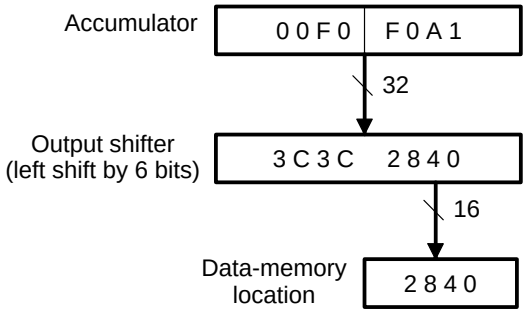


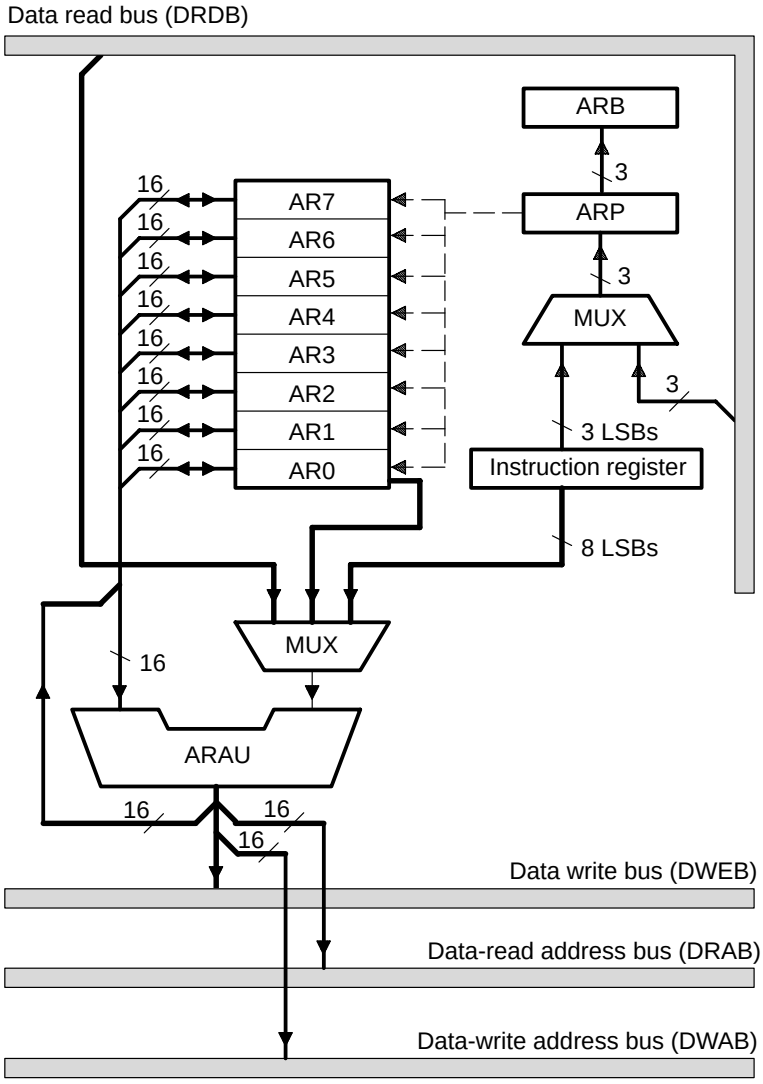
Figure 3–8. Shifting and Storing the Low Word of the Accumulator



3.4 Auxiliary Register Arithmetic Unit (ARAU)

The CPU also contains the ARAU, an arithmetic unit independent of the CALU. The main function of the ARAU is to perform arithmetic operations on eight auxiliary registers (AR7 through AR0) in parallel with operations occurring in the CALU. Figure 3–9 shows the ARAU and related logic.

Figure 3–9. ARAU and Related Logic



The eight auxiliary registers (AR7–AR0) provide flexible and powerful indirect addressing. Any location in the 64K data memory space can be accessed using a 16-bit address contained in an auxiliary register. For the details of indirect addressing, see Section 7.3 on page 7-9.

To select a specific auxiliary register, load the 3-bit auxiliary register pointer (ARP) of status register ST0 with a value from 0 through 7. The ARP can be loaded as a primary operation by the MAR instruction (which only performs modifications to the auxiliary registers and the ARP) or by the LST instruction (which can load a data-memory value to ST0 by way of the data read bus, DRDB). The ARP can be loaded as a secondary operation by any instruction that supports indirect addressing.

The register pointed to by the ARP is referred to as the *current auxiliary register* or *current AR*. During the processing of an instruction, the content of the current auxiliary register is used as the address at which the data-memory access will take place. The ARAU passes this address to the data-read address bus (DRAB) if the instruction requires a read from data memory, or it passes the address to the data-write address bus (DWAB) if the instruction requires a write to data memory. After the instruction uses the data value, the contents of the current auxiliary register can be incremented or decremented by the ARAU, which implements unsigned 16-bit arithmetic.

3.4.1 ARAU Functions

The ARAU performs the following operations:

- ❏ Increments or decrements an auxiliary register value by 1 or by an index amount (by way of any instruction that supports indirect addressing)
- ❏ Adds a constant value to an auxiliary register value (ADRK instruction) or subtracts a constant value from an auxiliary register value (SBRK instruction). The constant is an 8-bit value taken from the eight LSBs of the instruction word.
- ❏ Compares the content of AR0 with the content of the current AR and puts the result in the test/control flag bit (TC) of status register ST1 (CMPR instruction). The result is passed to TC by way of the data write bus (DWEB).

Normally, the ARAU performs its arithmetic operations in the decode phase of the pipeline (when the instruction specifying the operations is being decoded). This allows the address to be generated before the decode phase of the next instruction. There is an exception to this rule: During processing of the NORM instruction, the auxiliary register and/or ARP modification is done during the execute phase of the pipeline. For information on the operation of the pipeline, see Section 5.2 on page 5-7.

3.4.2 Auxiliary Register Functions

In addition to using the auxiliary registers to reference data-memory addresses, you can use them for other purposes. For example, you can:

- Use the auxiliary registers to support conditional branches, calls, and returns by using the CMPR instruction. This instruction compares the content of AR0 with the content of the current AR and puts the result in the test/control flag bit (TC) of status register ST1.
- Use the auxiliary registers for temporary storage by using the LAR instruction to load values into the registers and the SAR instruction to store AR values to data memory
- Use the auxiliary registers as software counters, incrementing or decrementing them as necessary

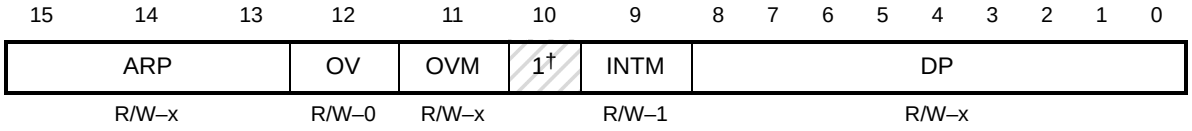
3.5 Status Registers ST0 and ST1

The ’C24x has two status registers, ST0 and ST1, which contain status and control bits. These registers can be stored into and loaded from data memory, thus allowing the status of the machine to be saved and restored for subroutines.

The LST (load status register) instruction writes to ST0 and ST1, and the SST (store status register) instruction reads from ST0 and ST1 (with the exception of the INTM bit, which is not affected by the LST instruction). Many of the individual bits of these registers can be set and cleared using the SETC and CLRC instructions. For example, the sign-extension mode is set with SETC SXM and cleared with CLRC SXM.

Figure 3–10 and Figure 3–11 show the organization of status registers ST0 and ST1, respectively. Several bits in the status registers are reserved; they are always read as logic 1s. The other bits are described in alphabetical order in Table 3–2.

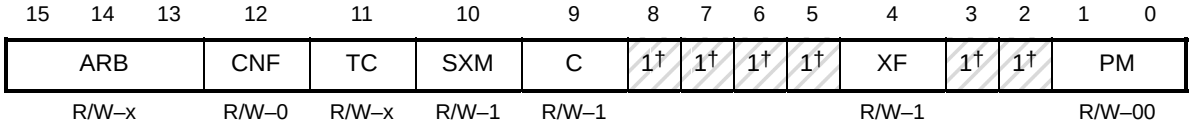
Figure 3–10. Status Register ST0



Note: R = Read access; W = Write access; value following dash (–) is value after reset (x means value not affected by reset).

† This reserved bit is always read as 1. Writes have no effect on it.

Figure 3–11. Status Register ST1



Note: R = Read access; W = Write access; value following dash (–) is value after reset (x means value not affected by reset).

† These reserved bits are always read as 1s. Writes have no effect on them.

Table 3–2. Bit Fields of Status Registers ST0 and ST1

Name	Description
ARB	**Auxiliary register pointer buffer.** Whenever the auxiliary register pointer (ARP) is loaded, the previous ARP value is copied to the ARB, except during an LST (load status register) instruction. When the ARB is loaded by an LST instruction, the same value is also copied to the ARP.
ARP	**Auxiliary register pointer.** This 3-bit field selects which auxiliary register (AR) to use in indirect addressing. When the ARP is loaded, the previous ARP value is copied to the ARB register, except during an LST (load status register) instruction. The ARP may be modified by memory-reference instructions using indirect addressing, and by the MAR (modify auxiliary register) and LST instructions. When the ARB is loaded by an LST instruction, the same value is also copied to the ARP. For more details on the use of ARP in indirect addressing, see Section 7.3, *Indirect Addressing Mode*, on page 7-9.
C	**Carry bit.** This bit is set to 1 if the result of an addition generates a carry, or cleared to 0 if the result of a subtraction generates a borrow. Otherwise, it is cleared after an addition or set after a subtraction, except if the instruction is ADD or SUB with a 16-bit shift. In these cases, ADD can only set and SUB only clear the carry bit, but cannot affect it otherwise. The single-bit shift and rotate instructions also affect this bit, as well as the SETC, CLRC, and LST instructions. The conditional branch, call, and return instructions can execute, based on the status of C. C is set to 1 on reset.
CNF	**On-chip DARAM configuration bit.** This bit determines whether reconfigurable dual-access RAM blocks are mapped to data space or to program space. The CNF bit may be modified by the SETC CNF, CLRC CNF, and LST instructions. Reset clears the CNF bit to 0. For more information about CNF and the dual-access RAM blocks, see Chapter 4, *Memory and I/O Space*s. CNF = 0 Reconfigurable dual-access RAM blocks are mapped to data space. CNF = 1 Reconfigurable dual-access RAM blocks are mapped to program space.
DP	**Data page pointer.** When an instruction uses direct addressing, the 9-bit DP field is concatenated with the seven LSBs of the instruction word to form a full 16-bit data-memory address. For more details, see Section 7.2, *Direct Addressing Mode*, on page 7-4. The LST and LDP (load DP) instructions can modify the DP field.
INTM	**Interrupt mode bit.** This bit enables or disables all maskable interrupts. INTM is set and cleared by the SETC INTM and CLRC INTM instructions, respectively. INTM has no effect on the nonmaskable interrupts $\overline{RS}$ and $\overline{NMI}$ or on interrupts initiated by software. INTM is unaffected by the LST (load status register) instruction. INTM is set to 1 when an interrupt trap is taken (except in the case of the TRAP instruction) and at reset. INTM = 0 All unmasked interrupts are enabled. INTM = 1 All maskable interrupts are disabled.
OV	**Overflow flag bit.** This bit holds a latched value that indicates whether overflow has occurred in the CALU. OV is set to 1 when an overflow occurs in the CALU. Once an overflow occurs, the OV bit remains set until it is cleared by a reset, a conditional branch on overflow (OV) or no overflow (NOV), or an LST instruction.

Table 3–2. Bit Fields of Status Registers ST0 and ST1 (Continued)

Name	Description
OVM	**Overflow mode bit.** OVM determines how overflows in the CALU are handled. The SETC and CLRC instructions set and clear this bit, respectively. An LST instruction can also be used to modify OVM. OVM = 0 Results overflow normally in the accumulator. OVM = 1 The accumulator is set to either its most positive or negative value upon encountering an overflow. (See subsection 3.3.2, *Accumulator*, on page 3-9.)
PM	**Product shift mode.** PM determines the amount that the PREG value is shifted on its way to the CALU or to data memory. Note that the content of the PREG remains unchanged; the value is copied to the product shifter and shifted there. PM is loaded by the SPM and LST instructions. The PM bits are cleared by reset. PM = 00 The multiplier's 32-bit product is passed to the CALU or to data memory with no shift. PM = 01 The output of the PREG is left shifted one place (with the LSBs zero filled) before being passed to the CALU or to data memory. PM = 10 The output of the PREG is left shifted four bits (with the LSBs zero filled) before being passed to the CALU or to data memory. PM = 11 This mode produces a right shift of six bits, sign extended.
SXM	**Sign-extension mode bit.** SXM does not affect the basic operation of certain instructions. For example, the ADDS instruction suppresses sign extension regardless of SXM. This bit is set by the SETC SXM instruction and cleared by the CLRC SXM instruction and may be loaded by the LST instruction. SXM is set to 1 by reset. SXM = 0 This mode suppresses sign extension. SXM = 1 This mode produces sign extension on data as it is passed into the accumulator from the input shifter.
TC	**Test/control flag bit.** The TC bit is set to 1 if a bit tested by BIT or BITT is a 1, if a compare condition tested by CMPR exists between the current auxiliary register and AR0, or if the exclusive-OR function of the two MSBs of the accumulator is true when tested by a NORM instruction. The conditional branch, call, and return instructions can execute, based on the condition of the TC bit. The TC bit is affected by the BIT, BITT, CMPR, LST, and NORM instructions.
XF	**XF pin status bit.** This bit determines the state of the XF pin, which is a general-purpose output pin. XF is set by the SETC XF instruction and cleared by the CLRC XF instruction. XF can also be modified with an LST instruction. XF is set to 1 by reset.

Chapter 4

Memory and I/O Spaces

Each ’C24x device has a 16-bit address line that accesses four individually selectable spaces (224K words total):

- A 64K-word program space
- A 64K-word local data space
- A 32K-word global data space
- A 64K-word I/O space

This chapter describes these four spaces and shows generic memory maps for the program, data, and I/O spaces. It also describes the ’C24x memory configuration options.

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4.2	**Program Memory**	**4-3**
4.3	**Local Data Memory**	**4-5**
4.4	**Global Data Memory**	**4-9**
4.5	**I/O Space**	**4-11**

4.1 Overview of the Memory and I/O Spaces

The ’C24x design is based on an enhanced Harvard architecture. The ’C24x has multiple memory spaces accessible on three parallel buses—the program address bus (PAB), the data-read address bus (DRAB), and the data-write address bus (DWAB). Each of the three buses access different memory spaces for different aspects of the device’s operation. Because the bus operations are independent, it is possible to access both the program and data spaces simultaneously. Within a given machine cycle, the CALU can execute as many as three concurrent memory operations.

The ’C24x address map is organized into four individually selectable spaces:

- **Program memory** (64K words) contains the instructions to be executed, as well as data used during program execution.
- **Local data memory** (64K words) holds data used by the instructions.
- **Global data memory** (32K words) shares data with other devices or serves as additional data space.
- **Input/output (I/O) space** (64K words) interfaces to external peripherals and may contain on-chip registers.

These spaces provide a total address space of 224K words. The ’C24x includes on-chip memory to aid in system performance and integration and a considerable number of addresses that can be used for external memory and I/O devices.

The advantages of operating from on-chip memory are:

- Higher performance than external memory (because the wait states required for slower external memories are avoided)
- Lower cost than external memory
- Lower power consumption than external memory

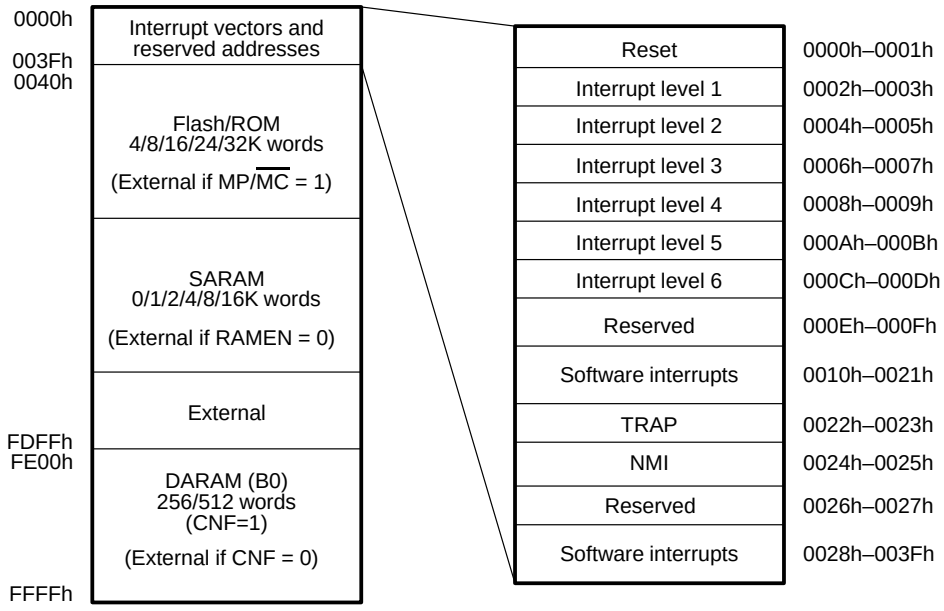
The advantage of operating from external memory is the ability to access a larger address space.

4.2 Program Memory

The program-memory space is where the application program code resides; it can also hold table information and immediate operands. The program-memory space addresses up to 64K 16-bit words. On all ’C24x devices, these words include on-chip DARAM. On-chip SARAM and on-chip ROM/flash EEPROM may be available on some of the devices. When the ’C24x generates an address outside the set of addresses configured to on-chip program memory, the device automatically generates an external access, asserting the appropriate control signals (if an external memory interface is present).

Figure 4–1 shows the program memory map.

Figure 4–1. Program Memory Map for ’C24x



Note: Flash/ROM memory includes the address range 0000h–003Fh.

4.2.1 Program Memory Configuration

Depending on which types of memory are on board a particular ’C24x, up to three factors contribute to the configuration of program memory:

- **CNF bit.** The CNF bit (bit 12) of status register ST1 determines whether the addresses for DARAM B0 are available for program space:
  - **CNF = 0.** There is no addressable on-chip program DARAM.
  - **CNF = 1.** The 256 or 512 words of DARAM B0 are configured for program use.

At reset, any words of program/data DARAM are mapped into local data space (CNF = 0).

- **MP/$\overline{MC}$ pin.** The level on the MP/$\overline{MC}$ pin determines whether program instructions are read from on-chip ROM/flash EEPROM (if available) after reset:
  - **MP/$\overline{MC}$ = 0.** The device is configured as a microcomputer. The on-chip ROM/flash EEPROM is accessible. The device fetches the reset vector from on-chip memory.
  - **MP/$\overline{MC}$ = 1.** The device is configured as a microprocessor. The device fetches the reset vector from external memory.

  Regardless of the value of MP/$\overline{MC}$, the 'C24x fetches its reset vector at location 0000h of program memory.

- **RAMEN pin.** The RAMEN signal (if available—check device data sheet) allows you to toggle a preset range of data/program addresses between on-chip SARAM (if available) and external memory:
  - **RAMEN = 1.** The preset range of addresses in both data space and program space are mapped to the same physical locations in the on-chip SARAM. For example, if 1000h were in the preset address range, 1000h in program memory and 1000h in data memory would point to the same physical location in the on-chip SARAM. Thus, the full block of on-chip SARAM is accessible for program and/or data space.

Note:

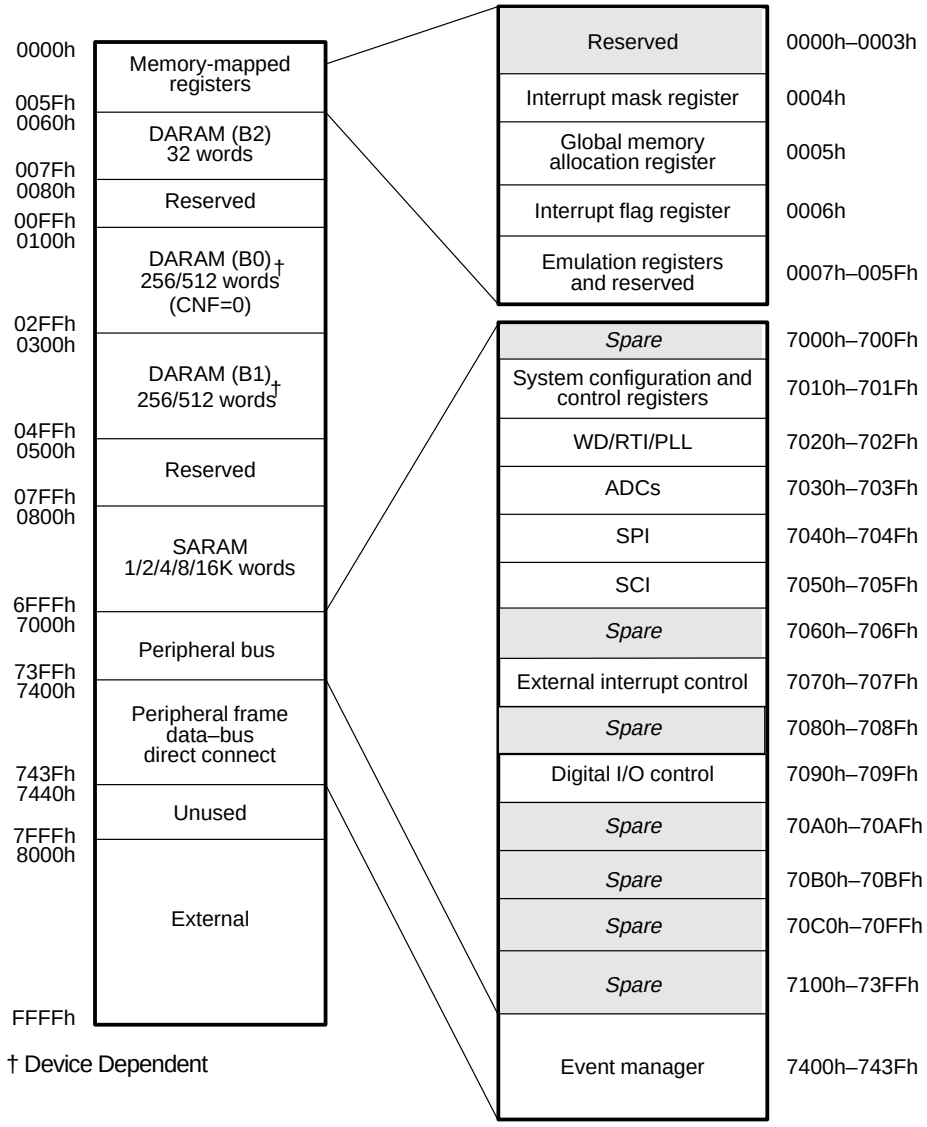
When RAMEN = 1, program memory and data memory share the same range of addresses. When writing data to these locations, be careful not to overwrite existing program instructions.

  - **RAMEN = 0.** The preset range of addresses is not shared by data memory and program memory. Instead, that range of addresses in data memory and the same range in program memory are mapped to external program memory. Thus, when RAMEN = 0, a block of addresses twice the size of the SARAM block is available for accessing external memory.

4.3 Local Data Memory

The local data-memory space addresses up to 64K 16-bit words. Figure 4–2 shows the data memory map for the ’C24x. The ’C24x devices each have three on-chip DARAM blocks: B0, B1, and B2. Block B0 is configurable as either data memory or program memory. Blocks B1 and B2 are available for data memory only. Some ’C24x devices also have an on-chip SARAM block that can be used for program and/or data memory.

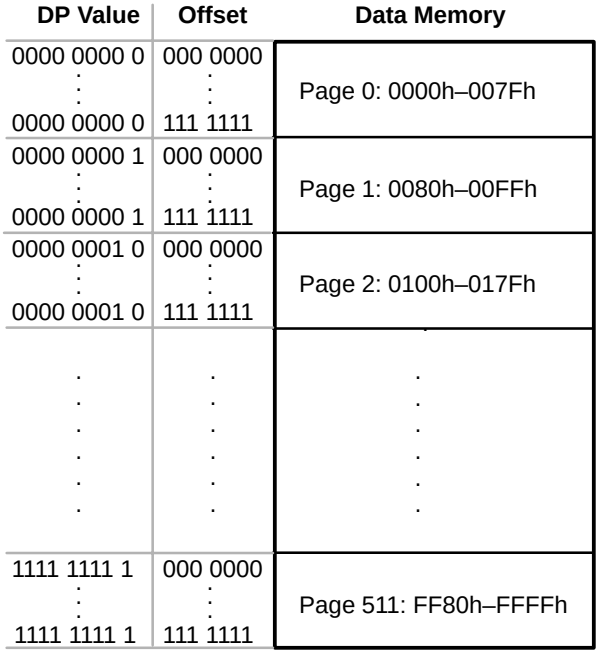
Figure 4–2. Data Memory Map for ’C24x



Data memory can be addressed with either of two addressing modes: direct-addressing or indirect-addressing. Addressing modes are described in detail in Chapter 7.

When direct addressing is used, data memory is addressed in blocks of 128 words called data pages. Figure 4–3 shows how these blocks are addressed. The entire 64K of data memory consists of 512 data pages labeled 0 through 511. The current data page is determined by the value in the 9-bit data page pointer (DP) in status register ST0. Each of the 128 words on the current page is referenced by a 7-bit offset, which is taken from the instruction that is using direct addressing. Therefore, when an instruction uses direct addressing, you must specify both the data page (with a preceding instruction) and the offset (in the instruction that accesses data memory).

Figure 4–3. Pages of Local Data Memory



4.3.1 Data Page 0 Address Map

The 64K words of local data memory include the device's memory-mapped registers, which reside at the top of data page 0 (addresses 0000h–007Fh). Note the following:

- Three registers that can be accessed with zero wait states:
  - Interrupt mask register (IMR)
  - Global memory allocation register (GREG)
  - Interrupt flag register (IFR)
- The test/emulation reserved area is used by the test and emulation systems for special information transfers.

Do Not Write to Test/Emulation Addresses

Writing to the test/emulation addresses can cause the device to change its operational mode and, therefore, affect the operation of an application.

- The scratch-pad RAM block (B2) includes 32 words of DARAM that provide for variable storage without fragmenting the larger RAM blocks, whether internal or external. This RAM block supports dual-access operations and can be addressed via any data-memory addressing mode.

Table 4–1 shows the address map of data page 0.

Table 4–1. Data Page 0 Address Map

Address	Name	Description
0000h–0003h	–	Reserved
0004h	IMR	Interrupt mask register
0005h	GREG	Global memory allocation register
0006h	IFR	Interrupt flag register
0023h–0027h	–	Reserved
002Bh–002Fh	–	Reserved for test/emulation
0060h–007Fh	B2	Scratch-pad RAM (DARAM B2)

4.3.2 Local Data Memory Configuration

Two factors may contribute to the configuration of data memory:

- **CNF bit.** The CNF bit (bit 12) of status register ST1 determines whether the on-chip DARAM B0 is mapped into local data space or into program space.
  - **CNF = 1.** DARAM B0 is used for program space.
  - **CNF = 0.** B0 is used for data space.

  At reset, B0 is mapped into local data space (CNF = 0).

- **RAMEN pin.** The RAMEN signal allows you to toggle a preset range of data/program addresses between on-chip SARAM (if available) and external memory:
  - **RAMEN = 1.** The preset range of addresses in both data space and program space are mapped to the same physical locations in the on-chip SARAM. For example, if 1000h were in the preset address range, 1000h in program memory and 1000h in data memory would point to the same physical location in the on-chip SARAM. Thus, the full block of on-chip SARAM is accessible for program and/or data space.

Note:

When RAMEN = 1, program memory and data memory share the same range of addresses. When writing data to these locations, be careful not to overwrite existing program instructions.

  - **RAMEN = 0.** The preset range of addresses is not shared by data memory and program memory. Instead, that range of addresses in data memory and the same range in program memory are mapped to external program memory. Thus, when RAMEN = 0, a block of addresses twice the size of the SARAM block is available for accessing external memory.

4.4 Global Data Memory

Addresses in the upper 32K words (8000h–FFFFh) of local data memory can be used for global data memory. The global memory allocation register (GREG) determines the size of the global data-memory space, which is between 256 and 32K words. The GREG is connected to the eight LSBs of the internal data bus and is memory-mapped to data-memory location 0005h. Table 4–2 shows the allowable GREG values and shows the corresponding address range set aside for global data memory. Any remaining addresses within 8000h–FFFFh are available for local data memory.

Note:

Choose only the GREG values listed in Table 4–2. Other values lead to fragmented memory maps.

Table 4–2. Global Data Memory Configurations

GREG Value		Local Memory		Global Memory	
High Byte	**Low Byte**	**Range**	**Words**	**Range**	**Words**
XXXX XXXX	0000 0000	0000h–FFFFh	65 536	–	0
XXXX XXXX	1000 0000	0000h–7FFFh	32 768	8000h–FFFFh	32 768
XXXX XXXX	1100 0000	0000h–BFFFh	49 152	C000h–FFFFh	16 384
XXXX XXXX	1110 0000	0000h–DFFFh	57 344	E000h–FFFFh	8 192
XXXX XXXX	1111 0000	0000h–EFFFh	61 440	F000h–FFFFh	4 096
XXXX XXXX	1111 1000	0000h–F7FFh	63 488	F800h–FFFFh	2 048
XXXX XXXX	1111 1100	0000h–FBFFh	64 512	FC00h–FFFFh	1 024
XXXX XXXX	1111 1110	0000h–FDFFh	65 024	FE00h–FFFFh	512
XXXX XXXX	1111 1111	0000h–FEFFh	65 280	FF00h–FFFFh	256

Note: X = Don't care

When a program accesses any data address, the 'C24x drives the $\overline{DS}$ signal low. If that address is within the range defined by the GREG as a global address, $\overline{BR}$ signal is also asserted. Because $\overline{BR}$ differentiates local and global accesses, the addresses configured by the GREG value are an additional data space. The external data-address range is extended by the selected amount of global space (up to 32K words).

As an example of configuring global memory, suppose you want to designate 8K data-memory addresses as global addresses. You would write the 8-bit value 11100000 to the GREG (see Figure 4–4). This would designate addresses E000h–FFFFh of data memory as global data addresses (see Figure 4–5).

Figure 4–4. GREG Register Set to Configure 8K for Global Data Memory

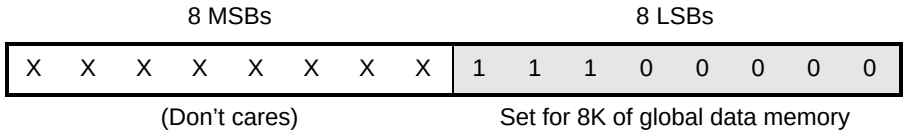
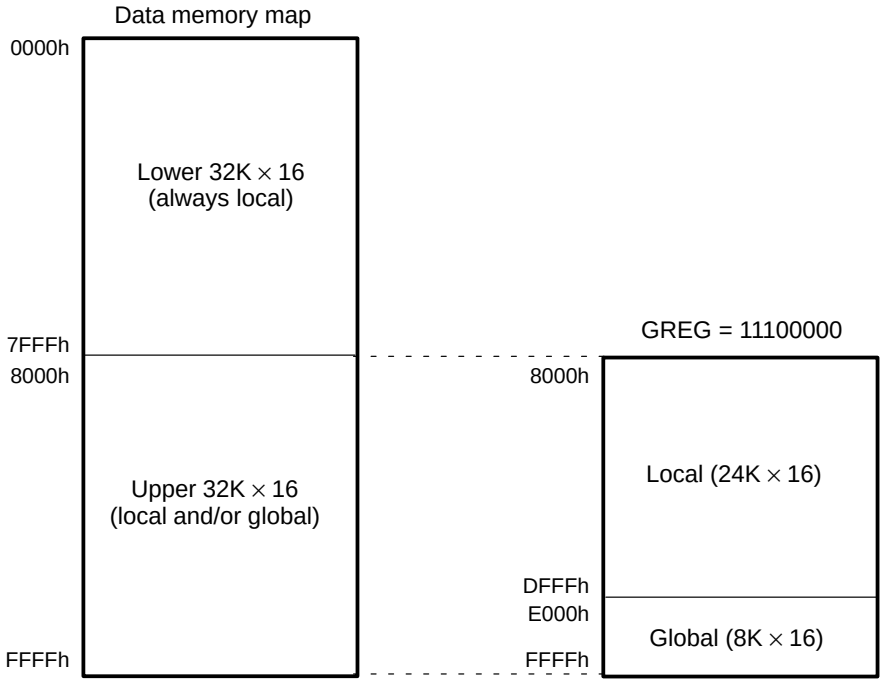


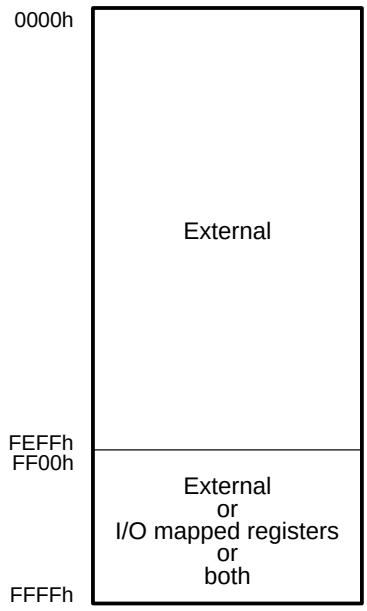
Figure 4–5. Global and Local Data Memory for GREG = 11100000



4.5 I/O Space

The I/O space memory addresses up to 64K 16-bit words. Figure 4–6 shows the I/O-space address map for the ’C24x.

Figure 4–6. I/O-Space Address Map for ’C24x



Note: See device-specific data sheets for information on I/O mapped registers.

Chapter 5

Program Control

This chapter discusses the processes and features involved in controlling the flow of a program on the ’C24x.

Program control involves controlling the order in which one or more blocks of instructions are executed. Normally, the flow of a program is sequential: the ’C24x executes instructions at consecutive program-memory addresses. At times, a program must branch to a nonsequential address and then execute instructions sequentially at that new location. For this purpose, the ’C24x supports branches, calls, returns, repeats, and interrupts. Interrupts are described in Chapter 6, *System Functions*.

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5.2	**Pipeline Operation**	**5-7**
5.3	**Branches, Calls, and Returns**	**5-8**
5.4	**Conditional Branches, Calls, and Returns**	**5-10**
5.5	**Repeating a Single Instruction**	**5-14**

5.1 Program-Address Generation

Program flow requires the processor to generate the next program address (sequential or nonsequential) while executing the current instruction. Program-address generation is illustrated in Figure 5–1 and summarized in Table 5–1.

Figure 5–1. Program-Address Generation Block Diagram

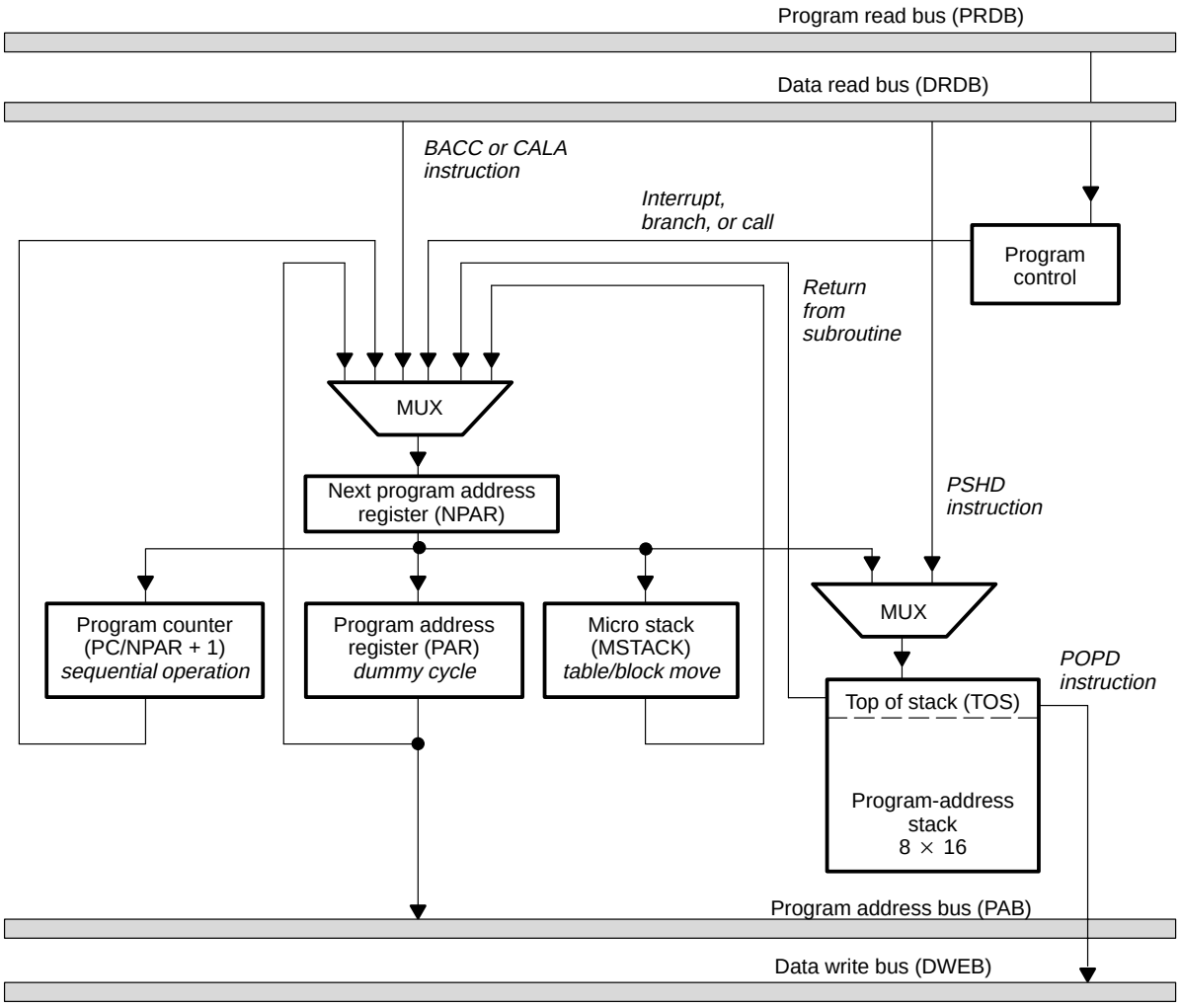


Table 5–1. Program-Address Generation Summary

Operation	Program-Address Source
Sequential operation	PC (contains program address +1)
Dummy cycle	PAR (contains program address)
Return from subroutine	Top of the stack (TOS)
Return from table move or block move	Microstack (MSTACK)
Branch or call to address specified in instruction	Branch or call instruction by way of the program read bus (PRDB)
Branch or call to address specified in lower half of the accumulator	Low accumulator by way of the data read bus (DRDB)
Branch to interrupt service routine	Interrupt vector location by way of the program read bus (PRDB)

The ’C24x program-address generation logic uses the following hardware:

- **Program counter (PC).** The ’C24x has a 16-bit program counter (PC) that addresses internal and external program memory when fetching instructions.
- **Program address register (PAR).** The PAR drives the program address bus (PAB). The PAB is a 16-bit bus that provides program addresses for both reads and writes.
- **Stack.** The program-address generation logic includes a 16-bit-wide, 8-level hardware stack for storing up to eight return addresses. In addition, you can use the stack for temporary storage.
- **Microstack (MSTACK).** Occasionally, the program-address generation logic uses the 16-bit-wide, 1-level MSTACK to store one return address.
- **Repeat counter (RPTC).** The 16-bit RPTC is used with the repeat (RPT) instruction to determine how many times the instruction following RPT is repeated.

5.1.1 Program Counter (PC)

The program-address generation logic uses the 16-bit program counter (PC) to address internal and external program memory. The PC holds the address of the next instruction to be executed. Through the program address bus (PAB), an instruction is fetched from that address in program memory and loaded into the instruction register. When the instruction register is loaded, the PC holds the next address.

The ’C24x can load the PC in a number of ways, to accommodate sequential and nonsequential program flow. Table 5–2 shows what is loaded to the PC according to the code operation performed.

Table 5–2. Address Loading to the Program Counter

Code Operation	Address Loaded to the PC
Sequential execution	The PC is loaded with PC + 1 if the current instruction has one word or PC + 2 if the current instruction has two words.
Branch	The PC is loaded with the long immediate value directly following the branch instruction.
Subroutine call and return	For a call, the address of the next instruction is pushed from the PC onto the stack, and then the PC is loaded with the long immediate value directly following the call instruction. A return instruction pops the return address back into the PC to return to the calling sequence of code.
Software or hardware interrupt	The PC is loaded with the address of the appropriate interrupt vector location. At this location is a branch instruction that loads the PC with the address of the corresponding interrupt service routine.
Computed GOTO	The content of the lower 16 bits of the accumulator is loaded into the PC. Computed GOTO operations can be performed using the BACC (branch to address in accumulator) or CALA (call subroutine at location specified by the accumulator) instructions.

5.1.2 Stack

The ’C24x has a 16-bit-wide, 8-level-deep hardware stack. The program-address generation logic uses the stack for storing return addresses when a subroutine call or interrupt occurs. When an instruction forces the CPU into a subroutine or an interrupt forces the CPU into an interrupt service routine, the return address is loaded to the top of the stack automatically; this event does not require additional cycles. When the subroutine or interrupt service routine is complete, a return instruction transfers the return address from the top of the stack to the program counter.

When the eight levels are not used for return addresses, the stack may be used for saving context data during a subroutine or interrupt service routine or for other storage purposes.

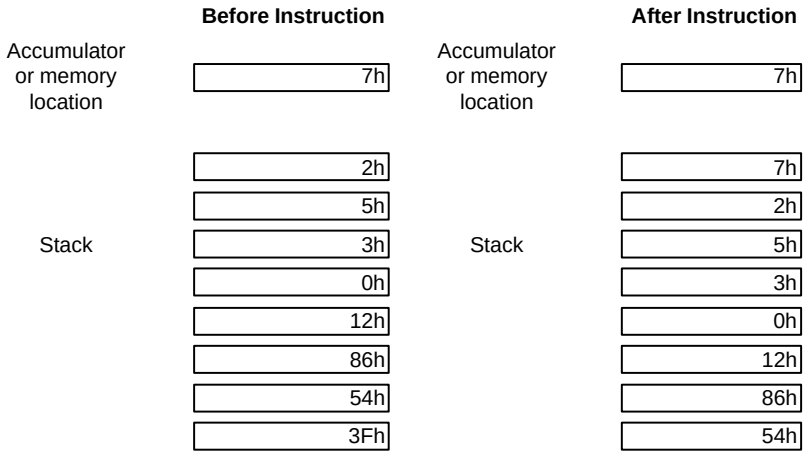
You can access the stack with two sets of instructions:

- **PUSH and POP.** The PUSH instruction copies the 16 LSBs of the accumulator to the top of the stack. The POP instruction copies the value on the top of the stack to the 16 LSBs of the accumulator.

- **PSHD and POPD.** These instructions allow you to build a stack in data memory for the nesting of subroutines or interrupts beyond eight levels. The PSHD instruction pushes a data-memory value onto the top of the stack. The POPD instruction pops a value from the top of the stack to data memory.

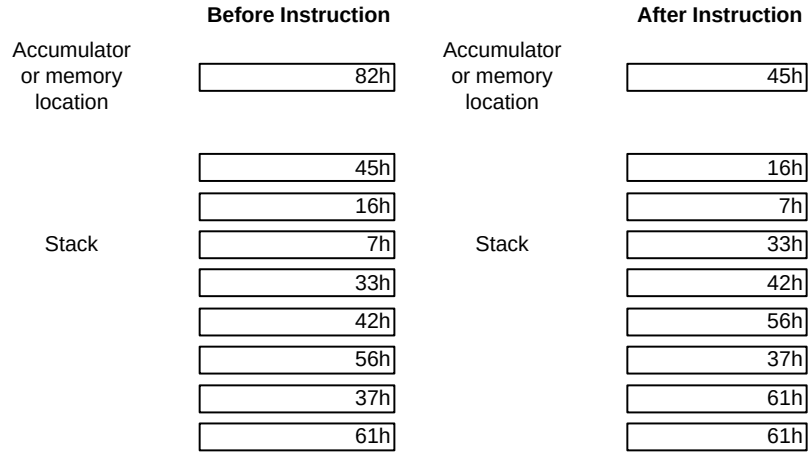
Whenever a value is pushed onto the top of the stack (by an instruction or by the address-generation logic), the content of each level is pushed down one level, and the bottom (eighth) location of the stack is lost. Therefore, data is lost (stack overflow occurs) if more than eight successive pushes occur before a pop. Figure 5–2 shows a push operation.

Figure 5–2. A Push Operation



Pop operations are the reverse of push operations. A pop operation copies the value at each level to the next higher level. Any pop after seven sequential pops yields the value that was originally at the bottom of the stack because, by then, the bottom value has been copied upward to all of the stack levels. Figure 5–3 shows a pop operation.

Figure 5–3. A Pop Operation



5.1.3 Microstack (MSTACK)

The program-address generation logic uses the 16-bit-wide, 1-level-deep MSTACK to store a return address before executing certain instructions. These instructions use the program-address generation logic to provide a second address in a 2-operand instruction. These instructions are: BLDD, BLPD, MAC, MACD, TBLR, and TBLW. When repeated, these instructions use the PC to increment the first operand address and can use the auxiliary register arithmetic unit (ARAU) to generate the second operand address. When these instructions are used, the return address (the address of the next instruction to be fetched) is pushed onto the MSTACK. Upon completion of the repeated instruction, the MSTACK value is popped back into the program-address generation logic. The MSTACK operations are not visible to you. Unlike the stack, the MSTACK can be used only by the program-address generation logic; there are no instructions that allow you to use the MSTACK for storage.

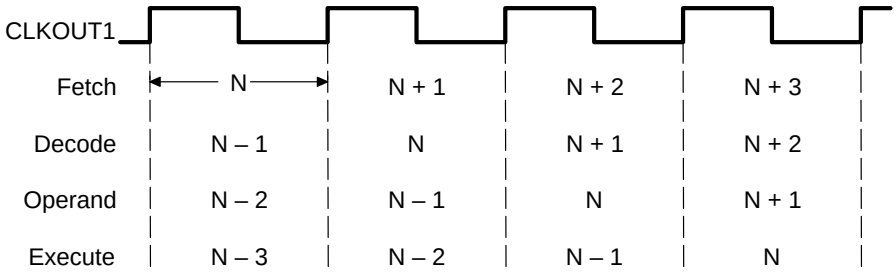
5.2 Pipeline Operation

Instruction pipelining consists of a sequence of bus operations that occur during the execution of an instruction. The ’C24x pipeline has four independent stages: instruction-fetch, instruction-decode, operand-fetch, and instruction-execute. Because the four stages are independent, these operations can overlap. During any given cycle, one to four different instructions can be active, each at a different stage of completion. Figure 5–4 shows the operation of the 4-level-deep pipeline for single-word, single-cycle instructions executing with no wait states.

The pipeline is essentially invisible to you, except in the following cases:

- A single-word, single-cycle instruction immediately following a modification of the global-memory allocation register (GREG) uses the previous global map.
- The NORM instruction modifies the auxiliary register pointer (ARP) and uses the current auxiliary register (the one pointed to by the ARP) during the execute phase of the pipeline. If the next two instruction words change the values in the current auxiliary register or the ARP, they will do so during the instruction decode phase of the pipeline (before the execution of NORM). This would cause NORM to use the wrong auxiliary register value and the following instructions to use the wrong ARP value.

Figure 5–4. Four-Level Pipeline Operation



The CPU is implemented using 2-phase static logic. The 2-phase operation of the ’C24x CPU consists of a master phase in which all commutation logic is executed, and a slave phase in which results are latched. Therefore, sequential operations require sequential master cycles. Although sequential operations require a deeper pipeline, 2-phase operation provides more time for the computational logic to execute. This allows the ’C24x to run at faster clock rates, despite having a deeper pipeline that imposes a penalty on branches and subroutine calls.

5.3 Branches, Calls, and Returns

Branches, calls, and returns break the sequential flow of instructions by transferring control to another location in program memory. A *branch* only transfers control to the new location. A *call* also saves the return address (the address of the instruction following the call) to the top of the hardware stack. Every called subroutine or interrupt service routine is concluded with a *return* instruction, which pops the return address off the stack and back into the program counter (PC).

The ’C24x has two types of branches, calls, and returns:

- **Unconditional.** An unconditional branch, call, or return is always executed. The unconditional branch, call, and return instructions are described in subsections 5.3.1, 5.3.2, and 5.3.3, respectively.
- **Conditional.** A conditional branch, call, or return is executed only if certain specified conditions are met. The conditional branch, call, and return instructions are described in detail in Section 5.4, *Conditional Branches, Calls, and Returns*, on page 5-10.

5.3.1 Unconditional Branches

When an unconditional branch is encountered, it is always executed. During the execution, the PC is loaded with the specified program-memory address and program execution begins at that address. The address loaded into the PC may come from either the second word of the branch instruction or the lower sixteen bits of the accumulator.

By the time the branch instruction reaches the execute phase of the pipeline, the next two instruction words have already been fetched. These two instruction words are flushed from the pipeline so that they are not executed, and then execution continues at the branched-to address. The unconditional branch instructions are B (branch) and BACC (branch to location specified by accumulator).

5.3.2 Unconditional Calls

When an unconditional call is encountered, it is always executed. When the call is executed, the PC is loaded with the specified program-memory address and program execution begins at that address. The address loaded into the PC may come from either the second word of the call instruction or the lower 16 bits of the accumulator. Before the PC is loaded, the return address is saved in the stack. After the subroutine or function is executed, a return instruction loads the PC with the return address from the stack, and execution resumes at the instruction following the call.

By the time the unconditional call instruction reaches the execute phase of the pipeline, the next two instruction words have already been fetched. These two instruction words are flushed from the pipeline so that they are not executed, the return address is stored to the stack, and then execution continues at the beginning of the called function. The unconditional call instructions are CALL and CALA (call subroutine at location specified by accumulator).

5.3.3 Unconditional Returns

When an unconditional return (RET) instruction is encountered, it is always executed. When the return is executed, the PC is loaded with the value at the top of the stack, and execution resumes at that address.

By the time the unconditional return instruction reaches the execute phase of the pipeline, the next two instruction words have already been fetched. The two instruction words are flushed from the pipeline so that they are not executed, the return address is taken from the stack, and then execution continues in the calling function.

5.4 Conditional Branches, Calls, and Returns

The ’C24x provides branch, call, and return instructions that execute only if one or more conditions are met. You specify the conditions as operands of the conditional instruction. Table 5–3 lists the conditions that you can use with these instructions and their corresponding operand symbols.

Table 5–3. Conditions for Conditional Calls and Returns

Operand Symbol	Condition	Description
EQ	ACC = 0	Accumulator equal to 0
NEQ	ACC ≠ 0	Accumulator not equal to 0
LT	ACC < 0	Accumulator less than 0
LEQ	ACC ≤ 0	Accumulator less than or equal to 0
GT	ACC > 0	Accumulator greater than 0
GEQ	ACC ≥ 0	Accumulator greater than or equal to 0
C	C = 1	Carry bit set to 1
NC	C = 0	Carry bit cleared to 0
OV	OV = 1	Accumulator overflow detected
NOV	OV = 0	No accumulator overflow detected
BIO	$\overline{\text{BIO}}$ low	$\overline{\text{BIO}}$ pin is low
TC	TC = 1	Test/control flag set to 1
NTC	TC = 0	Test/control flag cleared to 0

5.4.1 Using Multiple Conditions

Multiple conditions can be listed as operands of the conditional instructions. If multiple conditions are listed, all conditions must be met for the instruction to execute. Note that only certain combinations of conditions are meaningful. See Table 5–4. For each combination, the conditions must be selected from Group 1 and Group 2 as follows:

- **Group 1.** You can select up to two conditions. Each of these conditions must be from a different category (A or B); you cannot have two conditions from the same category. For example, you can test EQ and OV at the same time, but you cannot test GT and NEQ at the same time.

- **Group 2.** You can select up to three conditions. Each of these conditions must be from a different category (A, B, or C); you cannot have two conditions from the same category. For example, you can test TC, C, and BIO at the same time, but you cannot test C and NC at the same time.

Table 5–4. Groupings of Conditions

Group 1		Group 2		
Category A	**Category B**	**Category A**	**Category B**	**Category C**
EQ	OV	TC	C	BIO
NEQ	NOV	NTC	NC	
LT				
LEQ				
GT				
GEQ				

5.4.2 Stabilization of Conditions

A conditional instruction must be able to test the most recent values of the status bits. Therefore, the conditions cannot be considered stable until the fourth, or execution, stage of the pipeline, one cycle after the previous instruction has been executed. The pipeline controller stops the decoding of any instructions following the conditional instruction until the conditions are stable.

5.4.3 Conditional Branches

A branch instruction transfers program control to any location in program memory. Conditional branch instructions are executed only when one or more user-specified conditions are met (see Table 5–3 on page 5-10). If all the conditions are met, the PC is loaded with the second word of the branch instruction, which contains the address to branch to, and execution continues at this address.

By the time the conditions have been tested, the two instruction words following the conditional branch instruction have already been fetched in the pipeline. If all the conditions are met, these two instruction words are flushed from the pipeline so that they are not executed, and then execution continues at the branched-to address. If the conditions are *not* met, the two instruction words are executed instead of the branch. Because conditional branches use conditions determined by the execution of the previous instructions, a conditional branch takes one more cycle than an unconditional one.

The conditional branch instructions are BCND (branch conditionally) and BANZ (branch if currently selected auxiliary register is not equal to 0). The BANZ instruction is useful for implementing loops.

5.4.4 Conditional Calls

The conditional call (CC) instruction is executed only when the specified condition or conditions are met (see Table 5–3 on page 5-10). This allows your program to choose among multiple subroutines; based on the data being processed. If all the conditions are met, the PC is loaded with the second word of the call instruction, which contains the starting address of the subroutine. Before branching to the subroutine, the processor stores the address of the instruction following the call instruction—the return address—to the stack. The function must end with a return instruction, which takes the return address off the stack and forces the processor to resume execution of the calling program.

By the time the conditions of the conditional call instruction have been tested, the two instruction words following the call instruction have already been fetched in the pipeline. If all the conditions are met, these two instruction words are flushed from the pipeline so that they are not executed, and then execution continues at the beginning of the called function. If the conditions are *not* met, the two instructions are executed instead of the call. Because there is a wait cycle for conditions to become stable, the conditional call takes one more cycle than the unconditional one.

5.4.5 Conditional Returns

Returns are used in conjunction with calls and interrupts. A call or interrupt stores a return address to the stack and then transfers program control to a new location in program memory. The called subroutine or the interrupt service routine concludes with a return instruction, which pops the return address off the top of the stack and into the program counter (PC).

The conditional return instruction (RETC) is executed only when one or more conditions are met (see Table 5–3 on page 5-10). By using the RETC instruction, you can give a subroutine or interrupt service routine more than one possible return path. The path chosen then depends on the data being processed. In addition, you can use a conditional return to avoid conditionally branching to/around the return instruction at the end of the subroutine or interrupt service routine.

If all the conditions are met for execution of the RETC instruction, the processor loads the return address from the stack to the PC and resumes execution of the calling or interrupted program.

RETC, like RET, is a single-word instruction. However, because of the potential PC discontinuity, it operates with the same effective execution time as the conditional branch (BCND) and the conditional call (CC). By the time the conditions of the conditional return instruction have been tested, the two instruction words following the return instruction have already been fetched in the pipeline. If all the conditions are met, these two instruction words are flushed from the pipeline so that they are not executed, and then execution of the calling program continues. If the conditions are *not* met, the two instructions are executed instead of the return. Because there is a wait cycle for conditions to become stable, the conditional return takes one more cycle than the unconditional one.

5.5 Repeating a Single Instruction

The ’C24x repeat (RPT) instruction allows the execution of a single instruction N + 1 times, where N is specified as an operand of the RPT instruction. When RPT is executed, the repeat counter (RPTC) is loaded with N. RPTC is then decremented every time the repeated instruction is executed, until RPTC equals 0. RPTC can be used as a 16-bit counter when the count value is read from a data-memory location; if the count value is specified as a constant operand, it is in an 8-bit counter.

The repeat feature is useful with instructions such as NORM (normalize contents of accumulator), MACD (multiply and accumulate with data move), and SUBC (conditional subtract). When instructions are repeated, the address and data buses for program memory are free to fetch a second operand in parallel with the address and data buses for data memory. This allows instructions such as MACD and BLPD to effectively execute in a single cycle when repeated.

Chapter 6

System Functions

This chapter describes the device functions that are not specific to any peripheral:

- The peripheral interface transfers data between the CPU data bus and the peripheral bus, which is independent of the CPU.
- The system configuration registers provide software control and status information for functions that affect both the DSP core and certain peripherals.
- Hardware interrupts (including reset) require control by both CPU registers and peripheral registers.
- Power-down modes can affect both the CPU and the peripherals.

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6.1 Peripheral Interface

In order to support a large number of peripherals without compromising the electrical performance of the ’C2xx DSP CPU’s data bus, ’C24x devices have a separate peripheral bus which operates at a lower frequency than the CPU buses. Most peripherals are attached to this peripheral bus, although a few (for example, the event manager) interface directly to the CPU’s data bus. See the individual peripheral specifications and device data sheets for details. Up to 16 peripherals may be connected to the peripheral bus.

One of the functions of the peripheral interface is to interface the CPU to the peripheral bus.

The CPU is clocked at either two times (2× mode) or four times (4× mode) the clock rate of the peripheral bus. Because the peripheral bus runs slower than the CPU bus, peripheral bus reads and writes take multiple CPU cycles. The exact number of CPU cycles a peripheral access takes to complete depends on:

- Peripheral clock rate
- Phase of the peripheral clock in which the CPU initiates the peripheral access
- Type of access: read or write

Table 6–1 shows how many CPU clock cycles it takes to complete read and write accesses to peripherals connected to the peripheral bus. If, for example, the clocks are in 4x mode, a single peripheral read may take 5, 6, 7, or 8 CPU cycles, depending on the phase of the peripheral clock when the CPU initiates the peripheral access. If back-to-back accesses are performed, all accesses after the first will take eight cycles in 4× mode. Note that writes always take one cycle longer than reads. This is consistent with zero-wait-state external memory accesses and event manager accesses over the CPU’s data bus; these accesses take one cycle for a read and two cycles for a write.

Table 6–1. CPU Cycles to Complete Reads From and Writes to the Peripheral Bus

	2x Clock Mode		4x Clock Mode	
Type of Access	**Single Accesses (Cycles)**	**Back-to-Back Accesses (Cycles)**	**Single Accesses (Cycles)**	**Back-to-Back Accesses (Cycles)**
Read	3 or 4	4	5, 6, 7, or 8	8
Write	4 or 5	4	6, 7, 8, or 9	8

All CPU memory accesses are 16 bits wide. Reads from 8-bit peripherals are LSB aligned. The most significant eight bits of a write to an 8-bit peripheral are ignored. All peripherals are located in the CPU's data space; this allows the full instruction set to act upon the peripheral registers. I/O space is not used by on-chip peripherals. A 'C24x device can have no more than 16 on-chip peripherals attached to the peripheral bus. There is no such limitation on peripherals that interface directly to the CPU.

6.2 System Configuration Registers

The system configuration registers are shown in Figure 6–1 and described in subsections 6.2.1 through 6.2.3. Note these points about the register locations:

- All unimplemented (reserved) bits are read as indeterminate values (unless otherwise stated).
- Bit 0 of the peripheral address bus is not decoded; therefore, these 16-bit registers are accessible at each even address location and the following (odd) address location. For example, the register SYSIVR is nominally at location 701Eh, but it can also be accessed at address 701Fh.

Figure 6–1. System Configuration Registers

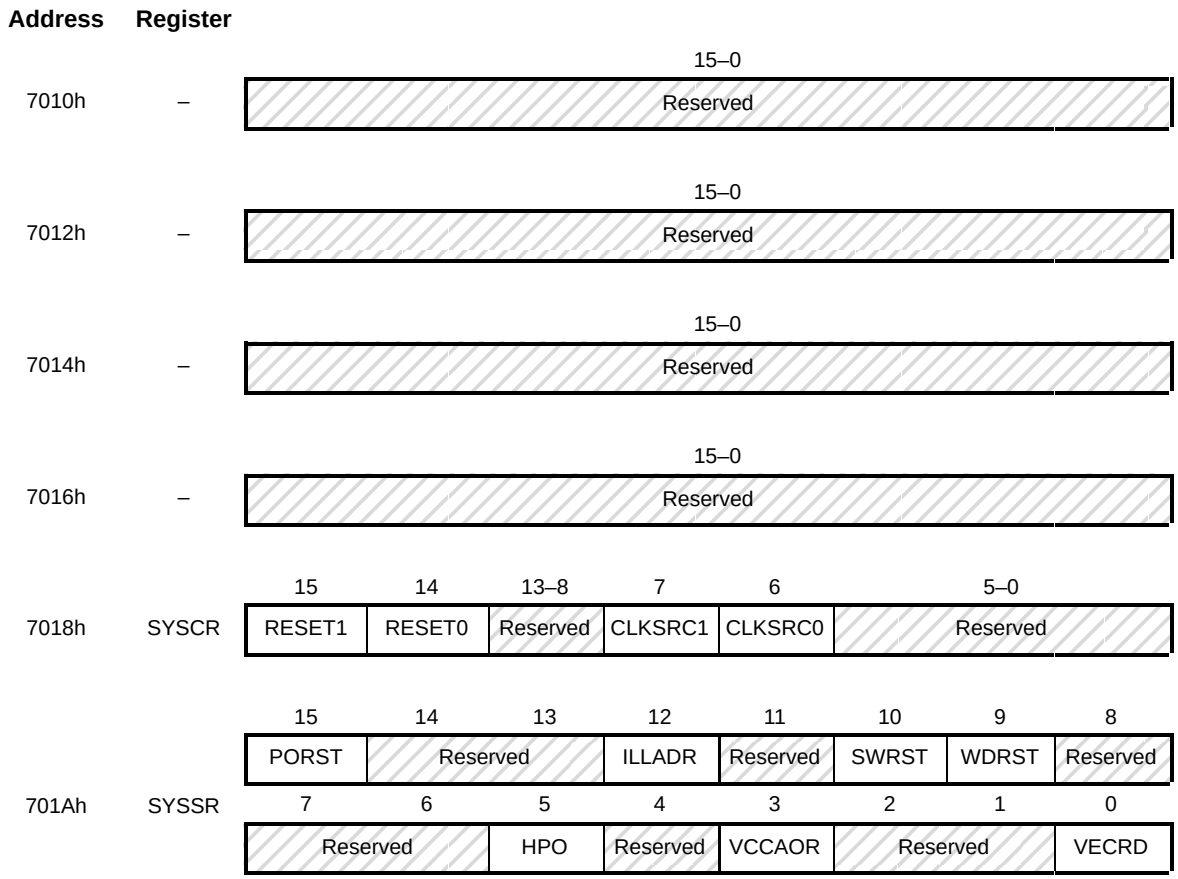


Figure 6–1. System Configuration Registers (Continued)

Address	Register	15–0
701Ch	–	Reserved
701Eh	SYSIVR	System interrupt vector register

6.2.1 System Control Register (SYSCR)

Figure 6–2. System Control Register (SYSCR) — Address 7018h

15	14	13–8	7	6	5–0
RESET1	RESET0	Reserved	CLKSRC1	CLKSRC0	Reserved
R/W–0	R/W–1		R/W–1*	R/W–1*	

Note: R = Read access, W = Write access, –n = Value after reset,
* = Not affected by reset, set to 1 by power-on reset

Bits 15–14 **RESET1, RESET0**. Software reset bits. These bits, which control the software reset function of the device, must be written to at the same time. Writing a 1 to RESET1 or a 0 to RESET0 causes a global reset to occur, as shown in the following table.

RESET1	RESET0	Resulting Action
0	0	Global reset
0	1	–
1	0	Global reset
1	1	Global reset

Bits 13–8 **Reserved**. Reads are indeterminate and writes have no effect.

Bits 7–6 **CLKSRC1, CLKSRC0**. CLKOUT-pin source select. These bits control the selection of the CLKOUT pin function.

CLKSRC1	CLKSRC0	CLKOUT Pin Function
0	0	Digital I/O mode (controlled by I/O register bits—see device data sheet).
0	1	WDCLK: Watchdog Timer Clock output mode (nominally 16 kHz).
1	0	SYSCLK: system clock.
1	1	CPUCLK: CPU clock output mode.

Bits 5–0 **Reserved**. Reads are indeterminate and writes have no effect.

6.2.2 System Status Register (SYSSR)

Bits 15, 12, 10 and 9 of the system status register indicate the cause of a reset. The reset service routine can read this register and use these bits to take the appropriate action according to the cause of reset. For example, if a power-on reset occurs, the clock module control registers may have to be reconfigured.

Figure 6–3. System Status Register (SYSSR) — Address 701Ah

15	14–13	12	11	10	9	8–6	5	4	3	2–1	0
PORST	Res	ILLADR	Res	SWRST	WDRST	Res	HPO	Res	VCCAOR	Res	VECRD
R/C–x		R/C–x		R/C–x	R/C–x		R/C–i		R–1		R–0

Note: R = Read access, C = Clear-only write access, –n = Value after reset (x means value unchanged by reset), –i = Value of V_{CCP} pin latch on rising edge of RESET

Bit 15 **PORST**. Power-on reset status bit. The occurrence of a power-on reset sets this bit. Depending on the device configuration, a power-on reset may be caused by either an on-chip low-voltage detect module indicating that V_{DD} is out of regulation, or from an off-chip power-on reset or low-voltage detection source connected to the PORST pin.

0 = No reset has occurred due to power-on reset or V_{DD} out of regulation.
1 = Reset due to power-on reset or V_{DD} out of regulation.

Bits 14–13 **Reserved**. Reads are indeterminate and writes have no effect.

Bit 12 **ILLADR**. Illegal-address reset status bit. Illegal address reset occurs when an unimplemented on-chip address location in data or program space is accessed. See the data sheet for each specific 'C24x device for details of which addresses are illegal on that device.

0 = No illegal address conditions
1 = Reset due to illegal address

Bit 11 **Reserved**. Reads are indeterminate and writes have no effect.

Bit 10 **SWRST**. Software reset status bit.

0 = No software reset
1 = Software reset occurred. (A 1 was written to bit 15 of the SYSCR, or a 0 was written to bit 14 of the SYSCR.)

Bit 9 **WDRST**. Watchdog reset status bit.

0 = No reset
1 = Reset due to Watchdog Timer overflow

Bits 8–6 **Reserved**. Reads are indeterminate and writes have no effect.

Bit 5 **HPO**. Hardware protect override. If the flash programming voltage pin (V_{CCP}) is at 5V on the trailing edge of the reset pin ($\overline{RS}$) and is held at that value, the HPO bit is set. (This only applies to 'F24x devices, which have on-chip flash EEPROM). This value is cleared by software or if the V_{CCP} pin level changes to 0V.

0 = Normal mode
1 = HPO mode: Flash EEPROM programming is enabled and the Watchdog can be disabled by setting the WDDIS bit in the WD control register. For details about this register, see *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices*.

Bit 4 **Reserved**. Reads are indeterminate and writes have no effect.

Bit 3 **VCCAOR**. Analog V_{CC} (V_{CCA}) out-of-regulation bit. This bit is only valid if the device has an on-chip low-voltage detect module.

0 = V_{CCA} is on and in regulation.
1 = V_{CCA} is off or out of regulation.

Bits 2–1 **Reserved**. Reads are indeterminate and writes have no effect.

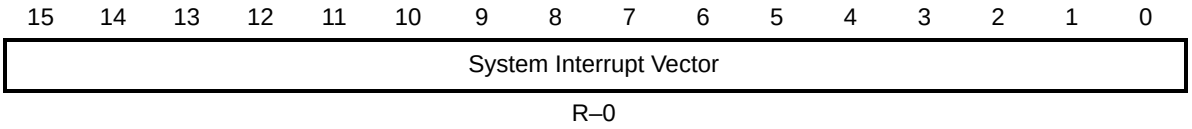
Bit 0 **VECRD**. Interrupt vector read pending bit. This bit is set when an interrupt vector is loaded into the SYSIVR (when the interrupt is acknowledged). It is cleared when the SYSIVR is read. This bit is used by the service routine of nonmaskable interrupt NMI (see subsection 6.3.7, *Programming an ISR for Nonmaskable Interrupt (NMI)* on page 6-30).

0 = No read of the interrupt vector register is pending.
1 = An interrupt vector has been latched but has not been read yet.

6.2.3 System Interrupt Vector Register (SYSIVR)

The system interrupt vector register is a read-only register.

Figure 6–4. System Interrupt Vector Register (SYSIVR) — Address 701Eh



Note: R = Read access, –n = Value after reset

Bits 15–8 Eight MSBs of the system interrupt vector. these bits are always read as 0s.

Bits 7–0 Eight LSBs of the system interrupt vector. These bits are loaded with the interrupt vector address offset value. This value is generated by a peripheral attached to the peripheral bus, in response to the acknowledgement of the corresponding maskable interrupt.

6.3 Interrupts

Interrupts are hardware- or software-driven signals that cause the ’C24x to suspend its main program and execute a subroutine. Typically, interrupts are generated by hardware devices that need to give data to or take data from the ’C24x (for example, A/D and D/A converters and other processors). Interrupts may also be used to signal that a particular event has taken place (for example, a timer has finished counting).

The ’C24x supports both software and hardware interrupts:

- A *software interrupt* is requested by an instruction (INTR, NMI, or TRAP).
- A *hardware interrupt* is requested by a signal from a physical device. Two types exist:
  - *External* hardware interrupts are triggered by signals at external interrupt pins. All these interrupts have programmable polarity and priority and are controlled by the external interrupt control registers.
  - *Internal* hardware interrupts are triggered by signals from the on-chip peripherals.

If hardware interrupts are triggered at the same time, the ’C24x services them according to a set priority ranking. Each of the ’C24x interrupts, whether hardware or software, can be placed in one of the following two categories:

- **Maskable interrupts.** These are hardware interrupts that can be blocked (masked) or enabled (unmasked) by software.
- **Nonmaskable interrupts.** These interrupts cannot be blocked. The ’C24x always responds to this type of interrupt and branches from the main program to a subroutine. The ’C24x nonmaskable interrupts include all software interrupts and two external hardware interrupts: reset ($\overline{RS}$) and NMI. Note that although $\overline{RS}$ is always active low, NMI has programmable polarity. For more information, see subsection 6.3.11, *External Interrupt Control Registers*, on page 6-37.

  For information about the reset signal and its effects on the ’C24x, see Section 6.4, *Reset Operation*, on page 6-48. The control register for NMI is described in subsection 6.3.2, *Nonmaskable Interrupt Operation*, on page 6-12.

Table 6–2 summarizes the interrupts available on the CPU. Other maskable interrupts are available through on-chip peripherals. The relationship between the maskable CPU interrupts (INT1–INT6) and the maskable peripheral interrupts is included in this section; however, for details on the peripheral interrupts available on a specific ’C24x device, see the data sheet for that device.

Table 6–2. 'C24x Interrupt Locations and Priorities

K†	Vector Location	Name	Priority	Function
0	0h	$\overline{RS}$	1 (highest)	Hardware reset (nonmaskable)
1	2h	INT1	4	Maskable interrupt level #1
2	4h	INT2	5	Maskable interrupt level #2
3	6h	INT3	6	Maskable interrupt level #3
4	8h	INT4	7	Maskable interrupt level #4
5	Ah	INT5	8	Maskable interrupt level #5
6	Ch	INT6	9	Maskable interrupt level #6
7	Eh		10	Reserved
8	10h	INT8	–	User-defined software interrupt
9	12h	INT9	–	User-defined software interrupt
10	14h	INT10	–	User-defined software interrupt
11	16h	INT11	–	User-defined software interrupt
12	18h	INT12	–	User-defined software interrupt
13	1Ah	INT13	–	User-defined software interrupt
14	1Ch	INT14	–	User-defined software interrupt
15	1Eh	INT15	–	User-defined software interrupt
16	20h	INT16	–	User-defined software interrupt
17	22h	TRAP	–	TRAP instruction vector
18	24h	NMI	3	Nonmaskable interrupt
19	26h		2	Reserved
20	28h	INT20	–	User-defined software interrupt
21	2Ah	INT21	–	User-defined software interrupt
22	2Ch	INT22	–	User-defined software interrupt
23	2Eh	INT23	–	User-defined software interrupt

† The K value is the operand used in an INTR instruction that branches to the corresponding interrupt vector location.

Table 6–2. 'C24x Interrupt Locations and Priorities (Continued)

K†	Vector Location	Name	Priority	Function
24	30h	INT24	–	User-defined software interrupt
25	32h	INT25	–	User-defined software interrupt
26	34h	INT26	–	User-defined software interrupt
27	36h	INT27	–	User-defined software interrupt
28	38h	INT28	–	User-defined software interrupt
29	3Ah	INT29	–	User-defined software interrupt
30	3Ch	INT30	–	User-defined software interrupt
31	3Eh	INT31	–	User-defined software interrupt

† The K value is the operand used in an INTR instruction that branches to the corresponding interrupt vector location.

6.3.1 Interrupt Operation: Three Phases

The 'C24x handles interrupts in three main phases:

1) **Receive interrupt request.** Suspension of the main program must be requested by software (program code) or hardware (a pin or an on-chip device).

2) **Acknowledge interrupt.** The 'C24x must acknowledge the interrupt request. If the interrupt is maskable, certain conditions must be met in order for the 'C24x to acknowledge it. For nonmaskable hardware interrupts and for software interrupts, acknowledgment is immediate.

3) **Execute interrupt service routine.** Once the interrupt is acknowledged, the 'C24x branches to its corresponding subroutine, called an interrupt service routine (ISR). The 'C24x follows the branch instruction you place at a predetermined address (the vector location) and executes the ISR you have written.

6.3.2 Nonmaskable Interrupt Operation

Hardware nonmaskable interrupts can be requested through two pins:

- **$\overline{\text{RS}}$ (reset).** This interrupt stops program flow, returns the processor to a predetermined state, and then begins program execution at address 0000h. For details of the reset operation, see Section 6.4, *Reset Operation*, on page 6-48.

- **NMI.** This interrupt is used as a soft reset. Unlike a hardware reset, NMI neither affects any of the modes of the device nor aborts a currently active instruction or memory operation. Although NMI uses the same logic as the maskable interrupts, it is not maskable. NMI happens regardless of the value of the INTM bit, and there is no mask bit for NMI. This interrupt can only be locked out by an already executing NMI or a reset. When NMI is activated (either by the NMI pin or by the NMI instruction), the processor switches program control to vector location 24h. In addition, maskable interrupts are disabled: the INTM bit of status register ST0 is set to 1.

  Note that although $\overline{\text{RS}}$ is always active low, NMI has programmable polarity. For more information, see subsection 6.3.11, *External Interrupt Control Registers*, on page 6-37.

Software interrupts (which are inherently nonmaskable) are requested by the following instructions:

- **INTR.** This instruction allows you to initiate any 'C24x interrupt, including user-defined interrupts INT8–INT16 and INT20–INT31. The instruction operand (K) indicates which interrupt vector location the CPU branches to. Table 6–2 (page 6-10) shows the operand K that corresponds to each vector location. When an INTR interrupt is acknowledged, the interrupt mode (INTM) bit of status register ST1 is set to 1 to disable maskable interrupts.

- **NMI.** This instruction forces a branch to interrupt vector location 24h, the same location used for the nonmaskable hardware interrupt NMI. Thus, you can either initiate NMI by driving the NMI pin active or by executing an NMI instruction. When the NMI instruction is executed, INTM is set to 1 to disable maskable interrupts.

- **TRAP.** This instruction forces the CPU to branch to interrupt vector location 22h. The TRAP instruction does not disable maskable interrupts (INTM is not set to 1); thus when the CPU branches to the interrupt service routine, that routine can be interrupted by the maskable hardware interrupts.

After acknowledging a nonmaskable interrupt, the CPU:

1) Stores the program counter (PC) value (the return address) to the top of the hardware stack.

2) Loads the PC with the address of the interrupt vector.

3) Fetches the branch instruction that you stored at the vector location.

   If the interrupt was a hardware interrupt or was requested by either the INTR or NMI instructions, the CPU also sets the INTM bit to 1 to disable maskable interrupts.

4) Executes the branch, which leads it to the address of your ISR.

5) Executes the ISR until a return instruction concludes the ISR.

   If INTM is 1, all maskable interrupts are disabled during the execution of the ISR.

6) Pops the return address off the stack and into the PC.

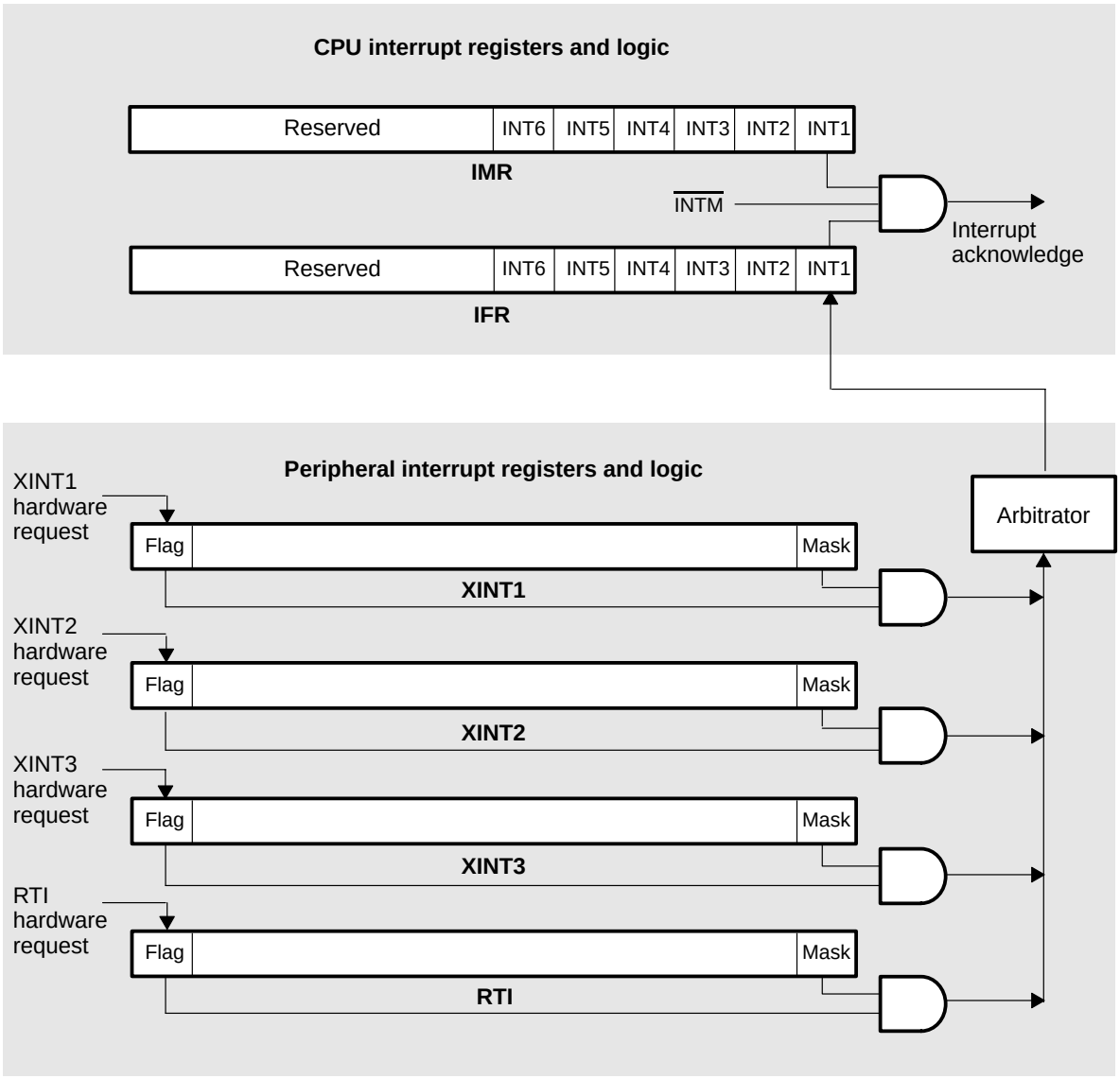
7) Continues executing the main program.

To determine which vector address has been assigned to each of the interrupts, refer to Table 6–2 (page 6-10). Each interrupt address has been spaced apart by two locations so that 2-word branch instructions can be accommodated in those locations.

6.3.3 Maskable Interrupt Structure

The CPU provides six maskable interrupt levels. Because a ’C24x device can have more than six maskable interrupt sources, each of the six interrupt levels can be shared by multiple interrupt sources. Figure 6–5 illustrates the structure used for receiving and acknowledging maskable interrupts. The figure shows four interrupt sources (XINT1, XINT2, XINT3, and RTI) sharing the interrupt level INT1. A similar situation exists for the other levels (INT2–INT6).

Figure 6–5. Example of Maskable Interrupt Structure



Path of a Maskable Interrupt Request

Each of the interrupt sources has its own control register with a flag bit and a mask bit (see subsection 6.3.12, *Type A, Type B, and Type C Interrupt Pins*, on page 6-39). When an interrupt signal is received, the flag bit in the corresponding control register is set, indicating that the interrupt has been requested. If the mask bit is also set, a signal is sent to arbitration logic, which may simultaneously receive similar signals from one or more other control registers. The arbitration logic compares the priority level of competing interrupt requests, and it passes the interrupt of highest priority to the CPU. The interrupt flag in the CPU's IFR that corresponds to the interrupt priority level on which the request was received is set. This indicates that the interrupt is pending. If the corresponding IMR bit is 1 and the INTM bit is 0, the CPU acknowledges the interrupt and executes the interrupt service routine (ISR).

Priorities of the Maskable Interrupts

All hardware interrupts are given a priority rank from 1 to 10 (1 being highest). The priorities are shown in Table 6–2 on page 6-10. When more than one of these hardware interrupts is pending, the interrupt of highest rank gets serviced first. The others are serviced in priority order after that. The maskable interrupt levels of the DSP core have the priorities shown in Table 6–3.

Table 6–3. Priorities of the Maskable Interrupt Levels in the DSP Core

Maskable Interrupt Level	Priority in the DSP Core
INT1	4
INT2	5
INT3	6
INT4	7
INT5	8
INT6	9

As an example of how the priority ranking is carried out, suppose INT1 and INT2 were both pending and not masked. INT1 would be acknowledged first, followed by INT2.

Each maskable interrupt level (INT1–INT6) is connected to multiple maskable interrupt sources, which also have set priority ranks. The source with highest priority has its interrupt request sent to the interrupt level first. Consider the interrupt sources in the example in Figure 6–5: XINT1, XINT2, XINT3, and RTI. Suppose they have the priority ranks with respect to INT1 listed in Table 6–4.

Table 6–4. Priority Ranking Under INT1

Maskable Interrupt Source	Priority Under INT1
XINT1	1
XINT2	2
XINT3	3
RTI	4

If all these sources had generated interrupt requests at the same time, XINT1 would get serviced first, then XINT2, followed by XINT3, and finally RTI.

6.3.4 CPU Interrupt Registers

There are two CPU registers for controlling interrupts:

- The interrupt flag register (IFR) contains flag bits that indicate when maskable interrupt requests have reached the CPU on levels INT1 through INT6.
- The interrupt mask register (IMR) contains mask bits that enable or disable each of the interrupt levels (INT1 through INT6).

Interrupt Flag Register (IFR)

The interrupt flag register (IFR), a 16-bit, memory-mapped register at address 0006h in data-memory space, is used to identify and clear pending interrupts. The IFR contains flag bits for all the maskable interrupts.

When a maskable interrupt is requested, the flag bit in the corresponding control register is set to 1. If the mask bit in that same control register is also 1, the interrupt request is sent to the CPU, setting the corresponding flag in the IFR. This indicates that the interrupt is pending or waiting for acknowledgement.

You can read the IFR to identify pending interrupts and write to the IFR to clear pending interrupts. To clear a single interrupt, write a 1 to the corresponding

IFR bit. All pending interrupts can be cleared by writing the current contents of the IFR back into the IFR. A device reset clears all IFR bits.

The following events also clear an IFR flag:

- ❑ The CPU acknowledges the interrupt.
- ❑ The 'C24x is reset.

Notes:

1) To clear an IFR bit, you must write a 1 to it, not a 0.

2) When a maskable interrupt is acknowledged, *only* the IFR bit is cleared automatically. The flag bit in the corresponding control register is *not* cleared. If an application requires that the control register flag be cleared, the bit must be cleared by software.

3) When an interrupt is requested by an INTR instruction and the corresponding IFR bit is set, the CPU does not clear the bit automatically. If an application requires that the IFR bit be cleared, the bit must be cleared by software.

The IFR is shown in Figure 6–6, and descriptions of the bits follow the figure.

Figure 6–6. Interrupt Flag Register (IFR) — Address 0006h

15–6	5	4	3	2	1	0
Reserved	INT6	INT5	INT4	INT3	INT2	INT1
0	R/W1C–0	R/W1C–0	R/W1C–0	R/W1C–0	R/W1C–0	R/W1C–0

Note: 0 = Always read as zeros, R = Read access, W1C = Write 1 to this bit to clear it, –n = Value after reset

Bits 15–6 **Reserved**. These bits are always read as 0s.

Bit 5 **INT6.** Interrupt 6 flag. This bit is the flag for interrupts connected to interrupt level INT6.

0 = No INT6 interrupt is pending.
1 = At least one INT6 interrupt is pending. Write a 1 to this bit to clear it to 0 and clear the interrupt request.

Bit 4 **INT5.** Interrupt 5 flag. This bit is the flag for interrupts connected to interrupt level INT5.

0 = No INT5 interrupt is pending.
1 = At least one INT5 interrupt is pending. Write a 1 to this bit to clear it to 0 and clear the interrupt request.

Bit 3 **INT4.** Interrupt 4 flag. This bit is the flag for interrupts connected to interrupt level INT4.

0 = No INT4 interrupt is pending.
1 = At least one INT4 interrupt is pending. Write a 1 to this bit to clear it to 0 and clear the interrupt request.

Bit 2 **INT3.** Interrupt 3 flag. This bit is the flag for interrupts connected to interrupt level INT3.

0 = No INT3 interrupt is pending.
1 = At least one INT3 interrupt is pending. Write a 1 to this bit to clear it to 0 and clear the interrupt request.

Bit 1 **INT2.** Interrupt 2 flag. This bit is the flag for interrupts connected to interrupt level INT2.

0 = No INT2 interrupt is pending.
1 = At least one INT2 interrupt is pending. Write a 1 to this bit to clear it to 0 and clear the interrupt request.

Bit 0 **INT1.** Interrupt 1 flag. This bit is the flag for interrupts connected to interrupt level INT1.

0 = No INT1 interrupt is pending.
1 = At least one INT1 interrupt is pending. Write a 1 to this bit to clear it to 0 and clear the interrupt request.

Interrupt Mask Register (IMR)

The IMR is a 16-bit, memory-mapped register located at address 0004h in data memory space. The IMR contains mask bits for all the maskable interrupt levels (INT1–INT6). Neither NMI nor $\overline{RS}$ is included in the IMR; thus, IMR has no effect on these interrupts.

You can read the IMR to identify masked or unmasked interrupt levels, and you can write to the IMR to mask or unmask interrupt levels. To unmask an interrupt level, set its corresponding IMR bit to 1. To mask an interrupt level, set its corresponding IMR bit to 0. When an interrupt is masked, it is not acknowledged, regardless of the value of the INTM bit. When an interrupt is unmasked, it is acknowledged if the corresponding IFR bit is 1 and the INTM bit is 0. At reset, the IMR bits are all set to 0, masking all the maskable interrupts.

The IMR is shown in Figure 6–7, and descriptions of the bits follow the figure.

Figure 6–7. Interrupt Mask Register (IMR) — Address 0004h

15–6	5	4	3	2	1	0
Reserved	INT6	INT5	INT4	INT3	INT2	INT1
0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0

Note: 0 = Always read as zeros, R = Read access, W = Write access, –n = Value after reset

Bits 15–6 **Reserved**. These bits are always read as 0s.

Bit 5 **INT6.** Interrupt 6 mask. This bit masks or unmasks interrupt level INT6.

0 = Level INT6 is masked.
1 = Level INT6 is unmasked.

Bit 4 **INT5.** Interrupt 5 mask. This bit masks or unmasks interrupt level INT5.

0 = Level INT5 is masked.
1 = Level INT5 is unmasked.

Bit 3 **INT4.** Interrupt 4 mask. This bit masks or unmasks interrupt level INT4.

0 = Level INT4 is masked.
1 = Level INT4 is unmasked.

Bit 2 **INT3.** Interrupt 3 mask. This bit masks or unmasks interrupt level INT3.

0 = Level INT3 is masked.
1 = Level INT3 is unmasked.

Bit 1 **INT2.** Interrupt 2 mask. This bit masks or unmasks interrupt level INT2.

0 = Level INT2 is masked.
1 = Level INT2 is unmasked.

Bit 0 **INT1.** Interrupt 1 mask. This bit masks or unmasks interrupt level INT1.

0 = Level INT1 is masked.
1 = Level INT1 is unmasked.

6.3.5 Maskable Interrupt Acknowledgement and Servicing

After an interrupt has been requested by hardware or software, the CPU must decide whether to acknowledge the request. If the CPU acknowledges the interrupt, the CPU executes its ISR.

Acknowledging Maskable Interrupts

Software interrupts and nonmaskable hardware interrupts are acknowledged immediately. Maskable hardware interrupts are acknowledged only after certain conditions are met:

- **Priority is highest.** When more than one hardware interrupt is requested at the same time, the 'C24x services each interrupt according to a set priority ranking (see Table 6–2 on page 6-10).
- **INTM bit is 0.** The interrupt mode (INTM) bit, bit 9 of status register ST0, enables or disables all maskable interrupts:
  - When INTM = 0, all unmasked interrupts are enabled.
  - When INTM = 1, all maskable interrupts are disabled.

  INTM is set to 1 automatically when the CPU acknowledges an interrupt (except when initiated by the TRAP instruction). INTM can also be set to 1 by a hardware reset or by execution of a disable-interrupts instruction (SETC INTM). INTM is reset to 0 by executing the enable-interrupts instruction (CLRC INTM). INTM has no effect on reset, NMI, or software-interrupts. Also, INTM is unaffected by the LST (load status register) instruction.

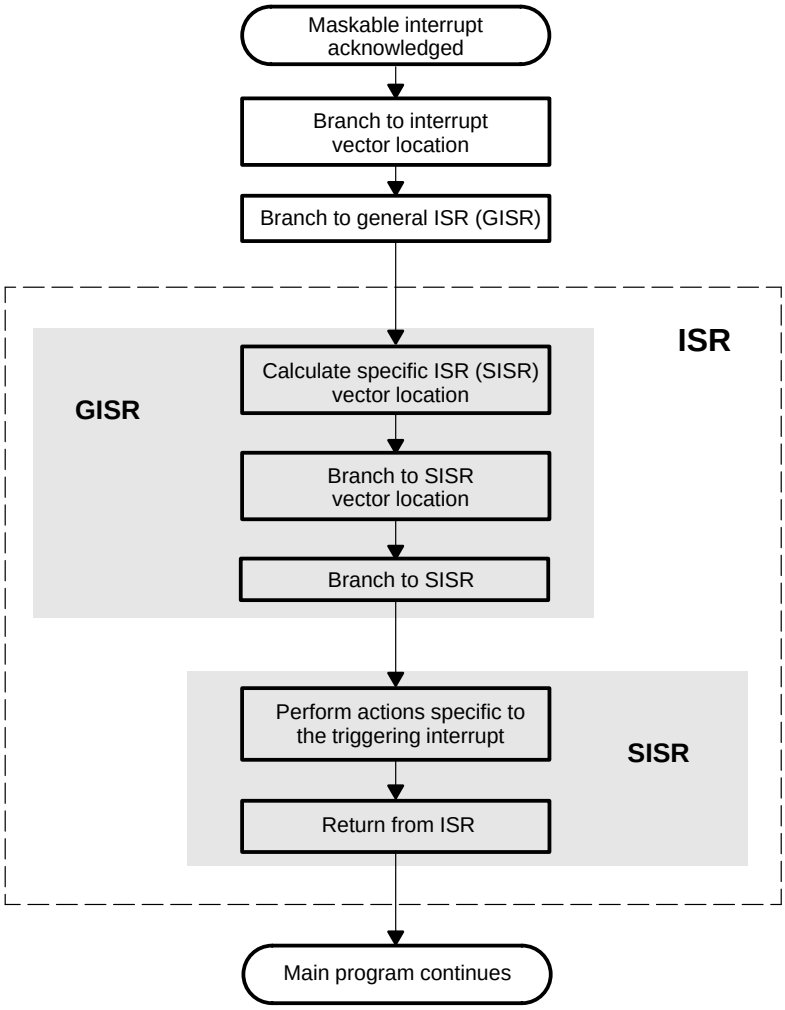
  INTM does not modify the interrupt mask register (IMR) or the interrupt flag register (IFR).
- **IMR mask bit is 1.** Each of the maskable interrupt levels has a mask bit in the IMR. To unmask an interrupt level, set its IMR bit to 1. For more details about the IMR, see *Interrupt Mask Register (IMR)* on page 6-18.

When the CPU acknowledges a maskable hardware interrupt, it jams the instruction bus with the INTR instruction. This instruction forces the PC to the appropriate address from which the CPU fetches the software vector.

Two-Part Interrupt Service Routine (ISR)

The ’C2xx CPU has six interrupt levels, but ’C24x devices provide a means of creating more than six ISRs. Figure 6–8 illustrates the process of servicing an interrupt request. For each of the six interrupt levels, the CPU branches to a corresponding general interrupt service routine (GISR). For example, when an interrupt request on priority level INT1 is responded to, the CPU branches to and executes GISR1. The GISR, after performing any necessary context saves, will identify and then branch to the specific interrupt service routine (SISR). The SISR performs the actions specific to the triggering interrupt and then returns program control to the interrupted program sequence.

Figure 6–8. Interrupt Service Routine Flow Chart



Branching to the GISR. Table 6–5 shows the addresses and contents of the interrupt vector locations for the maskable interrupt levels (for a table of all the interrupt vector locations, see Table 6–2 on page 6-10). When an interrupt is acknowledged through one of these interrupt levels, the CPU branches to the corresponding vector address and follows the branch at that address to the GISR. For example, if an interrupt is acknowledged through INT3, the program counter (PC) value is stored to the stack, and then the PC is loaded with program-memory address 0006h. Locations 0006h and 0007h contain a branch instruction that takes the CPU to the GISR.

Table 6–5. 'C24x Maskable Interrupt Vector Table

Interrupt Level	Interrupt Vector Location	Contents of Vector Location
INT1	0002h	Branch to GISR1
INT2	0004h	Branch to GISR2
INT3	0006h	Branch to GISR3
INT4	0008h	Branch to GISR4
INT5	000Ah	Branch to GISR5
INT6	000Ch	Branch to GISR6

Branching to the SISR. When a peripheral interrupt request is acknowledged (this includes the external interrupt control registers), the peripheral generates a vector address offset (vector) that corresponds to this interrupt event. This vector is usually latched in the system interrupt vector register (SYSIVR), although some peripherals, including the event manager, may keep the vector in a register in the peripheral. For details on the location of the interrupt vector register (IVR) for each interrupt level and also the vector value for each interrupt event, see the device data sheet. The GISR must read the value stored in the IVR and use it to generate the branch target address to the SISR. Several methods of doing this are described in subsection 6.3.8, *Additional Tasks of ISRs*, on page 6-31.

Interrupt Vectors

Table 6–6 shows an edited example of an interrupt vector table for a 'C24x device. Note that most peripherals use the system interrupt vector register (SYSIVR) to fetch the vector address offset, but the event manager has its own register for each of its three interrupt priority level requests:

- EVIVRA for interrupts in Group A
- EVIVRB for interrupts in Group B
- EVIVRC for interrupts in Group C

Table 6–6. Example Interrupt Locations and Priorities

Interrupt Name	Overall Priority	DSP-Core Interrupt and Address	Peripheral Vector Address	Peripheral Vector Address Offset	Maskable?	Source	Function
$\overline{\text{RS}}$	1 Highest	$\overline{\text{RS}}$ 0000h	N/A		No	Pin, S/W, Watchdog Timer, etc.	External, system reset (RESET)
RESERVED	2	0026h	N/A	N/A	No	Core	Emulator trap
NMI	3	NMI 0024h	N/A	N/A	No	External pin	External user interrupt
XINT1	4	INT1 0002h	SYSIVR	0002h	Yes	External pins	High-priority external user interrupts
XINT2	5	(System)	701Eh	0011h			
XINT3	6			001Fh			
RTI	10			0010h	Yes	Watchdog Timer	Watchdog timer interrupt
$\overline{\text{PDPINT}}$	11	INT2 0004h (EV INTA)	EVIVRA 7432h	0020h	Yes	External pin	Power drive protection interrupt
CMP1INT	12	(Group A)		0021h	Yes	Event Manager	Compare 1 interrupt
etc.							
TPINT2	22	INT3 0006h (EV INTB)	EVIVRB 7434h	002Bh	Yes	Event Manager	Timer 2 period interrupt
TCINT2	23	(Group B)		002Ch	Yes	Event Manager	Timer 2 compare interrupt
etc.							

Table 6–6. Example Interrupt Locations and Priorities (Continued)

Interrupt Name	Overall Priority	DSP-Core Interrupt and Address	Peripheral Vector Address	Peripheral Vector Address Offset	Maskable?	Source	Function
CAPINT1	30	INT4 0008h (EV INTC)	EVIVRC 7431h	0033h	Yes	Event Manager	Capture 1 interrupt
CAPINT2	31	(Group C)		0034h	Yes	Event Manager	Capture 2 interrupt
CAPINT3	32			0035h	Yes	Event Manager	Capture 3 interrupt
CAPINT4	33			0036h	Yes	Event Manager	Capture 4 interrupt
etc.							
ADCINT	37	INT6 000Ch (System)	SYSIVR 701Eh	0004h	Yes	ADC	Analog-to-Digital Converter interrupt
XINT1	38			0002h	Yes	External pins	Low-priority external user interrupts
XINT2	39			0011h	Yes		
XINT3	40			001Fh	Yes		
RESERVED	41	– 000Eh	N/A		Yes	DSP core	Used for analysis
TRAP	N/A	– 0022h	N/A		N/A		TRAP instruction vector

Phantom Interrupt Vector

The phantom interrupt vector is an interrupt system-integrity feature. In the event that an interrupt is acknowledged but no peripheral responds by loading an interrupt vector address offset value into the interrupt vector register (IVR), the phantom vector (0000h) is loaded into the IVR instead so that this fault can be handled in a controlled manner.

Two causes of phantom interrupts are:

- Execution of the INTR instruction with an argument in the range of 1 to 6. This is a software request to service one of the six maskable interrupt levels (INT1, INT2, INT3, INT4, INT5, or INT6).
- A glitch on an interrupt request line.

In either case, when the interrupt is acknowledged, no peripheral loads a vector into the IVR. Loading the IVR with the phantom interrupt vector ensures that the DSP branches to a known location.

6.3.6 Programming ISRs for Maskable Interrupts

This subsection lists three methods for creating an interrupt service routine (ISR) for a maskable interrupt:

- **Method 1:** A typical ISR
- **Method 2:** An ISR that is designed to reduce the latency between the time the interrupt is requested and the time the event-specific portion of the ISR is executed
- **Method 3:** An ISR that has very little latency. It can be used only if one peripheral interrupt source is connected to an interrupt request priority level or if only one of many interrupt sources connected to an interrupt request priority level is enabled.

Method 1: Typical ISR

The method described here is the typical implementation of an ISR for the ’C24x. Methods 2 and 3, described below are slightly more difficult to program. It is best to implement this typical ISR first and then develop a more complex routine if the latency for certain events is too high.

In Method 1, the ISR for a maskable interrupt is divided into two segments:

1) **General ISR (GISR).** When an interrupt on one of the six maskable interrupt levels (INT1–INT6) is acknowledged, the GISR reads the relevant interrupt vector register (IVR), shifts the value left by one bit, adds an offset to the value, and then branches to a peripheral interrupt vector table. From this table, it fetches and executes the appropriate branch to the specific ISR (SISR).
2) **Specific ISR (SISR).** The SISR performs actions specific to the event that caused the interrupt and then returns program control to the interrupted code sequence.

Table 6–7 shows an example of Method 1.

Table 6–7. Example of Method 1 ISR

Cycle Count	Address	Assembly Language Code		
		;CPU interrupt vector table		
	...		...	
0	0006h	INT3	B	;Branch to address of general ISR
	0007h		GISR3	;for INT3.
	...		...	
4	GISR3 Addr	GISR3	LACC xIVR,1	;Load accumulator with ;contents of interrupt vector ;register (xIVR) shifted by 1.
4+n†	GISR3 Addr+1		ADD offset	;Add offset to accumulator.
5+n†	GISR3 Addr+2		BACC	;Branch to address in accumulator ;(2*IVR+offset).
	...		...	
		;Peripheral interrupt vector table		
	...		...	
9+n†	2*IVR+offset		B	;Branch to specific ISR for
	(2*IVR+offset)+1		SISRx	;event that requested interrupt.
	...		...	
13+n†	SISRx Addr	SISRx	...	;Perform event–specific actions.
	SISRx Addr+1		...	;...
	...		...	;...
	SISRx Addr+n		RET	;Return from ISR.

† For the peripheral interface, n = 2; for the event manager, n = 1

In the example, the following events take place:

1) An acknowledged peripheral interrupt asserts INT3 and causes the device to enter its interrupt vector table at address 0006h.

2) From addresses 0006h and 0007h, the device reads a branch that leads it to GISR3.

3) GISR3 loads the accumulator with the contents of the relevant IVR shifted by 1 (that is, multiplied by 2). Note the following:

   ❑ The LACC instruction uses direct addressing to access the IVR value. For this to work, the data page pointer (DP in status register ST0) must be set to point to the page in data memory that contains the IVR.

- ❑ You may want to save the accumulator value before loading the accumulator with the shifted IVR value.
- ❑ The incoming vector has been multiplied by two because a peripheral interrupt vector table must support a 2-word branch instruction for each peripheral interrupt.

4) GISR3 adds to the accumulator an offset that corresponds to the start of the peripheral interrupt vector table. The accumulator now contains the address that needs to be accessed in the peripheral interrupt vector table.

5) GISR3 branches to the address in the accumulator.

6) From the peripheral interrupt vector location, the device reads a branch that leads to the SISR for the event that requested the interrupt.

7) The SISR is executed. The SISR concludes with a return instruction, which returns program control to the interrupted code sequence.

Method 2: Minimum Latency ISR for Multiple Events per Interrupt Level

This method is similar to Method 1, but has reduced latency because one of the branches (the branch to the peripheral interrupt vector table) is bypassed. Instead, the general ISR (GISR) branches directly to the specific ISR (SISR) To do this, the GISR shifts the value from the IVR by more than 2 and uses the result as the branch target. It may be difficult to locate all the SISRs in program space without leaving some holes in memory. It is best to use Method 1 initially and then Method 2 for some interrupts if the higher latency of Method 1 is unacceptable.

The Method 2 ISR is implemented as follows:

1) **General ISR (GISR).** When an interrupt on one of the six maskable interrupt levels (INT1–INT6) is acknowledged, the GISR reads the relevant interrupt vector register (IVR), shifts the value left by a predetermined amount, and then branches to the specific ISR (SISR). The shift amount is chosen such that the accumulator will contain the address of the SISR. For example, if three interrupt sources were tied to INT2, and the SISR for each of the events were not greater than 16 words long, then a shift of 4 would be suitable.

2) **Specific ISR (SISR).** The SISR performs actions specific to the event that caused the interrupt and then returns program control to the interrupted code sequence.

Table 6–8 shows an example of Method 2.

Table 6–8. Example of Method 2 ISR

Cycle Count	Address	Assembly Language Code		
		;CPU interrupt vector table		
	...		...	
0	0004h	INT2	B	;Branch to GISR
	0005h		GISR2	;for INT2.
	...		...	
4	GISR2 Addr	GISR2	LACC xIVR,shift	;Load accumulator with ;contents of interrupt vector ;register (xIVR) shifted by ;an amount that will result in ;SISRx address in accumulator.
4+n†	GISR2 Addr+1		BACC	;Branch to address in ;accumulator.
	...		...	
8+n†	SISRx Addr	SISRx	...	;Perform event–specific actions.
	SISRx Addr+1		...	;...
	...		...	;...
	SISRx Addr+n		RET	;Return from ISR.

† For the peripheral interface, n = 2; for the event manager, n = 1.

In the example, the following events take place:

1) An acknowledged peripheral interrupt asserts INT2 and causes the device to enter it's interrupt vector table at address 0004h.

2) From addresses 0004h and 0005h, the device reads a branch that leads it to GISR2.

3) GISR2 loads the accumulator with the contents of the relevant IVR shifted by an amount that results in the accumulator's holding the address of the SISR. Note the following:

   - ❏ The LACC instruction uses direct addressing to access the IVR value. For this to work, the data page pointer (DP in status register ST0) must be set to point to the page in data memory that contains the IVR.
   - ❏ You may want to save the accumulator value before loading the accumulator with the shifted IVR value.

4) GISR2 branches to the address in the accumulator (the start address of the SISR).

5) The SISR is executed. The SISR concludes with a return instruction, which returns program control to the interrupted code sequence.

Method 3: ISR for Single Event per Interrupt Level

A device can be configured such that only one event can assert a particular maskable interrupt.

Either: Both of these conditions must be met:

- ❑ Only one peripheral is connected to the maskable interrupt level (INT1, INT2, INT3, INT4, INT5, or INT6).
- ❑ That one peripheral has only one event that can cause an interrupt request (and, thus, has only one interrupt vector).

Or: Only one of the many events tied to a particular interrupt priority level would be enabled.

In either situation, there is no need for a 2-part ISR like the one in Method 1 or the one in Method 2. There is one simple ISR and, thus, only one branch instruction. The address of the ISR is known; it does not have to be calculated in a routine. The GISR and the SISR become one and the same.

Table 6–9 shows an example of Method 3. In the example, the following events take place:

1) An acknowledged peripheral interrupt asserts INT1 and causes the device to enter its interrupt vector table at address 0002h.

2) From addresses 0002h and 0003h, the device reads a branch that leads it directly to the SISR.

3) The SISR is executed and concludes with a return instruction, which returns program control to the interrupted code sequence.

Table 6–9. Example of Method 3 ISR

Cycle Count	Address	Assembly Language Code		
		;CPU interrupt vector table		
	...		...	
0	0002h	INT1	B	;Branch to ISR1.
	0003h		ISR1	
	...		...	
4	SISRx Addr	ISR1	...	;Perform event–specific actions.
	SISRx Addr+1		...	;...
	...		...	;...
	SISRx Addr+n		RET	;Return from ISR.

6.3.7 Programming an ISR for Nonmaskable Interrupt (NMI)

Generally, there cannot be more than one event capable of generating an NMI request, and in cases where they do, they all share the same ISR. Thus, NMI does not require loading an interrupt vector register (IVR), and the branch at vector location 24h is the only branch necessary for the ISR.

An NMI can interrupt a maskable ISR after the vector has been loaded into the IVR but before the maskable ISR has read the IVR. Because NMI does not cause a loading of the IVR, an NMI does not cause overwriting of the maskable interrupt's vector address offset. However, an enabled interrupt immediately following the NMI ISR could overwrite the value in the IVR before the interrupted ISR has read it. Therefore, the ISR for NMI must check the VECRD bit (bit 0) in the system status register (SSR). If VECRD = 1, a read of the IVR is pending, and the NMI ISR should not reenable maskable interrupts.

The ISR for NMI could be implemented as shown in Table 6–10.

Table 6–10. One Implementation of an ISR for NMI

Address	Assembly Language Code			
	; CPU interrupt vector table			
...		...		
0024h	NMI	B		;Branch to NMI ISR.
0025h		NMI_ISR		
...		...		
NMI ISR Addr	NMI_ISR	...		;Start of ISR
...		...		;Body of ISR
...		BIT	SSR,15	;Test bit 0 (VECRD) of SYSSR0
...		RETC	TC	;If set, return from ISR
...				;(without enabling interrupts)
...		CLRC	INTM	;If not set, enable interrupts and
...		RET		;return from ISR.

6.3.8 Additional Tasks of ISRs

While performing the tasks requested by an interrupt, the ISR may also be:

- ❏ Saving and restoring register values
- ❏ Managing ISRs within ISRs

Saving and Restoring Register Values

Only the incremented program counter value is stored automatically before the CPU enters an ISR. You must design the ISR to save and then restore any other important register values or control bit values. You can use a common routine or routines individualized for each interrupt to secure the context of the processor during interrupt processing.

You can manage stack storage as long as the stack does not exceed the memory space. This stack is also used for subroutine calls; the ’C24x supports subroutine calls within the ISR. The PSHD and POPD instructions can transfer data-memory values to and from the stack.

Managing ISRs Within ISRs

The ’C24x hardware stack allows you to have ISRs within ISRs. When considering nesting ISRs like this, keep the following in mind:

- If you want the ISR to be interrupted by a maskable interrupt, the ISR must unmask the interrupt by setting the appropriate individual mask bit and the appropriate IMR bit and globally reenabling interrupts by clearing the INTM bit (CLRC INTM).
- The appropriate interrupt vector register must have been read by the GISR to obtain the *vector address offset* before globally reenabling interrupts. Otherwise, a subsequent interrupt can cause the old vector value to be overwritten.
- The hardware stack is limited to eight levels. Each time an interrupt is serviced or a subroutine is entered, the return address is pushed onto the hardware stack. This provides a way to return to the previous context after the interrupt service routine. The stack contains eight locations, allowing interrupts or subroutines to be nested up to eight levels deep. (One level of the stack is reserved for debugging, to be used for break-point/single-step operations. If debugging is not used, this extra level is available for internal use.) If your software requires more than eight levels of stack, you can use the POPD and PSHD instructions to extend the stack into data memory.
- If you do not nest ISRs, you can avoid stack overflow. The ’C24x has a feature that allows you to prevent unintentional nesting. If an interrupt occurs during the execution of a CLRC INTM instruction, the device always completes CLRC INTM as well as the following instruction before the pending interrupt is processed. This ensures that a return (RET) placed immediately after the CLRC INTM can be executed before the next interrupt is processed. The processor removes the previous return address from the stack before it adds the new return address.
- If you want an ISR to occur *within* the current ISR rather than after the current ISR, place the CLRC INTM instruction more than one instruction before the return (RET) instruction.

6.3.9 Interrupt Latency

The length of an interrupt latency—the delay between when an interrupt request is made and when it is serviced—depends on many factors. This subsection describes the factors that determine minimum latency and then describes factors that may cause additional latency. The maximum latency is a function of wait states and pipeline protection.

There are several components to interrupt latency on the ’C24x:

- Peripheral interface synchronization time
- CPU response time
- ISR branching time

Peripheral Interface Synchronization Time

Peripheral Interface synchronization time is the time it takes for the interrupt request from the peripheral to be recognized by the peripheral interface, arbitrated and converted into a request to the DSP. This takes up to one SYSCLK cycle for internal interrupts from on-chip peripherals (peripheral bus operates at the SYSCLK rate); in other words, in divide-by-two clock mode, *two* CPUCLK cycles, in divide-by-four mode, *four* CPUCLK cycles. External interrupts (NMI, INTx) have a two-SYSCLK-cycle synchronization delay.

Interrupt requests from the event manager peripheral are not arbitrated by the peripheral interface; there is a one-CPUCLK-cycle delay for the CPU to recognize the interrupt.

CPU Response Time

CPU response time is the time it takes for the CPU to recognize the enabled interrupt request, acknowledge the interrupt, clear its pipeline, and begin retrieving the first instruction from the CPU’s interrupt vector table. The minimum CPU latency is *four* CPUCLK cycles. If a higher priority maskable interrupt is requested during this minimum latency period, it is masked until the ISR for the interrupt being serviced is completed. NMIs are not maskable and are serviced before the current ISR is completed.

Latency is longer if the interrupt request occurs during multicycle operations or other operations that cannot be interrupted. If a higher priority interrupt occurs during this additional latency period, it is serviced before the original lower priority interrupt, assuming both are enabled.

Here are details about the effects of multicycle instructions:

- **Memory access using wait states.** An instruction that writes to or reads from external memory may be delayed by wait states caused by the external READY pin or the on-chip wait-state generator. These wait states may affect the instruction being executed at the time the interrupt is requested, and they may affect the interrupt itself if the interrupt vector must be fetched from external memory.
- **Repeat loop.** When repeated with RPT, instructions run parallel operations in the pipeline, and the context of these additional parallel operations cannot be saved in an interrupt service routine. To protect the context of the repeated instruction, the CPU locks out all interrupts except reset until the RPT loop completes.

Note:

Reset ($\overline{RS}$) is not delayed by multicycle instructions. An NMI can be delayed by multicycle instructions.

If one interrupt is being serviced, there are other factors that delay the servicing of a new interrupt:

- A return address (incremented program counter value) is forced onto the hardware stack every time the CPU follows another interrupt service routine or other subroutine. The ’C24x has a feature that helps you keep the hardware stack from overflowing. Interrupts cannot be processed between the CLRC INTM (enable maskable interrupts) instruction and the next instruction in a program sequence. This ensures that a return instruction that directly follows CLRC INTM is executed before an interrupt is processed. The return instruction pops the previous return address off the top of the stack before the new return address is pushed onto the stack. If the interrupt were to occur before the return, the new return address would be added to the hardware stack, even if the stack were already full.
- Interrupts are also blocked after a RET instruction until at least one instruction at the return address is executed.

ISR Branching Time

ISR branching time is the time it takes to execute all the necessary branches to get to the event-specific portion of the ISR. This length of time varies depending on how you have implemented the ISR. For the simplest situation in which only one branch to the ISR is required, the minimum branching time is four DSP cycles. See subsection 6.3.8, *Additional Tasks of ISRs*, on page 6-31 for three different methods of implementing ISRs for maskable interrupts.

6.3.10 Summary of Interrupt Operation

Once an interrupt has been passed to the CPU, the CPU operates in the following manner (see Figure 6–9):

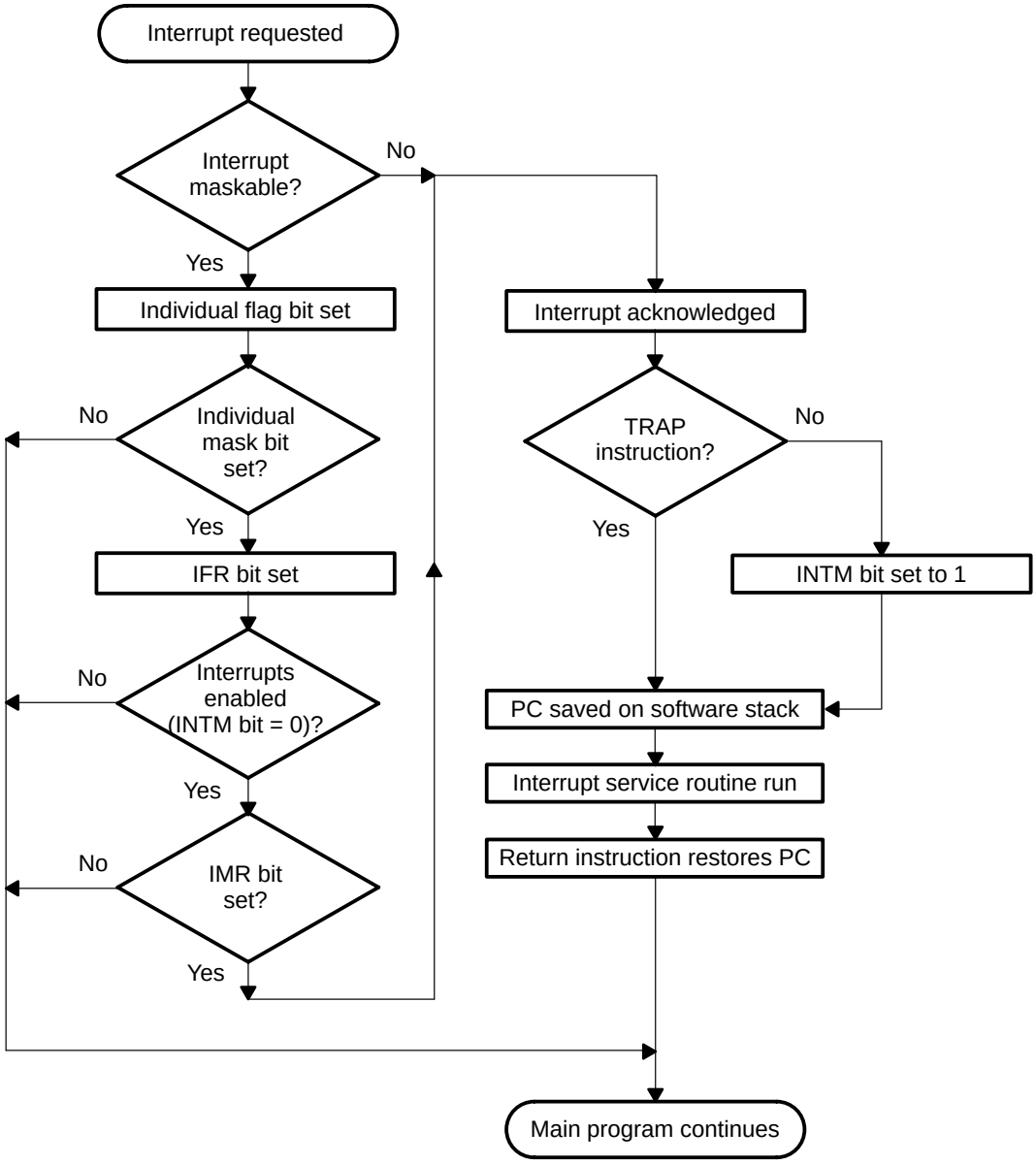
- If a maskable interrupt is requested:
  1) The flag bit in the individual control register is set. If the individual mask bit is also set, the corresponding IFR bit is set.
  2) Once the IFR bit is set, the acknowledgement conditions (INTM bit = 0 and IMR bit = 1) are tested. If the conditions are true, the CPU services the interrupt, generating an interrupt acknowledge signal; otherwise, it ignores the interrupt and continues with the current code sequence.
  3) When the interrupt has been acknowledged, the IFR bit is cleared to 0 and the INTM bit is set to 1 (to block other maskable interrupts). The flag bit in the corresponding control register is *not* cleared.
  4) The return address (incremented PC value) is saved on the stack.
  5) The CPU branches to and executes the interrupt service routine (ISR). The ISR is concluded by a return instruction, which pops the return address off the stack. The CPU continues with the interrupt code sequence.
- If a nonmaskable interrupt is requested:
  1) The CPU immediately acknowledges the interrupt, generating an interrupt acknowledge signal.

Note:

When an interrupt is requested by an INTR instruction and the corresponding IFR bit is set, the CPU does not clear the bit automatically. If an application requires that the IFR bit be cleared, the bit must be cleared by software.

  2) If the interrupt was requested by the $\overline{RS}$ pin, the NMI pin, the NMI instruction, or the INTR instruction, the the INTM bit is set to 1 to block maskable hardware interrupts. If the interrupt was requested by the TRAP instruction, the INTM bit is *not* set to 1.
  3) The return address (incremented PC value) is saved on the stack.
  4) The CPU branches to and executes the ISR. The ISR is concluded by a return instruction, which pops the return address of the stack. The CPU continues with the interrupted code sequence.

Figure 6–9. Interrupt Operation Flow Chart



6.3.11 External Interrupt Control Registers

A ’C24x device has up to six external interrupt pins with software programmable polarity and, in most cases, priority. These pins are programmed using Type A, B, and C interrupt control registers. Up to 14 additional interrupt pins can be supported using the two power module interrupt registers. In some ’C24x device configurations, the power module control registers may be used to generate interrupts in response to events occurring in certain peripherals. In all cases, refer to the device data sheet for details about the implementation of external interrupts on a particular ’C24x device.

There are three types of external interrupts: Type A, Type B, and Type C. A device will typically have two of each type, although they may not all be used. The actual number and type of external interrupt pins is dependent upon the device configuration. Figure 6–10 shows the available external interrupt control registers, including the Power Module interrupt registers.

Figure 6–10. External Interrupt Control Registers

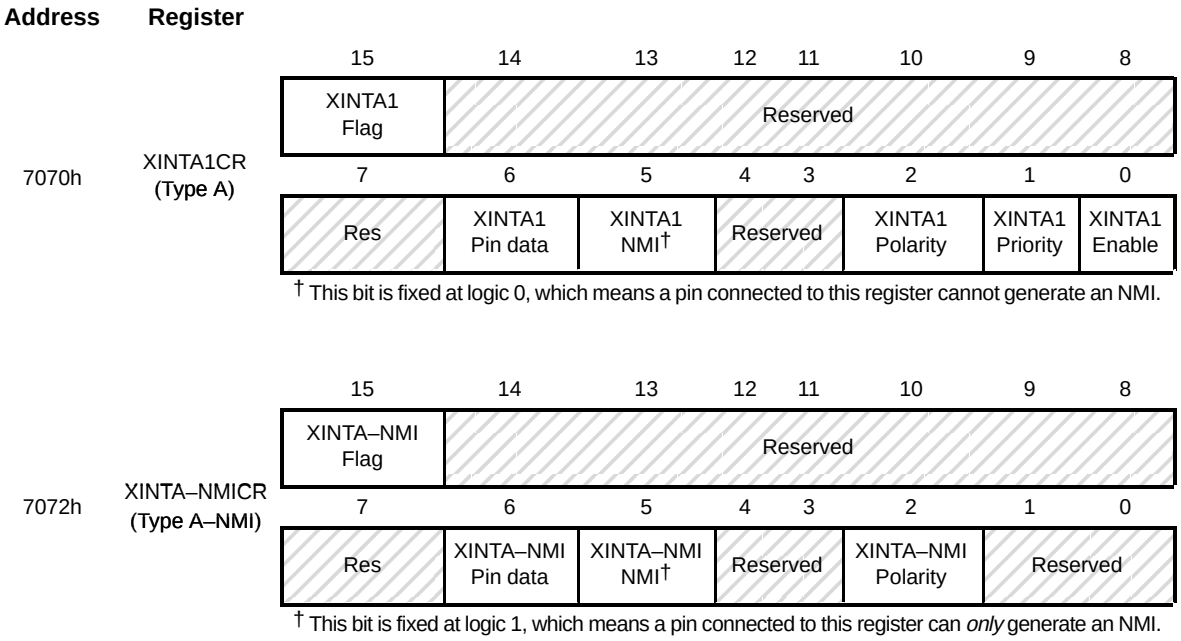


Figure 6–10. External Interrupt Control Registers (Continued)

Address **Register**

7074h XINTB1CR (Type B)

15	14	13	12	11	10	9	8
XINTB1 Flag	Reserved						
7	6	5	4	3	2	1	0
Res	XINTB1 Pin data	XINTB1 NMI	XINTB1 Data dir	XINTB1 Data out	XINTB1 Polarity	XINTB1 Priority	XINTB1 Enable

7076h XINTB2CR (Type B)

15	14	13	12	11	10	9	8
XINTB2 Flag	Reserved						
7	6	5	4	3	2	1	0
Res	XINTB2 Pin data	XINTB2 NMI	XINTB2 Data dir	XINTB2 Data out	XINTB2 Polarity	XINTB2 Priority	XINTB2 Enable

7078h XINTC1CR (Type C)

15	14	13	12	11	10	9	8
XINTC1 Flag	Reserved						
7	6	5	4	3	2	1	0
Res	XINTC1 Pin data	Res	XINTC1 Data dir	XINTC1 Data out	XINTC1 Polarity	XINTC1 Priority	XINTC1 Enable

Figure 6–10. External Interrupt Control Registers (Continued)

Address **Register**

707Ah XINTC2CR (Type C)

15	14	13	12	11	10	9	8
XINTC2 Flag	Reserved						
7	6	5	4	3	2	1	0
Res	XINTC2 Pin data	Res	XINTC2 Data dir	XINTC2 Data out	XINTC2 Polarity	XINTC2 Priority	XINTC2 Enable

707Ch PMINT1CR

15	14	13	12	11	10	9	8
PMINT1 Flag	PM INT STS 13	PM INT STS 12	PM INT STS 11	PM INT STS 10	PM INT STS 9	PM INT STS 8	PM INT STS 7
7	6	5	4	3	2	1	0
PMINT1 Enable	PM INT ENA 13	PM INT ENA 12	PM INT ENA 11	PM INT ENA 10	PM INT ENA 9	PM INT ENA 8	PM INT ENA 7

707Eh PMINT2CR

15	14	13	12	11	10	9	8
PMINT2 Flag	PM INT STS 6	PM INT STS 5	PM INT STS 4	PM INT STS 3	PM INT STS 2	PM INT STS 1	PM INT STS 0
7	6	5	4	3	2	1	0
PMINT2 Enable	PM INT ENA 6	PM INT ENA 5	PM INT ENA 4	PM INT ENA 3	PM INT ENA 2	PM INT ENA 1	PM INT ENA 0

6.3.12 Type A, Type B, and Type C Interrupt Pins

Table 6–11 summarizes the external interrupt pin types. Note the following:

- ❏ Type A interrupt pins allow for digital input only.
- ❏ One Type A pin is a maskable interrupt pin, the other is a nonmaskable interrupt pin.
- ❏ Type B pins can be programmed as maskable or nonmaskable interrupts.
- ❏ Type C inputs can only be maskable.
- ❏ Type B and Type C interrupt pins can be used as digital input or output.

- ❑ All three types can be active on rising or falling edges; the polarity is programmable.
- ❑ All three types of interrupts can be configured for low or high priority interrupt requests.
- ❑ All interrupt pins are configured to digital inputs on reset.

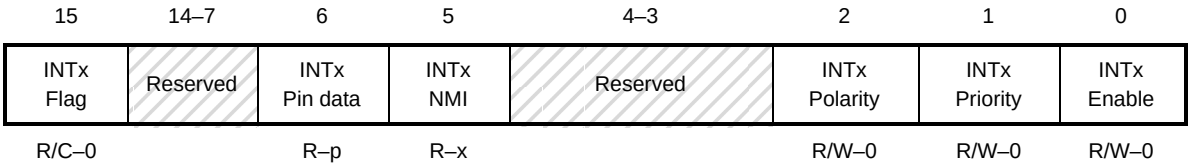
Table 6–11. External Interrupt Pin Types

Pin Type	NMI Capability?	Number Available	Digital I/O
Type A – NMI	Yes – hardwired	1	Input only
Type A	No – hardwired	1	Input only
Type B	Yes – programmable	2	I/O
Type C	No	2	I/O

Type A Interrupt Pins

Type A interrupt pins can be used as nonmaskable interrupts, normal maskable interrupts, or digital input pins. Type A control registers have the general form shown in Figure 6–11.

Figure 6–11. Type A Interrupt Control Register



Note: R = Read access, W = Write access, C = Clear-only write access, –n = Value after reset (x means value unchanged by reset), –p = Logic level of pin

Bit 15 **INTx Flag.** Interrupt x flag. This read/clear bit indicates whether the selected transition has been detected on the pin for interrupt x. This bit is set whether or not the interrupt is enabled. You can use this bit for software polling to see if the selected edge has occurred. This bit is cleared by software or a system reset. This bit need not be cleared when this pin is used as an interrupt; the interrupt occurs once for each selected edge on the interrupt pin, even though this bit is already set. Clearing this bit, however, clears a pending request from the pin.

0 = No transition detected
1 = Transition detected

Bits 14–7 **Reserved.** Reads are undefined; writes have no effect.

Bit 6 **INTx Pin data.** Interrupt pin data bit. This read-only bit reflects the current level on the interrupt pin, regardless of how the interrupt pin is configured.

0 = Pin is a logic 0
1 = Pin is a logic 1

Bit 5 **INTx NMI.** Nonmaskable interrupt enable bit. This read-only bit determines whether this pin can generate a nonmaskable interrupt. On most 'C24x devices, one Type A interrupt register has this bit hardwired to a logic 0 level; the other register has this bit hardwired to a logic 1 level.

0 = Pin is for a regular interrupt or a digital input
1 = Pin is for a nonmaskable interrupt

Bits 4–3 **Reserved.** Reads are undefined; writes have no effect.

Bit 2 **INTx Polarity.** Interrupt polarity bit. This read/write bit determines whether interrupts are generated on the rising edge or the falling edge of a signal on the pin.

0 = Interrupt generated on a falling edge (high to low transition)
1 = Interrupt generated on a rising edge (low to high transition)

Bit 1 **INTx Priority.** Interrupt priority bit. This read/write bit determines which interrupt priority is requested. This bit has no effect if the NMI bit is set. See the device data sheet for details about which interrupt level corresponds to high priority and which interrupt level corresponds to low priority.

0 = High priority
1 = Low priority

Bit 0 **INTx Enable.** Interrupt enable bit. This read/write bit enables or disables the maskable interrupt. This bit has no effect if the NMI bit is set.

0 = Disable interrupt (use pin as digital input)
1 = Enable interrupt

Type B Interrupt Pins

Type B interrupt pins can be used as nonmaskable interrupts, normal interrupts, digital output or digital input pins. The general form for a Type B control registers is shown in Figure 6–12.

Figure 6–12. Type B Interrupt Control Register

15	14–7	6	5	4	3	2	1	0
INTx Flag	Reserved	INTx Pin data	INTx NMI	INTx Data dir	INTx Data out	INTx Polarity	INTx Priority	INTx Enable
R/C–0		R–p	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0

Note: R = Read access, W = Write access, C = Clear-only write access, –n = Value after reset, –p = Logic level of pin

Bit 15 **INTx Flag.** Interrupt flag. This read/clear bit indicates if the selected transition has been detected. This bit is set whether or not the interrupt is enabled. You can use this bit for software polling to see if the selected edge has occurred. This bit can only be cleared by software or a system reset. This bit need not be cleared when this pin is used as an interrupt. The interrupt occurs once for each selected edge on the interrupt pin, even though this bit is already set. Clearing this bit, however, clears a pending request from the pin.

0 = No transition detected
1 = Transition detected

Bits 14–7 **Reserved.** Reads are undefined; writes have no effect.

Bit 6 **INTx Pin data.** Interrupt pin data bit. This read-only bit reflects the current level on the interrupt pin, regardless of how the interrupt pin is configured.

0 = Pin is a logic 0
1 = Pin is a logic 1

Bit 5 **INTx NMI.** Nonmaskable interrupt enable bit. This read/write bit determines whether this pin can generate a nonmaskable interrupt.

0 = Pin is for a regular interrupt or a digital I/O
1 = Pin is for a nonmaskable interrupt

Bit 4 **INTx Data dir.** Interrupt pin data direction bit. When the interrupt pin is not enabled for interrupts, this read/write bit determines whether the pin is a digital input or a digital output.

0 = Pin is an input
1 = Pin is an output

Bit 3 **INTx Data out.** Interrupt pin output data bit. This read/write bit determines whether the logic level on the pin is low or high when the pin is used as a digital output pin.

0 = Pin level is low (when pin used as digital output)
1 = Pin level is high (when pin used as digital output)

Bit 2 **INTx Polarity.** Interrupt polarity bit. This read/write bit determines whether interrupts are generated on the rising edge or the falling edge of a signal on the pin.

0 = Interrupt generated on a falling edge (high to low transition)
1 = Interrupt generated on a rising edge (low to high transition)

Bit 1 **INTx Priority.** Interrupt priority bit. This read/write bit determines which interrupt priority is requested. This bit has no effect if the NMI bit is set. See the device data sheet for details about which interrupt level corresponds to high priority and which interrupt level corresponds to low priority.

0 = High priority
1 = Low priority

Bit 0 **INTx Enable.** Interrupt enable bit. This read/write bit enables or disables the maskable interrupt. This bit has no effect if the NMI bit is set.

0 = Disable interrupt (use pin as digital input or output)
1 = Enable interrupt

Type C Interrupt Pins

Type C interrupt pins can be used as normal interrupts, digital output, or digital input pins. Figure 6–13 shows the general form of a Type C control register.

Figure 6–13. Type C Interrupt Control Register

15	14–7	6	5	4	3	2	1	0
INTx Flag	Reserved	INTx Pin data	Reserved	INTx Data dir	INTx Data out	INTx Polarity	INTx Priority	INTx Enable
R/C–0		R–p	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0

Note: R = Read access, W = Write access, C = Clear-only write access, –n = Value after reset, –p = Logic level of pin

Bit 15 **INTx Flag.** Interrupt flag. This read/clear bit indicates if the selected transition has been detected. This bit is set whether or not the interrupt is enabled. You can use this bit for software polling to see if the selected edge has occurred. This bit can only be cleared by software or a system reset. This bit need not be cleared when this pin is used as an interrupt. The interrupt occurs once for each selected edge on the interrupt pin, even though this bit is already set. Clearing this bit, however, clears a pending request from the pin.

0 = No transition detected
1 = Transition detected

Bits 14–7 **Reserved.** Reads are undefined; writes have no effect.

Bit 6 **INTx Pin data.** Interrupt pin data bit. This read-only bit reflects the current level on the interrupt pin, regardless of how the interrupt pin is configured.

0 = Pin is a logic 0
1 = Pin is a logic 1

Bits 5 **Reserved.** Reads are undefined; writes have no effect.

Bit 4 **INTx Data dir.** Interrupt pin data direction bit. When the interrupt pin is not enabled for interrupts, this read/write bit determines whether the pin is a digital input or a digital output.

0 = Pin is an input
1 = Pin is an output

Bit 3 **INTx Data out.** Interrupt pin output data bit. This read/write bit determines whether the logic level on the pin is low or high when the pin is used as a digital output pin.

0 = Pin level is low (when pin used as digital output)
1 = Pin level is high (when pin used as digital output)

Bit 2 **INTx Polarity.** Interrupt polarity bit. This read/write bit determines whether interrupts are generated on the rising edge or the falling edge of a signal on the pin.

0 = Interrupt generated on a falling edge (high to low transition)
1 = Interrupt generated on a rising edge (low to high transition)

Bit 1 **INTx Priority.** Interrupt priority bit. This read/write bit determines which interrupt priority is requested. See the device data sheet for details about which interrupt level corresponds to high priority and which interrupt level corresponds to low priority.

0 = High priority
1 = Low priority

Bit 0 **INTx Enable.** Interrupt enable bit. This read/write bit enables or disables the maskable interrupt.

0 = Disable interrupt (use pin as digital input or output)
1 = Enable interrupt

Summary of External Interrupt Pin Functions

Table 6–12. External Interrupt Pin Functions and Corresponding Bit Settings

	Control Register Bit Values					
Pin Used As	**NMI Bit†**	**Data Out**	**Data Dir‡**	**Polarity§**	**Priority**	**Int Enable**
Nonmaskable interrupt	1	X	X	0, 1	X	X
Interrupt high priority	0	X	X	0, 1	0	1
Interrupt low priority	0	X	X	0, 1	1	1
Digital output – 0	0	0	1	X	X	0
Digital output – 1	0	1	1	X	X	0
Digital input	0	X	0	X	X	0

† Type C interrupts do not have an NMI bit; assume a value of 0.
‡ Type A interrupts do not have a data direction bit.
§ A polarity value of 1 indicates rising edge and a value of 0 indicates falling edge.

6.3.13 Power Module Interrupts

Power module interrupts may be used to interface to signals from internal peripherals such as linear module fault condition signals (high or low active) that need to request an interrupt. These interrupts may also be used with signals coming from external pins.

Each interrupt signal has one enable bit and one interrupt flag. Each set of seven internal interrupts has a single interrupt vector. The interrupt priority level is determined at device fabrication; it is not programmable. Each power module interrupt flag can have either an active high input or an active low input. See the device data sheet for details.

The bits available in a power module interrupt register are shown in Figure 6–14 and Figure 6–15.

Figure 6–14. PM INT Flag Bits

15	14	13	12	11	10	9	8
PM INT Flag	PM INT Status 6	PM INT Status 5	PM INT Status 4	PM INT Status 3	PM INT Status 2	PM INT Status 1	PM INT Status 0
R/C–0	R/C–0	R/C–0	R/C–0	R/C–0	R/C–0	R/C–0	R/C–0

Note: R = Read access, W = Write access, C = Clear-only write access, –n = Value after reset

Figure 6–15. PM INT Enable Bits

7	6	5	4	3	2	1	0
PM INT ENA	PM INT Enable 6	PM INT Enable 5	PM INT Enable 4	PM INT Enable 3	PM INT Enable 2	PM INT Enable 1	PM INT Enable 0
R/W–0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0	R/W–0

Note: R = Read access, W = Write access, –n = Value after reset

Bit 15 **PM INT Flag.** Power module interrupt flag. This bit is set any time one of the power module interrupt sources sees the appropriate edge when the corresponding PM INT Enable x bit is set. This flag can only be cleared by writing a 0; writing a 1 has no effect.

0 = Power module interrupt event has not occurred since the flag was last cleared
1 = Power module interrupt event has occurred since the flag was last cleared

Bits 14–8 **PM INT Status x.** Power module interrupt status flags. These read-only bits reflect the status of the input source signals to the power module interrupts. If the source is in its active state, causing an interrupt, this bit contains a 1; otherwise, it contains a 0.

0 = Power module interrupt is inactive
1 = Power module interrupt is active

Bit 7 **PM INT ENA.** Power module interrupt enable bit. This bit designates whether the seven power module interrupts associated with this register are able to generate an interrupt request to the CPU. If this bit is cleared, none of the seven related interrupts in bits 14–8 can cause an interrupt request to be generated. If this bit is set, an active and enabled interrupt can be acknowledged. This bit provides a quick means to disable all power module interrupts associated with a register. The wakeup signal associated with this interrupt is also disabled when the interrupt is disabled (see Section 6.5, *Power-Down Modes*, on page 6-51).

0 = All power module interrupts are disabled
1 = All power module interrupts are enabled

Bits 6–0 **PM INT Enable x.** Power module interrupt enable bits. These bits specify whether the power module interrupt sources are enabled to set the PM INT Flag (bit 15). To allow an interrupt from a particular power module interrupt input, the corresponding PM INT Enable x bit must be set, as well as the PM INT ENA bit.

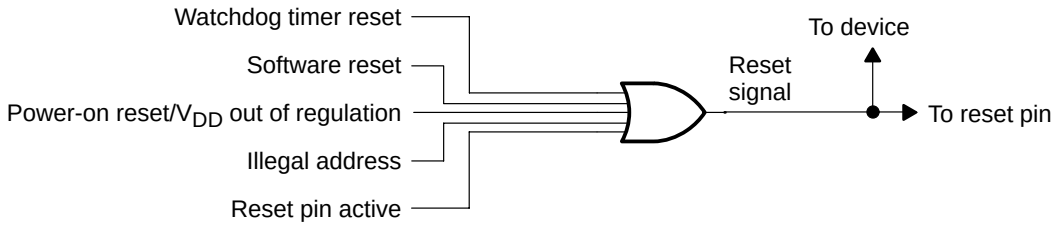
0 = Power module interrupt x is disabled
1 = Power module interrupt x is enabled

6.4 Reset Operation

The reset ($\overline{RS}$) pin generates a nonmaskable external interrupt that can be used at any time to put the ’C24x into a known state. Reset is the highest priority interrupt; no other interrupt takes precedence over a reset. Reset is typically applied after power up when the machine is in an unknown state. Because the reset signal aborts memory operations and initializes status bits, the system should be reinitialized after each reset. The NMI interrupt can be used for soft resets because it neither aborts memory operations nor initializes status bits.

Depending on the device configuration, there are up to five causes of a device reset, as shown in Figure 6–16. Four of these causes are internally generated; the other cause, the $\overline{RS}$ pin, is controlled externally.

Figure 6–16. Reset Signals



The five possible reset signals are generated as follows:

- **Watchdog timer reset.** A Watchdog timer-generated reset occurs if the Watchdog timer (if present) overflows or an improper value is written to either the Watchdog key register or the Watchdog control register. (Note that when the device is powered on, the Watchdog timer is automatically active.)
- **Software-generated reset.** This is implemented with the system control register (SCR). Clearing the RESET0 bit (bit 14) or setting the RESET1 bit (bit 15) causes a system reset.
- **Power-on reset/V_{DD} out of regulation.** This reset action is generated by either of two sources:
  - The external power-on reset pin (PORESET)
  - The low voltage detect circuitry, if present on the device. This circuitry sends a reset signal if the ’C24x device is operating with V_{DD} outside the recommended operating range.
- **Illegal address.** The system and peripheral module control register frame address map contains unimplemented address locations in the ranges labeled ‘Illegal’ in a device data sheet. Any access to an address located in the Illegal ranges generates an illegal-address reset.

- **Reset pin active.** To generate an external reset pulse on the $\overline{RS}$ pin, a low level pulse duration of as little as a few nanoseconds is usually effective; however, pulses of one SYSCLK cycle are necessary to ensure that the device recognizes the reset signal. A typical reset circuit required for a 'C24x device consists of a 10-kΩ pull-up resistor from the $\overline{RS}$ pin to V_{DD}.

Once a reset source is activated, the external $\overline{RS}$ pin is driven (active) low for a minimum of eight SYSCLK cycles. This allows the 'C24x to reset external devices connected to the $\overline{RS}$ pin. (The $\overline{RS}$ pin is an open collector I/O pin and should have a pull-up resistor attached.) Additionally, if a V_{DD} out-of-regulation condition occurs or this $\overline{RS}$ pin is held low, the reset logic holds the device in a reset state for as long as these actions are active.

When a reset signal is received, the program determines the source of the reset by reading the contents of the system status register (SYSSR). The SYSSR contains one status bit for each of the five internal sources that can cause a reset.

The occurrence of a reset condition causes the 'C24x to terminate execution and affects various registers and status bits. During a reset, RAM contents remain unchanged, and all control bits that are affected by a reset are initialized. Then processor execution begins at location 0, which normally contains a branch instruction to the system initialization routine.

When a 'C24x reset occurs, the following actions take place:

- A logic 0 is loaded into the CNF (configuration control) bit in status register ST1; this maps dual-access RAM block 0 into the data space.
- The program counter is cleared to 0.
- The INTM (interrupt mode) bit is set to 1, disabling all maskable interrupts. ($\overline{RS}$ and NMI are not maskable.) Also, the interrupt flag register (IFR) and interrupt mask register (IMR) are cleared.
- The status bits are loaded with the following values: OV = 0, XF = 1, SXM = 1, PM = 0, CNF = 0, INTM = 1, and C = 1. (The other status bits remain undefined and should be initialized after a reset.)
- The global memory allocation register (GREG) is cleared to make all memory local.
- The repeat counter (RPTC) is cleared.
- The wait states (if the device has an external memory interface) are set for the maximum duration.

- The peripheral register bits are initialized as described in the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices*.

No other CPU registers or status bits (such as the accumulator, the DP, the ARP, and the auxiliary registers) are initialized.

6.5 Power-Down Modes

A ’C24x device can have up to four power-down modes that reduce the operating power of the ’C24x device by stopping the clocks (and, thus, the activity and power consumption) of the CPU and various on-chip peripherals. While the ’C24x is in a power-down mode, all of its internal contents are maintained and operation continues unaltered when the power-down mode is terminated with an interrupt. The content of all on-chip RAM remains unchanged. However, if the power-down mode is terminated with a reset, the contents of some registers are changed (the register contents that are always changed during a reset).

There are three different clock domains that can be shut off during power down:

- **CPU clock domain.** All clocks in the CPU and memories, except the interrupt control registers
- **System clock domain.** All peripheral clocks (CPUCLK or SYSCLK), the clocks for the CPU’s interrupt control registers and the analog module clock (ACLK)
- **Watchdog clock.** The nominally 16 kHz clock used to increment the Watchdog Timer and Real Time Interrupt module (WDCLK)

Note:

The terms CPUCLK and CPU clock domain, SYSCLK and system clock domain are *not* interchangeable.

The four types of possible power-down modes on a ’C24x device have decreasing levels of power consumption and increasing delays to exit the low power mode. All of the possible power-down modes may not be implemented on a ’C24x device. A low-power mode is entered when the CPU executes an IDLE instruction. Which of the four possible low power modes is entered is determined by the PLLPM(1:0) bits in the CKCR0 register in the clock module. See the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices* for more details on this register and on the power-down modes. The power-down modes in order of decreasing power and increasing startup time are:

- **Idle1** mode stops the clocks to the CPU (CPU clock domain) but the clocks for all peripherals (system clock domain) continue to run. Exit from Idle1 occurs immediately following any interrupt or a reset.

- **Idle2** mode stops the clocks in both the CPU clock domain and the system clock domain. Exit from Idle2 occurs immediately following a wakeup interrupt or a reset. The Watchdog clock continues to run and eventually times out, causing a reset.
- **PLL power-down** mode powers down the PLL (if enabled). The Watchdog clock continues to run. Exit from this mode can be caused by a wakeup interrupt or a reset. The Watchdog clock continues to run and eventually times out, causing a reset. The device does not start full-speed operation until the PLL has powered up and reattained lock (this can be hundreds of microseconds; see the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices* for more details).
- **Oscillator power-down** mode shuts off power to the oscillator (if enabled) and is the lowest power mode available. No clocks are running on the device. A wakeup interrupt or reset causes the device to exit this low-power mode. No clock runs until the oscillator has powered back up (the power-up time is on the order of milliseconds; see the *TMS320C24x DSP Controllers Reference Set, Volume 2: Peripheral Library and Specific Devices* for more details). The device does not start full-speed operation until the PLL has powered up and reattained lock (this can be additional hundreds of microseconds).

See the relevant device data sheet for details of low-power mode power consumption for each device. Table 6–13 shows the status of the CPU and peripheral clocks during each of the four power-down modes.

Table 6–13. Power-Down Modes

Mode	CPU Clock Domain	System Clock Domain	Watchdog Clock	Oscillator
Idle1	Off	On	On	On
Idle2	Off	Off	On	On
PLL power down	Off	Off	On	On
Oscillator power down	Off	Off	Off	Off

6.5.1 Setting and Entering the Power-Down Modes

All the power-down modes are initiated by the execution of the IDLE instruction. The mode entered then depends on the value in the PLLPM field of clock module control register 0 (CKCR0). Table 6–14 shows how the two PLLPM bits of CKCR0 determine the power-down mode.

Table 6–14. Setting the Power-Down Mode with the PLLPM Bits

PLLPM bits	Power-Down Mode
00	Idle1
01	Idle2
10	PLL power down
11	Oscillator power down

6.5.2 Exiting the Power-Down Modes

In any one of the four power-down modes (Idle1, Idle2, PLL power down, and oscillator power down), the CPU clock domain is off. Therefore, software interrupts cannot be generated to take the processor out of power-down. Interrupts can only be generated at external pins or by on-chip peripherals.

Reset. Reset signals terminate power-down modes as follows:

- ❏ The reset pin ($\overline{RS}$) causes the device to exit *any* power-down mode.
- ❏ Watchdog timeout reset causes the device to exit Idle1, Idle2, or PLL power down. The Watchdog timer is not incrementing during oscillator power down because the Watchdog clock is stopped.

External interrupts. Any of the power-down modes terminate when the CPU receives any one of these interrupts at a pin:

- ❏ NMI
- ❏ XINTn (any external interrupt controlled by the external interrupt control registers, if unmasked).

Wakeup interrupts. The external interrupts XINTn and NMI are also wakeup interrupts. This means that when the clocks are shut off, there is a combinatorial logic path from these pins to restart the clocks.

The external interrupts, XINTn, are maskable interrupts and can only completely bring the processor out of the power-down mode if they are unmasked. If they are masked, they wake up the device by starting all clocks, but the device remains in the IDLE state.

Certain peripherals are also able to generate wakeup interrupts when the clocks in the system clock domain are shut off. For example, some communication ports may be able to generate a wakeup interrupt in response to receiving a character.

Oscillator power down mode is only terminated by an external interrupt (NMI or XINTn). In this mode, the oscillator is off (no clocks on the device are active); thus, no interrupts can be generated by on-chip peripherals, and the Watchdog timer cannot generate a time-out signal.

Peripheral interrupts. The Idle1, Idle2, and PLL power-down modes can be terminated by various peripheral interrupts under the right conditions. In Idle1 mode, all the clocks for the peripheral devices are still running; therefore, any unmasked peripheral interrupt terminates Idle1 mode. In the Idle2 and PLL power-down modes, the clocks in the system clock domain are off. As a result, only unmasked peripheral interrupts not timed by the system clock can bring the processor out of the Idle2 and PLL power-down modes.

In oscillator power-down mode, the peripheral clocks and the Watchdog timer clocks are off. Only unmasked peripheral interrupts not timed by either of those clocks can terminate oscillator power-down mode.

Table 6–15 summarizes the state of processors in each power-down mode and lists the interrupts that do and do not terminate each mode.

Table 6–15. Power-Down Modes and Their Termination

Mode	CPU Clock Domain	System Clock Domain	Watchdog Clock	PLL	Oscillator	Terminated By	Not Terminated By
Idle1	Off	On	On	On	On	Reset ($\overline{RS}$) Watchdog reset NMI (at pin) XINT's (at pin, unmasked) Any peripheral interrupt (unmasked)	Masked interrupts
Idle2	Off	Off	On	On	On	Reset ($\overline{RS}$) Watchdog reset NMI (at pin) XINT's (at pin, unmasked) Peripheral wakeup interrupts.	Masked interrupts Peripheral interrupts dependent on the system clock.
PLL power down (PPD)	Off	Off	On	Off	On	Reset ($\overline{RS}$) Watchdog reset NMI (at pin) XINT's (at pin, unmasked) Peripheral wakeup interrupts	Masked interrupts Peripheral interrupts dependent on the system clock
Oscillator power down (OPD)	Off	Off	Off	Off	Off	Reset ($\overline{RS}$) NMI (at pin) XINT's (at pin, unmasked) Peripheral wakeup interrupts	Masked interrupts Peripheral interrupts

After Exiting Power-Down

There are two items to consider when deciding how to wake the processor:

- If you use reset or NMI, the CPU immediately executes the corresponding interrupt service routine.
- If you use a maskable hardware interrupt, the next action depends on the interrupt mode (INTM) bit of status register ST0:
  - **INTM = 0:** The interrupt is enabled, and the CPU executes the corresponding interrupt service routine.
  - **INTM = 1:** The interrupt is disabled, and the CPU continues with the instruction after IDLE.

If you do not want the CPU to take an interrupt service routine before continuing with the interrupted program sequence:

- Do not use reset or NMI to bring the processor out of power-down.
- Make sure your program sets INTM to 1 (SETC INTM) before IDLE is executed.

If you want the CPU to take the interrupt service routine before continuing:

- Make sure your program clears INTM to 0 (CLRC INTM) before IDLE is executed.
- Make sure you enable all relevant interrupt sources, both locally (in the peripheral mask/enable registers) and globally (in the CPU's interrupt mask register).

6.5.3 Summary of Power-Down Mode Operation

When the IDLE instruction is executed:

1) The program counter is incremented once, so that when the power-down mode is exited, the next instruction is the one that follows the IDLE instruction, except in the case of reset.

2) The ’C24x enters the power-down mode selected by the PLLPM bits and remains in that low-power state until it receives a proper hardware interrupt (as described in subsection 6.5.2, *Exiting the Power-Down Modes*, on page 6-53).

3) Upon receipt of the proper interrupt, the ’C24x exits the power-down mode.

4) If you use NMI to wake the processor, the CPU executes the corresponding interrupt service routine before continuing with the interrupted program sequence.

   If you use a maskable interrupt, the next action depends on the value of the INTM bit:

   - **INTM = 0:** Maskable interrupts are enabled; the CPU first executes the interrupt service routine of the interrupt that brought it out of power-down. Then it continues with the instruction after the IDLE instruction.
   - **INTM = 1:** Maskable interrupts are disabled; the CPU continues execution at the instruction after IDLE.

Table 6–16 summarizes the four power-down modes, including the approximate power level for each. (For more accurate power values, see the data sheet for your particular ’C24x device.) In addition, the table shows the status of the ’C24x when not in a power-down mode. See Table 6–15 for the list of interrupts that do and do not terminate each mode.

Table 6–16. Power-Down Modes/Run Mode Summary

Mode	PLLPM	+	IDLE	Power Level	CPU Clock Domain	System Clock Domain	Watchdog Clock	PLL	Oscillator
Run	XX	+	No IDLE	> 40 mA	On	On	On	On	On
Idle 1	00	+	IDLE	~15 mA	Off	On	On	On	On
Idle 2	01	+	IDLE	~4 mA	Off	Off	On	On	On
PPD	10	+	IDLE	~1 mA	Off	Off	On	Off	On
OPD	11	+	IDLE	< 30 μA	Off	Off	Off	Off	Off

Chapter 7

Addressing Modes

This chapter explains the three basic memory addressing modes used by the ’C24x instruction set. The three modes are:

- Immediate addressing mode
- Direct addressing mode
- Indirect addressing mode

In immediate addressing, a constant to be manipulated by the instruction is supplied directly as an operand of that instruction. Two types of immediate addressing are available—short and long. In short-immediate addressing, an 8-, 9-, or 13-bit operand is included in the instruction word. Long-immediate addressing uses a 16-bit operand.

When you need to access data memory, you can use direct or indirect addressing. Direct addressing concatenates seven bits of the instruction word with the nine bits of the data-memory page pointer (DP) to form the 16-bit data memory address. Indirect addressing accesses data memory through one of eight 16-bit auxiliary registers.

Topic **Page**

7.1 Immediate Addressing Mode 7-2

7.2 Direct Addressing Mode 7-4

7.3 Indirect Addressing Mode 7-9

7.1 Immediate Addressing Mode

In immediate addressing, the instruction word contains a constant to be manipulated by the instruction. The ’C24x supports two types of immediate addressing:

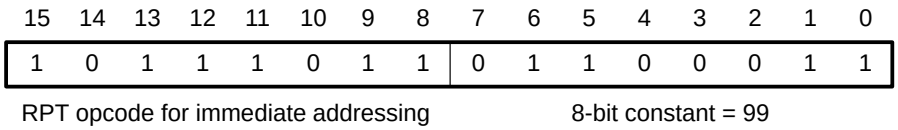
- **Short-immediate addressing.** Instructions that use short-immediate addressing take an 8-bit, 9-bit, or 13-bit constant as an operand. Short-immediate instructions require a single instruction word, with the constant embedded in that word.
- **Long-immediate addressing.** Instructions that use long-immediate addressing take a 16-bit constant as an operand and require two instruction words. The constant is sent as the second instruction word. This 16-bit value can be used as an absolute constant or as a 2s-complement value.

In Example 7–1, the immediate operand is contained as a part of the RPT instruction word. For this RPT instruction, the instruction register will be loaded with the value shown in Figure 7–1. Immediate operands are preceded by the symbol #.

Example 7–1. RPT Instruction Using Short-Immediate Addressing

```
RPT #99      ;Execute the instruction that follows RPT
             ;100 times.
```

Figure 7–1. Instruction Register Contents for Example 7–1

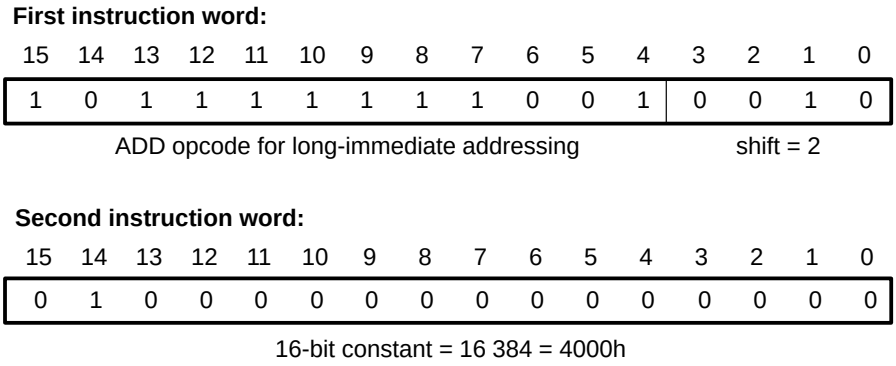


In Example 7–2, the immediate operand is contained in the second instruction word. The instruction register receives, consecutively, the two 16-bit values shown in Figure 7–2.

Example 7–2. ADD Instruction Using Long-Immediate Addressing

```
ADD   #16384,2  ;Shift the value 16384 left by two bits
                ;and add the result to the accumulator.
```

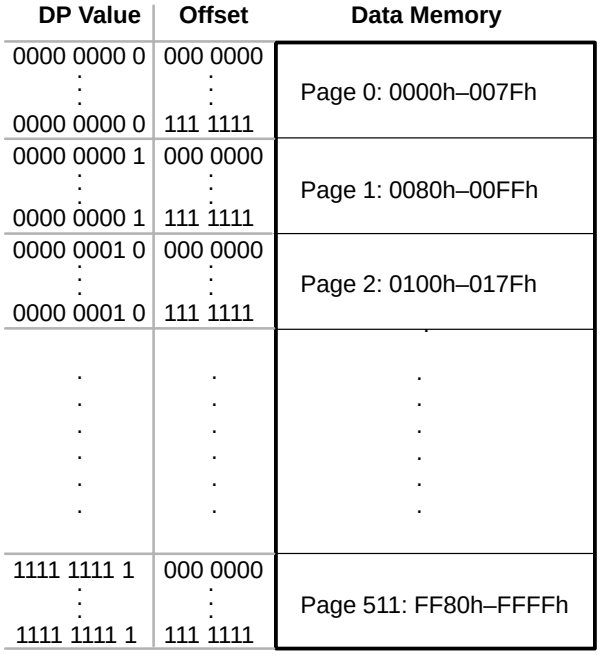
Figure 7–2. Two Words Loaded Consecutively to the Instruction Register in Example 7–2



7.2 Direct Addressing Mode

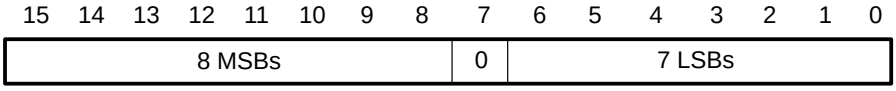
In the direct addressing mode, data memory is addressed in blocks of 128 words called data pages. The entire 64K of data memory consists of 512 data pages labeled 0 through 511, as shown in Figure 7–3. The current data page is determined by the value in the 9-bit data page pointer (DP) in status register ST0. For example, if the DP value is 0 0000 0000_2, the current data page is 0. If the DP value is 0 0000 0010_2, the current data page is 2.

Figure 7–3. Pages of Data Memory



In addition to the data page, the processor must know the particular word being referenced on that page. This is determined by a 7-bit offset (see Figure 7–3). The offset is supplied by the seven least significant bits (LSBs) of the instruction register, which holds the opcode for the next instruction to be executed. In direct addressing mode, the content of the instruction register has the format shown in Figure 7–4.

Figure 7–4. Instruction Register (IR) Contents in Direct Addressing Mode



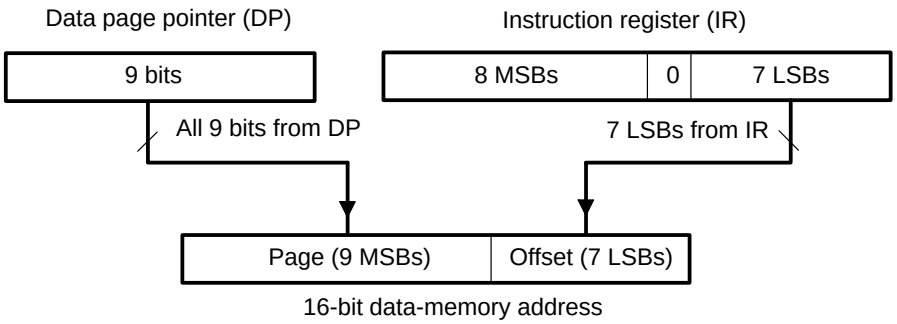
8 MSBs Bits 15 through 8 indicate the instruction type (for example, ADD) and also contain any information regarding a shift of the data value to be accessed by the instruction.

0 **Direct/indirect indicator.** Bit 7 contains a 0 to define the addressing mode as direct.

7 LSBs Bits 6 through 0 indicate the offset for the data-memory address referenced by the instruction.

To form a complete 16-bit address, the processor concatenates the DP value and the seven LSBs of the instruction register, as shown in Figure 7–5. The DP supplies the nine most significant bits (MSBs) of the address (the page number), and the seven LSBs of the instruction register supply the seven LSBs of the address (the offset). For example, to access data address 003Fh, you specify data page 0 (DP = 0000 0000 0) and an offset of 011 1111. Concatenating the DP and the offset produces the 16-bit address 0000 0000 0011 1111, which is 003Fh or decimal 63.

Figure 7–5. Generation of Data Addresses in Direct Addressing Mode



Initialize the DP in All Programs

It is critical that all programs initialize the DP. The DP is not initialized by reset and is undefined after power up. The ’C24x development tools use default values for many parameters, including the DP. However, programs that do not explicitly initialize the DP can execute improperly, depending on whether they are executed on a ’C24x device or with a development tool.

7.2.1 Using Direct Addressing Mode

When you use direct addressing mode, the processor uses the DP to find the data page and uses the seven LSBs of the instruction register to find a particular address on that page. Always do the following:

1) **Set the data page.** Load the appropriate value (from 0 to 511) into the DP. The DP register can be loaded by the LDP instruction or by any instruction that can load a value to ST0. The LDP instruction loads the DP directly without affecting the other bits of ST0, and it clearly indicates the value loaded into the DP. For example, to set the current data page to 32 (addresses 1000h–107Fh), you can use:

```
LDP #32   ;Initialize data page pointer
```

2) **Specify the offset.** Supply the 7-bit offset as an operand of the instruction. For example, if you want the ADD instruction to use the value at the second address of the current data page, you would write:

```
ADD 1h ;Add to accumulator the value in the current
   ;data page, offset of 1.
```

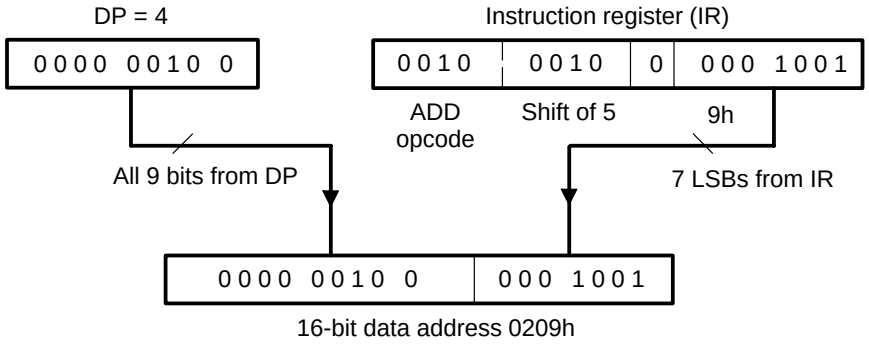
You do not have to set the data page prior to every instruction that uses direct addressing. If all the instructions in a block of code access the same data page, you can simply load the DP at the front of the block. However, if various data pages are being accessed throughout the block of code, be sure the DP is changed whenever a new data page should be accessed.

7.2.2 Examples of Direct Addressing

In Example 7–3, the first instruction loads the DP with $0\ 0000\ 0100_2$ to set the current data page to 4. The ADD instruction then references a data memory address that is generated as shown following the program code. Before the ADD instruction is executed, the opcode is loaded into the instruction register. Together, the DP and the seven LSBs of the instruction register form the complete 16-bit address, $0000\ 0010\ 0000\ 1001_2$ (0209h).

Example 7–3. Using Direct Addressing with ADD (Shift of 0 to 15)

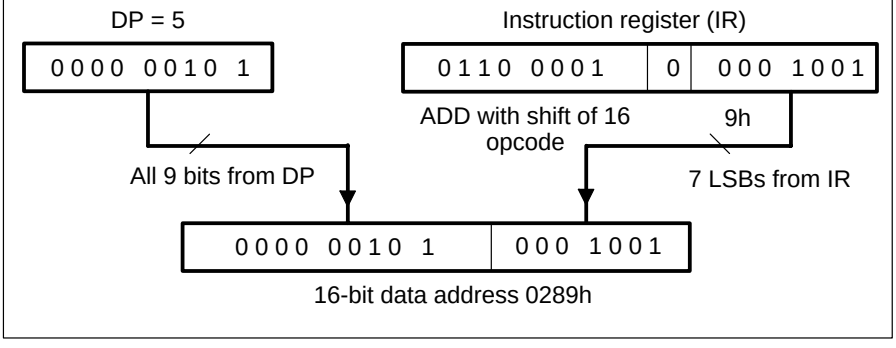
```
LDP #4    ;Set data page to 4 (addresses 0200h–027Fh).
ADD 9h,5  ;The contents of data address 0209h are
          ;left shifted 5 bits and added to the
          ;contents of the accumulator.
```



In Example 7–4, the ADD instruction references a data memory address that is generated as shown following the program code. For any instruction that performs a shift of 16, the shift value is not embedded directly in the instruction word; instead, all eight MSBs contain an opcode that not only indicates the instruction type, but also a shift of 16. The eight MSBs of the instruction word indicate an ADD with a shift of 16.

Example 7–4. Using Direct Addressing with ADD (Shift of 16)

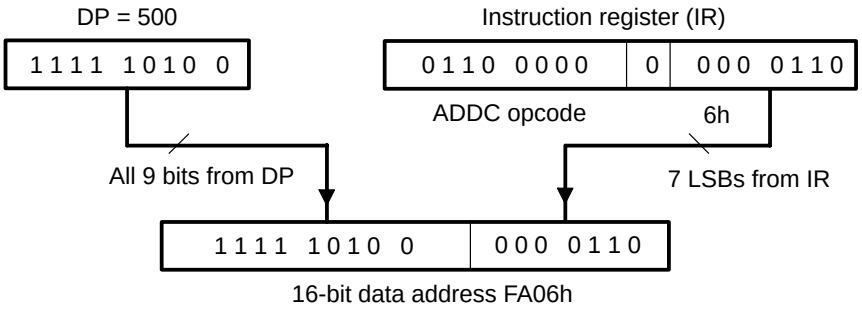
```
LDP #5     ;Set data page to 5 (addresses 0280h–02FFh).
ADD 9h,16  ;The contents of data address 0289h are
           ;left shifted 16 bits and added to the
           ;contents of the accumulator.
```



In Example 7–5, the ADDC instruction references a data memory address that is generated as shown following the program code. Note that if an instruction does not perform shifts, like the ADDC instruction does not, all eight MSBs of the instruction contain the opcode for the instruction type.

Example 7–5. Using Direct Addressing with ADDC

```
LDP #500  ;Set data page to 500 (addresses FA00h–FA7Fh).
ADDC 6h   ;The contents of data address FA06h
          ;and the value of the carry bit (C) are
          ;added to the contents of the accumulator.
```



7.3 Indirect Addressing Mode

Eight auxiliary registers (AR0–AR7) provide flexible and powerful indirect addressing. Any location in the 64K data memory space can be accessed using a 16-bit address contained in an auxiliary register.

7.3.1 Current Auxiliary Register

To select a specific auxiliary register, load the 3-bit auxiliary register pointer (ARP) of status register ST0 with a value from 0 to 7. The ARP can be loaded as a primary operation by the MAR instruction or by the LST instruction. The ARP can be loaded as a secondary operation by any instruction that supports indirect addressing.

The register pointed to by the ARP is referred to as the *current auxiliary register* or *current AR*. During the processing of an instruction, the content of the current auxiliary register is used as the address at which the data-memory access takes place. The ARAU passes this address to the data-read address bus (DRAB) if the instruction requires a read from data memory, or it passes the address to the data-write address bus (DWAB) if the instruction requires a write to data memory. After the instruction uses the data value, the contents of the current auxiliary register can be incremented or decremented by the ARAU, which implements unsigned 16-bit arithmetic.

Normally, the ARAU performs its arithmetic operations in the decode phase of the pipeline (when the instruction specifying the operations is being decoded). This allows the address to be generated before the decode phase of the next instruction. There is an exception to this rule: during processing of the NORM instruction, the auxiliary register and/or ARP modification is done during the execute phase of the pipeline. For information on pipeline operation, see Section 5.2 on page 5-7.

7.3.2 Indirect Addressing Options

The ’C24x provides four types of indirect addressing options:

- **No increment or decrement.** The instruction uses the content of the current auxiliary register as the data memory address but neither increments nor decrements the content of the current auxiliary register.
- **Increment or decrement by 1.** The instruction uses the content of the current auxiliary register as the data memory address and then increments or decrements the content of the current auxiliary register by one.
- **Increment or decrement by an index amount.** The value in AR0 is the index amount. The instruction uses the content of the current auxiliary register as the data memory address and then increments or decrements the content of the current auxiliary register by the index amount.

- **Increment or decrement by an index amount using reverse carry.** The value in AR0 is the index amount. After the instruction uses the content of the current auxiliary register as the data-memory address, that content is incremented or decremented by the index amount. The addition or subtraction, in this case, is done with the carry propagation reversed for fast Fourier transforms (FFTs).

These four option types provide the seven indirect addressing options listed in Table 7–1. The table also shows the instruction operand that corresponds to each indirect addressing option and gives an example of how each option is used.

Table 7–1. Indirect Addressing Operands

Operand	Option	Example
*	No increment or decrement	**LT *** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR.
*+	Increment by 1	**LT *+** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR and then adds 1 to the content of the current AR.
*–	Decrement by 1	**LT *–** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR and then subtracts 1 from the content of the current AR.
*0+	Increment by index amount	**LT *0+** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR and then adds the content of AR0 to the content of the current AR.
*0–	Decrement by index amount	**LT *0–** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR and then subtracts the content of AR0 from the content of the current AR.
*BR0+	Increment by index amount, adding with reverse carry	**LT *BR0+** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR and then adds the content of AR0 to the content of the current AR, adding with reverse carry propagation.
*BR0–	Decrement by index amount, subtracting with reverse carry	**LT *BR0–** loads the temporary register (TREG) with the content of the data memory address referenced by the current AR and then subtracts the content of AR0 from the content of the current AR, subtracting with bit reverse carry propagation.

All increments or decrements are performed by the auxiliary register arithmetic unit (ARAU) in the same cycle during which the instruction is being decoded in the pipeline.

The bit-reversed indexed addressing allows efficient I/O operations by resequencing the data points in a radix-2 FFT program. The direction of carry propagation in the ARAU is reversed when the address is selected, and AR0 is added to or subtracted from the current auxiliary register. A typical use of this addressing mode requires that AR0 first be set to a value corresponding to half of the array's size, and that the current AR value be set to the base address of the data (the first data point).

7.3.3 Next Auxiliary Register

In addition to updating the current auxiliary register, a number of instructions can also specify the *next auxiliary register* or *next AR*. This register will be the current auxiliary register when the instruction execution is complete. The instructions that allow you to specify the next auxiliary register load the ARP with a new value. When the ARP is loaded with that value, the previous ARP value is loaded into the auxiliary register pointer buffer (ARB). Example 7–6 illustrates the selection of a next auxiliary register, as well as other indirect addressing features discussed so far.

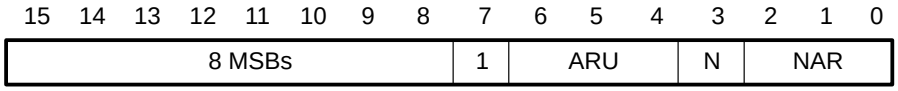
Example 7–6. Selecting a New Current Auxiliary Register

```
MAR*,AR1     ;Load the ARP with 1 to make AR1 the
             ;current auxiliary register.
LT  *+,AR2   ;AR2 is the next auxiliary register.
             ;Load the TREG with the content of the
             ;address referenced by AR1, add one to
             ;the content of AR1, then make AR2 the
             ;current auxiliary register.
MPY*         ;Multiply TREG by content of address
             ;referenced by AR2.
```

7.3.4 Indirect Addressing Opcode Format

Figure 7–6 shows the format of the instruction word loaded into the instruction register when you use indirect addressing. The opcode fields are described following Figure 7–6.

Figure 7–6. Instruction Register Content in Indirect Addressing



8 MSBs Bits 15 through 8 indicate the instruction type (for example, LT) and also contain any information regarding data shifts.

1 **Direct/indirect indicator.** Bit 7 contains a 1 to define the addressing mode as indirect.

ARU **Auxiliary register update code.** Bits 6 through 4 determine whether and how the current auxiliary register is incremented or decremented. See Table 7–2.

N **Next auxiliary register indicator.** Bit 3 specifies whether the instruction changes the ARP value.

- N = 0 The content of the ARP remains unchanged.
- N = 1 The content of NAR is loaded into the ARP, and the old ARP value is loaded into the auxiliary register buffer (ARB) of status register ST1.

NAR **Next auxiliary register value.** Bits 2 through 0 contain the value of the next auxiliary register. NAR is loaded into the ARP if N = 1.

Table 7–2. Effects of the ARU Code on the Current Auxiliary Register

ARU Code			
6	**5**	**4**	**Arithmetic Operation Performed on Current AR**
0	0	0	No operation on current AR
0	0	1	Current AR – 1 → current AR
0	1	0	Current AR + 1 → current AR
0	1	1	Reserved
1	0	0	Current AR – AR0 → current AR [reverse carry propagation]
1	0	1	Current AR – AR0 → current AR
1	1	0	Current AR + AR0 → current AR
1	1	1	Current AR + AR0 → current AR [reverse carry propagation]

Table 7–3 shows the opcode field bits and the notation used for indirect addressing. It also shows the corresponding operations performed on the current auxiliary register and the ARP.

Table 7–3. Field Bits and Notation for Indirect Addressing

Instruction Opcode Bits												
15	**–**	**8**	**7**	**6**	**5**	**4**	**3**	**2**	**1**	**0**	**Operand(s)**	**Operation**
←	8 MSBs	→	1	0	0	0	0	←NAR→			*****	No manipulation of current AR
←	8 MSBs	→	1	0	0	0	1	←NAR→			*,AR*n*	NAR → ARP
←	8 MSBs	→	1	0	0	1	0	←NAR→			***–**	Current AR – 1 → current AR
←	8 MSBs	→	1	0	0	1	1	←NAR→			*–,AR*n*	Current AR – 1 → current AR NAR → ARP
←	8 MSBs	→	1	0	1	0	0	←NAR→			*+	Current AR + 1 → current AR
←	8 MSBs	→	1	0	1	0	1	←NAR→			*+,AR*n*	Current AR + 1 → current AR NAR → ARP
←	8 MSBs	→	1	1	0	0	0	←NAR→			*BR0–	Current AR – *rc*AR0 → current AR †
←	8 MSBs	→	1	1	0	0	1	←NAR→			*BR0–,AR*n*	Current AR – *rc*AR0 → current AR NAR → ARP †
←	8 MSBs	→	1	1	0	1	0	←NAR→			*0–	Current AR – AR0 → current AR
←	8 MSBs	→	1	1	0	1	1	←NAR→			*0–,AR*n*	Current AR – AR0 → current AR NAR → ARP
←	8 MSBs	→	1	1	1	0	0	←NAR→			*0+	Current AR + AR0 → current AR
←	8 MSBs	→	1	1	1	0	1	←NAR→			*0+,AR*n*	Current AR + AR0 → current AR NAR → ARP
←	8 MSBs	→	1	1	1	1	0	←NAR→			*BR0+	Current AR + *rc*AR0 → current AR †
←	8 MSBs	→	1	1	1	1	1	←NAR→			*BR0+,AR*n*	Current AR + *rc*AR0 → current AR NAR → ARP †

† Bit-reversed addressing mode

Legend:

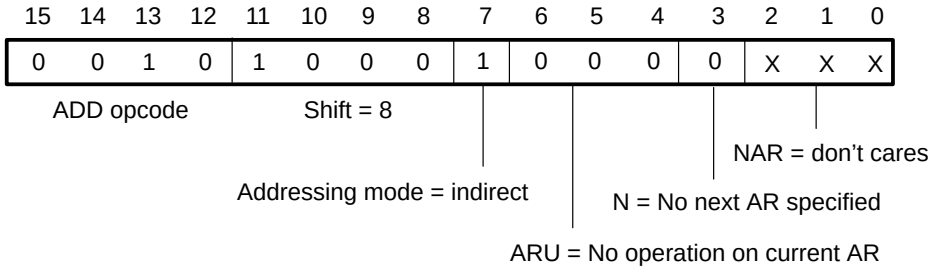
rc	Reverse carry propagation
NAR	Next AR
n	0, 1, 2, ..., or 7
8 MSBs	Eight bits determined by instruction type and (sometimes) shift information
→	Is loaded into

7.3.5 Examples of Indirect Addressing

In Example 7–7, when the ADD instruction is fetched from program memory, the instruction register is loaded with the value shown.

Example 7–7. Indirect Addressing—No Increment or Decrement

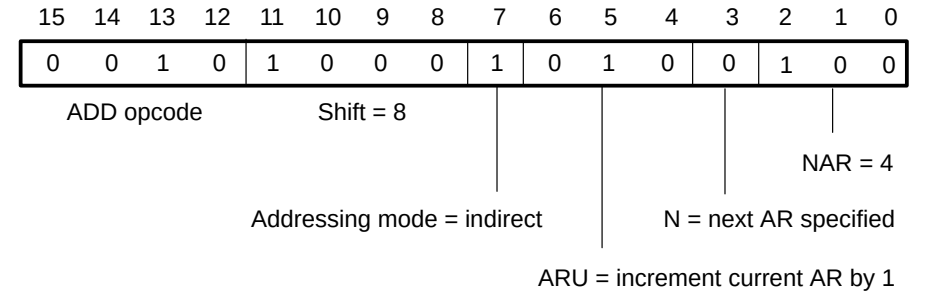
```
ADD *,8      ;Add to the accumulator the content of the
             ;data-memory address referenced by the
             ;current auxiliary register. The data
             ;is left shifted 8 bits before being added.
```



In Example 7–8, when the ADD instruction is fetched from program memory, the instruction register is loaded with the value shown.

Example 7–8. Indirect Addressing—Increment by 1

```
ADD *+,8,AR4    ;Operates as in Example 7–7, but
                ;in addition, the current auxiliary
                ;register is incremented by one, and
                ;AR4 is chosen as the next auxiliary
                ;register.
```



Example 7–9. Indirect Addressing—Decrement by 1

```
ADD *–,8     ;Operates as in Example 7–7, but in
             ;addition, the current auxiliary register
             ;is decremented by one.
```

Example 7–10. Indirect Addressing—Increment by Index Amount

```
ADD *0+,8   ;Operates as in Example 7–7, but in
            ;addition, the content of register AR0
            ;is added to the current auxiliary
            ;register.
```

Example 7–11. Indirect Addressing—Decrement by Index Amount

```
ADD *0–,8    ;Operates as in Example 7–7, but in
             ;addition, the content of register AR0
             ;is subtracted from the current auxiliary
             ;register.
```

Example 7–12. Indirect Addressing—Increment by Index Amount With Reverse Carry Propagation

```
ADD *BR0+,8 ;Operates as in Example 7–10, except that
            ;the content of register AR0 is added to
            ;the current auxiliary register with
            ;reverse carry propagation.
```

Example 7–13. Indirect Addressing—Decrement by Index Amount With Reverse Carry Propagation

```
ADD *BR0–,8 ;Operates as in Example 7–11, except that
            ;the content of register AR0 is subtracted
            ;from the current auxiliary register with
            ;reverse carry propagation.
```

7.3.6 Modifying Auxiliary Register Content

The LAR, ADRK, SBRK, and MAR instructions are specialized instructions for changing the content of an auxiliary register (AR):

- The LAR instruction loads an AR.
- The ADRK instruction adds an immediate value to an AR; SBRK subtracts an immediate value.
- The MAR instruction can increment or decrement an AR value by 1 or by an index amount.

However, you are not limited to these four instructions. Auxiliary registers can be modified by any instruction that supports indirect addressing operands. (Indirect addressing can be used with all instructions except those that have immediate operands or no operands.)

Chapter 8

Assembly Language Instructions

Note:

The instruction set for the TMS320C24x is identical to that of the TMS320C2xx. All references to ’C2xx devices in this chapter also apply to ’C24x devices.

This chapter describes the ’C24x assembly language instructions. This instruction set supports numerically intensive signal-processing operations as well as general-purpose applications, such as multiprocessing and high-speed control. The ’C2xx instruction set is compatible with the ’C2x instruction set; code written for the ’C2x can be reassembled to run on the ’C2xx. The ’C5x instruction set is a superset of that of the ’C2xx; thus, code written for the ’C2xx can be upgraded to run on a ’C5x.

Topic **Page**

8.1 Instruction Set Summary 8-2

8.2 How To Use the Instruction Descriptions 8-12

8.3 Instruction Descriptions 8-19

8.1 Instruction Set Summary

This section provides six tables (Table 8–1 to Table 8–6) that summarize the instruction set according to the following functional headings:

- Accumulator, arithmetic, and logic instructions (see Table 8–1 on page 8-5)
- Auxiliary register and data page pointer instructions (see Table 8–2 on page 8-7)
- TREG, PREG, and multiply instructions (see Table 8–3 on page 8-8)
- Branch instructions (see Table 8–4 on page 8-9)
- Control instructions (see Table 8–5 on page 8-10)
- I/O and memory operations (see Table 8–6 on page 8-11)

Within each table, the instructions are arranged alphabetically. The number of words that an instruction occupies in program memory is specified in column three of each table; the number of cycles that an instruction requires to execute is in column four. All instructions are assumed to be executed from internal program memory (RAM) and internal data dual-access memory. The cycle timings are for single-instruction execution, not for repeat mode. Additional information about each instruction is presented in the individual instruction descriptions in Section 8.2 on page 8-12.

For your reference, here are the definitions of the symbols used in the six summary tables:

ACC The accumulator

AR The auxiliary register

ARX A 3-bit value used in the LAR and SAR instructions to designate which auxiliary register will be loaded (LAR) or have its contents stored (SAR)

BITX A 4-bit value (called the bit code) that determines which bit of a designated data memory value will be tested by the BIT instruction

CM A 2-bit value. The CMPR instruction performs a comparison specified by the value of CM:

If CM = 00, test whether current AR = AR0
If CM = 01, test whether current AR < AR0
If CM = 10, test whether current AR > AR0
If CM = 11, test whether current AR ≠ AR0

IAAA AAAA (One I followed by seven As) The I at the left represents a bit that reflects whether direct addressing (I = 0) or indirect addressing (I = 1) is being used. When direct addressing is used, the seven As are the seven least significant bits (LSBs) of a data memory address. For indirect addressing, the seven As are bits that control auxiliary register manipulation (see Section 7.3, *Indirect Addressing Mode*, on page 7-9).

IIII IIII (Eight Is) An 8-bit constant used in short immediate addressing

I IIII IIII (Nine Is) A 9-bit constant used in short immediate addressing for the LDP instruction

I IIII IIII IIII (Thirteen Is) A 13-bit constant used in short immediate addressing for the MPY instruction

I NTR# A 5-bit value representing a number from 0 to 31. The INTR instruction uses this number to change program control to one of the 32 interrupt vector addresses.

PM A 2-bit value copied into the PM bits of status register ST1 by the SPM instruction

SHF A 3-bit left-shift value

SHFT A 4-bit left-shift value

TP A 2-bit value used by the conditional execution instructions to represent four conditions:

$\overline{\text{BIO}}$ pin low	TP = 00
TC bit =1	TP = 01
TC bit = 0	TP = 10
No condition	TP = 11

ZLVC ZLVC Two 4-bit fields — each representing the following conditions:

ACC = 0	Z
ACC < 0	L
Overflow	V
Carry	C

A conditional instruction contains two of these 4-bit fields. The 4-LSB field of the instruction is a mask field. A 1 in the corresponding mask bit indicates that condition is being tested. For example, to test for ACC ≥ 0, the Z and L fields are set, and the V and C fields are not set. The Z field is set to test the condition ACC = 0, and the L field is reset to test the condition ACC ≥ 0.The second 4-bit field (bits 4 – 7) indicates the state of the conditions to test. The conditions possible with these eight bits are shown in the descriptions for the BCND, CC, and RETC instructions.

+ 1 word The second word of a 2-word opcode. This second word contains a 16-bit constant. Depending on the instruction, this constant is a long immediate value, a program memory address, or an address for an I/O port or an I/O-mapped register.

Table 8–1. Accumulator, Arithmetic, and Logic Instructions

Mnemonic	Description	Words	Cycles	Opcode
ABS	Absolute value of ACC	1	1	1011 1110 0000 0000
ADD	Add to ACC with shift of 0 to 15, direct or indirect	1	1	0010 SHFT IAAA AAAA
	Add to ACC with shift 0 to 15, long immediate	2	2	1011 1111 1001 SHFT + 1 word
	Add to ACC with shift of 16, direct or indirect	1	1	0110 0001 IAAA AAAA
	Add to ACC, short immediate	1	1	1011 1000 IIII IIII
ADDC	Add to ACC with carry, direct or indirect	1	1	0110 0000 IAAA AAAA
ADDS	Add to low ACC with sign-extension suppressed, direct or indirect	1	1	0110 0010 IAAA AAAA
ADDT	Add to ACC with shift (0 to 15) specified by TREG, direct or indirect	1	1	0110 0011 IAAA AAAA
AND	AND ACC with data value, direct or indirect	1	1	0110 1110 IAAA AAAA
	AND with ACC with shift of 0 to 15, long immediate	2	2	1011 1111 1011 SHFT + 1 word
	AND with ACC with shift of 16, long immediate	2	2	1011 1110 1000 0001 + 1 word
CMPL	Complement ACC	1	1	1011 1110 0000 0001
LACC	Load ACC with shift of 0 to 15, direct or indirect	1	1	0001 SHFT IAAA AAAA
	Load ACC with shift of 0 to 15, long immediate	2	2	1011 1111 1000 SHFT + 1 word
	Load ACC with shift of 16, direct or indirect	1	1	0110 1010 IAAA AAAA
LACL	Load low word of ACC, direct or indirect	1	1	0110 1001 IAAA AAAA
	Load low word of ACC, short immediate	1	1	1011 1001 IIII IIII
LACT	Load ACC with shift (0 to 15) specified by TREG, direct or indirect	1	1	0110 1011 IAAA AAAA
NEG	Negate ACC	1	1	1011 1110 0000 0010
NORM	Normalize the contents of ACC, indirect	1	1	1010 0000 IAAA AAAA
OR	OR ACC with data value, direct or indirect	1	1	0110 1101 IAAA AAAA

Table 8–1. Accumulator, Arithmetic, and Logic Instructions (Continued)

Mnemonic	Description	Words	Cycles	Opcode
	OR with ACC with shift of 0 to 15, long immediate	2	2	1011 1111 1100 SHFT + 1 word
	OR with ACC with shift of 16, long immediate	2	2	1011 1110 1000 0010 + 1 word
ROL	Rotate ACC left	1	1	1011 1110 0000 1100
ROR	Rotate ACC right	1	1	1011 1110 0000 1101
SACH	Store high ACC with shift of 0 to 7, direct or indirect	1	1	1001 1SHF IAAA AAAA
SACL	Store low ACC with shift of 0 to 7, direct or indirect	1	1	1001 0SHF IAAA AAAA
SFL	Shift ACC left	1	1	1011 1110 0000 1001
SFR	Shift ACC right	1	1	1011 1110 0000 1010
SUB	Subtract from ACC with shift of 0 to 15, direct or indirect	1	1	0011 SHFT IAAA AAAA
	Subtract from ACC with shift of 0 to 15, long immediate	2	2	1011 1111 1010 SHFT + 1 word
	Subtract from ACC with shift of 16, direct or indirect	1	1	0110 0101 IAAA AAAA
	Subtract from ACC, short immediate	1	1	1011 1010 IIII IIII
SUBB	Subtract from ACC with borrow, direct or indirect	1	1	0110 0100 IAAA AAAA
SUBC	Conditional subtract, direct or indirect	1	1	0000 1010 IAAA AAAA
SUBS	Subtract from ACC with sign-extension suppressed, direct or indirect	1	1	0110 0110 IAAA AAAA
SUBT	Subtract from ACC with shift (0 to 15) specified by TREG, direct or indirect	1	1	0110 0111 IAAA AAAA
XOR	Exclusive OR ACC with data value, direct or indirect	1	1	0110 1100 IAAA AAAA
	Exclusive OR with ACC with shift of 0 to 15, long immediate	2	2	1011 1111 1101 SHFT + 1 word
	Exclusive OR with ACC with shift of 16, long immediate	2	2	1011 1110 1000 0011 + 1 word
ZALR	Zero low ACC and load high ACC with rounding, direct or indirect	1	1	0110 1000 IAAA AAAA

Table 8–2. Auxiliary Register Instructions

Mnemonic	Description	Words	Cycles	Opcode
ADRK	Add constant to current AR, short immediate	1	1	0111 1000 IIII IIII
BANZ	Branch on current AR not 0, indirect	2	4 (condition true) 2 (condition false)	0111 1011 1AAA AAAA + 1 word
CMPR	Compare current AR with AR0	1	1	1011 1111 0100 01CM
LAR	Load specified AR from specified data location, direct or indirect	1	2	0000 0ARX IAAA AAAA
	Load specified AR with constant, short immediate	1	2	1011 0ARX IIII IIII
	Load specified AR with constant, long immediate	2	2	1011 1111 0000 1ARX + 1 word
MAR	Modify current AR and/or ARP, indirect (performs no operation when direct)	1	1	1000 1011 IAAA AAAA
SAR	Store specified AR to specified data location, direct or indirect	1	1	1000 0ARX IAAA AAAA
SBRK	Subtract constant from current AR, short immediate	1	1	0111 1100 IIII IIII

Table 8–3. TREG, PREG, and Multiply Instructions

Mnemonic	Description	Words	Cycles	Opcode
APAC	Add PREG to ACC	1	1	1011 1110 0000 0100
LPH	Load high PREG, direct or indirect	1	1	0111 0101 IAAA AAAA
LT	Load TREG, direct or indirect	1	1	0111 0011 IAAA AAAA
LTA	Load TREG and accumulate previous product, direct or indirect	1	1	0111 0000 IAAA AAAA
LTD	Load TREG, accumulate previous product, and move data, direct or indirect	1	1	0111 0010 IAAA AAAA
LTP	Load TREG and store PREG in accumulator, direct or indirect	1	1	0111 0001 IAAA AAAA
LTS	Load TREG and subtract previous product, direct or indirect	1	1	0111 0100 IAAA AAAA
MAC	Multiply and accumulate, direct or indirect	2	3	1010 0010 IAAA AAAA + 1 word
MACD	Multiply and accumulate with data move, direct or indirect	2	3	1010 0011 IAAA AAAA + 1 word
MPY	Multiply TREG by data value, direct or indirect	1	1	0101 0100 IAAA AAAA
	Multiply TREG by 13-bit constant, short immediate	1	1	110I IIII IIII IIII
MPYA	Multiply and accumulate previous product, direct or indirect	1	1	0101 0000 IAAA AAAA
MPYS	Multiply and subtract previous product, direct or indirect	1	1	0101 0001 IAAA AAAA
MPYU	Multiply unsigned, direct or indirect	1	1	0101 0101 IAAA AAAA
PAC	Load ACC with PREG	1	1	1011 1110 0000 0011
SPAC	Subtract PREG from ACC	1	1	1011 1110 0000 0101
SPH	Store high PREG, direct or indirect	1	1	1000 1101 IAAA AAAA
SPL	Store low PREG, direct or indirect	1	1	1000 1100 IAAA AAAA
SPM	Set product shift mode	1	1	1011 1111 0000 00PM
SQRA	Square and accumulate previous product, direct or indirect	1	1	0101 0010 IAAA AAAA
SQRS	Square and subtract previous product, direct or indirect	1	1	0101 0011 IAAA AAAA

Table 8–4. Branch Instructions

Mnemonic	Description	Words	Cycles	Opcode
B	Branch unconditionally, indirect	2	4	0111 1001 1AAA AAAA + 1 word
BACC	Branch to address specified by ACC	1	4	1011 1110 0010 0000
BANZ	Branch on current AR not 0, indirect	2	4 (condition true) 2 (condition false)	0111 1011 1AAA AAAA + 1 word
BCND	Branch conditionally	2	4 (conditions true) 2 (any condition false)	1110 00TP ZLVC ZLVC + 1 word
CALA	Call subroutine at location specified by ACC	1	4	1011 1110 0011 0000
CALL	Call subroutine, indirect	2	4	0111 1010 1AAA AAAA + 1 word
CC	Call conditionally	2	4 (conditions true) 2 (any condition false)	1110 10TP ZLVC ZLVC + 1 word
INTR	Soft interrupt	1	4	1011 1110 011I NTR#
NMI	Nonmaskable interrupt	1	4	1011 1110 0101 0010
RET	Return from subroutine	1	4	1110 1111 0000 0000
RETC	Return conditionally	1	4 (conditions true) 2 (any condition false)	1110 11TP ZLVC ZLVC
TRAP	Software interrupt	1	4	1011 1110 0101 0001

Table 8–5. Control Instructions

Mnemonic	Description	Words	Cycles	Opcode
BIT	Test bit, direct or indirect	1	1	0100 BITX IAAA AAAA
BITT	Test bit specified by TREG, direct or indirect	1	1	0110 1111 IAAA AAAA
CLRC	Clear C bit	1	1	1011 1110 0100 1110
	Clear CNF bit	1	1	1011 1110 0100 0100
	Clear INTM bit	1	1	1011 1110 0100 0000
	Clear OVM bit	1	1	1011 1110 0100 0010
	Clear SXM bit	1	1	1011 1110 0100 0110
	Clear TC bit	1	1	1011 1110 0100 1010
	Clear XF bit	1	1	1011 1110 0100 1100
IDLE	Idle until interrupt	1	1	1011 1110 0010 0010
LDP	Load data page pointer, direct or indirect	1	2	0000 1101 IAAA AAAA
	Load data page pointer, short immediate	1	2	1011 110I IIII IIII
LST	Load status register ST0, direct or indirect	1	2	0000 1110 IAAA AAAA
	Load status register ST1, direct or indirect	1	2	0000 1111 IAAA AAAA
NOP	No operation	1	1	1000 1011 0000 0000
POP	Pop top of stack to low ACC	1	1	1011 1110 0011 0010
POPD	Pop top of stack to data memory, direct or indirect	1	1	1000 1010 IAAA AAAA
PSHD	Push data memory value on stack, direct or indirect	1	1	0111 0110 IAAA AAAA
PUSH	Push low ACC onto stack	1	1	1011 1110 0011 1100
RPT	Repeat next instruction, direct or indirect	1	1	0000 1011 IAAA AAAA
	Repeat next instruction, short immediate	1	1	1011 1011 IIII IIII
SETC	Set C bit	1	1	1011 1110 0100 1111
	Set CNF bit	1	1	1011 1110 0100 0101
	Set INTM bit	1	1	1011 1110 0100 0001
	Set OVM bit	1	1	1011 1110 0100 0011
	Set SXM bit	1	1	1011 1110 0100 0111
	Set TC bit	1	1	1011 1110 0100 1011
	Set XF bit	1	1	1011 1110 0100 1101
SPM	Set product shift mode	1	1	1011 1111 0000 00PM
SST	Store status register ST0, direct or indirect	1	1	1000 1110 IAAA AAAA
	Store status register ST1, direct or indirect	1	1	1000 1111 IAAA AAAA

Table 8–6. I/O and Memory Instructions

Mnemonic	Description	Words	Cycles	Opcode
BLDD	Block move from data memory to data memory, direct/indirect with long immediate source	2	3	1010 1000 IAAA AAAA + 1 word
	Block move from data memory to data memory, direct/indirect with long immediate destination	2	3	1010 1001 IAAA AAAA + 1 word
BLPD	Block move from program memory to data memory, direct/indirect with long immediate source	2	3	1010 0101 IAAA AAAA + 1 word
DMOV	Data move in data memory, direct or indirect	1	1	0111 0111 IAAA AAAA
IN	Input data from I/O location, direct or indirect	2	2	1010 1111 IAAA AAAA + 1 word
OUT	Output data to port, direct or indirect	2	3	0000 1100 IAAA AAAA + 1 word
SPLK	Store long immediate to data memory location, direct or indirect	2	2	1010 1110 IAAA AAAA + 1 word
TBLR	Table read, direct or indirect	1	3	1010 0110 IAAA AAAA
TBLW	Table write, direct or indirect	1	3	1010 0111 IAAA AAAA

8.2 How To Use the Instruction Descriptions

Section 8.3 contains detailed information on the instruction set. The description for each instruction presents the following categories of information:

- ❏ Syntax
- ❏ Operands
- ❏ Opcode
- ❏ Execution
- ❏ Status Bits
- ❏ Description
- ❏ Words
- ❏ Cycles
- ❏ Examples

8.2.1 Syntax

Each instruction begins with a list of the available assembler syntax expressions and the addressing mode type(s) for each expression. For example, the description for the ADD instruction begins with:

ADD *dma* [, *shift*]	Direct addressing
ADD *dma*, **16**	Direct with left shift of 16
ADD *ind* [, *shift* [, **AR***n*]]	Indirect addressing
ADD *ind*, **16** [, **AR***n*]	Indirect with left shift of 16
ADD #*k*	Short immediate addressing
ADD #*lk* [, *shift*]	Long immediate addressing

These are the notations used in the syntax expressions:

italic symbols

Italic symbols in an instruction syntax represent variables.

Example: For the syntax
ADD *dma*
you may use a variety of values for *dma*. Samples with this syntax follow:

```
ADD DAT
ADD 15
```

boldface characters

Boldface characters in an instruction syntax must be typed as shown.

Example: For the syntax
ADD *dma*, **16**
you may use a variety of values for *dma*, but the word ADD and the number 16 must be typed as shown. Samples with this syntax follow:

```
ADD 7h, 16
ADD X, 16
```

[, x] Operand x is optional.

Example: For the syntax
ADD *dma*, [, *shift*]
you must supply *dma*, as in the instruction:
`ADD 7h`
and you have the option of adding a *shift* value, as in the instruction:
`ADD 7h, 5`

[, x1 [, x2]] Operands x1 and x2 are optional, but you cannot include x2 without also including x1.

Example: For the syntax
ADD *ind*, [, *shift* [, **AR***n*]]
you must supply *ind*, as in the instruction:
`ADD *+`
You have the option of including *shift*, as in the instruction:
`ADD *+, 5`
If you wish to include **AR***n*, you must also include *shift*, as in:
`ADD *+, 0, AR2`

The # symbol is a prefix for constants used in immediate addressing. For short- or long- immediate operands, it is used in instructions where there is ambiguity with other addressing modes.

Example: `RPT #15` uses short immediate addressing. It causes the next instruction to be repeated 16 times. But `RPT 15` uses direct addressing. The number of times the next instruction repeats is determined by a value stored in memory.

Finally, consider this code example:

```
MoveData  BLDD   DAT5, #310h  ;move data at address
                               ;referenced by DAT5 to address
                               ;310h.
```

Note the optional MoveData label is used as a reference in front of the instruction mnemonic. Place labels either before the instruction mnemonic on the same line or on the preceding line in the first column. (Be sure there are no spaces in your labels.) An optional comment field can conclude the syntax expression. At least one space is required between fields (label, mnemonic, operand, and comment).

8.2.2 Operands

Operands can be constants, or assembly-time expressions referring to memory, I/O ports, register addresses, pointers, shift counts, and a variety of other constants. The operands category for each instruction description defines the variables used for and/or within operands in the syntax expressions. For example, for the ADD instruction, the syntax category gives these syntax expressions:

ADD *dma* [, *shift*]	Direct addressing
ADD *dma*, **16**	Direct with left shift of 16
ADD *ind* [, *shift* [, **AR***n*]]	Indirect addressing
ADD *ind*, **16** [, **AR***n*]	Indirect with left shift of 16
ADD #*k*	Short immediate addressing
ADD #*lk* [, *shift*]	Long immediate addressing

The operands category defines the variables *dma*, *shift*, *ind*, *n*, *k*, and *lk*. For *ind*, an indirect addressing variable, you supply one of the following seven symbols:

* *+ *– *0+ *0– *BR0+ *BR0–

These symbols are defined in subsection 7.3.2, *Indirect Addressing Options*, on page 7-9.

8.2.3 Opcode

The opcode category breaks down the various bit fields that make up each instruction word. When one of the fields contains a constant value derived directly from an operand, it has the same name as that operand. The contents of fields that do not directly relate to operands have other names; the opcode category either explains these names directly or refers you to a section of this book that explains them in detail. For example, these opcodes are given for the ADDC instruction:

ADDC *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	0	0	0	dma						

ADDC *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in Section 7.3, *Indirect Addressing Mode* (page 7-9).

The field called dma contains the value *dma*, which is defined in the operands category. The contents of the fields ARU, N, and NAR are derived from the operands *ind* and *n* but do not directly correspond to those operands; therefore, a note directs you to the appropriate section for more details.

8.2.4 Execution

The execution category presents an instruction operation sequence that describes the processing that takes place when the instruction is executed. If the execution event or events depend on the addressing mode used, the execution category specifies which events are associated with which addressing modes. Here are notations used in the execution category:

(r)	The content of register or location r. *Example:* (ACC) represents the value in the accumulator.
x → y	Value x is assigned to register or location y. *Example:* (data-memory address) → ACC means: The content of the specified data-memory address is put into the accumulator.
r(n:m)	Bits n through m of register or location r. *Example:* ACC(15:0) represents bits 15 through 0 of the accumulator.
(r(n:m))	The content of bits n through m of register or location r. *Example:* (ACC(31:16)) represents the content of bits 31 through 16 of the accumulator.
nnh	Indicates that nn represents a hexadecimal number.

8.2.5 Status Bits

The bits in status registers ST0 and ST1 affect the operation of certain instructions and are affected by certain instructions. The status bits category of each instruction description states which of the bits (if any) affect the execution of the instruction and which of the bits (if any) are affected by the instruction.

8.2.6 Description

The description category explains what happens during instruction execution and its effect on the rest of the processor or on memory contents. It also discusses any constraints on the operands imposed by the processor or the assembler. This description parallels and supplements the information given in the execution category.

8.2.7 Words

The words category specifies the number of memory words required to store the instruction (one or two). When the number of words depends on the addressing mode used for an instruction, the words category specifies which addressing modes require one word and which require two words.

8.2.8 Cycles

The cycles category of each instruction description contains tables showing the number of processor machine cycles (CLKOUT1 periods) required for the instruction to execute in a given memory configuration when executed as a single instruction or when repeated with the RPT instruction. For example:

Cycles for a Single Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1	1+p
External	1+d	1+d	1+d	2+d+p

Cycles for a Repeat (RPT) Execution of an Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

The column headings in these tables indicate the program source location, defined as follows:

ROM	The instruction executes from internal program ROM.
DARAM	The instruction executes from internal dual-access program RAM.
SARAM	The instruction executes from internal single-access program RAM.
External	The instruction executes from external program memory.

If an instruction requires memory operand(s), the rows in the table indicate the location(s) of the operand(s), as defined here:

DARAM	The operand is in internal dual-access RAM.
SARAM	The operand is in internal single-access RAM.
External	The operand is in external memory.

For the RPT mode execution, *n* indicates the number of times a given instruction is repeated by an RPT instruction. Additional cycles (wait states) can be generated for program-memory, data-memory, and I/O accesses by the wait-state generator or by the external READY signal. These additional wait states are represented in the tables by the following variables:

p	Program-memory wait states. Represents the number of additional clock cycles the device waits for external program memory to respond to a single access.
d	Data-memory wait states. Represents the number of additional clock cycles the device waits for external data memory to respond to a single access.
io	I/O wait states. Represents the number of additional clock cycles the device waits for an external I/O device to respond to a single access.
n	Number of repetitions (where $n > 2$ to fill the pipeline). Represents the number of times a repeated instruction is executed.

If there are multiple accesses to one of the spaces, the variable is preceded by the appropriate integer multiple. For example, two accesses to external program memory would require 2p wait states. The above variables may also use the subscripts *src, dst,* and *code* to indicate source, destination, and code, respectively.

The internal single-access memory on each ’C24x processor is divided into 2K-word blocks contiguous in address space. All ’C24x processors support parallel accesses to these internal single-access RAM blocks. Furthermore, one single access block allows only one access per cycle. Thus, the processor can read/write on single-access RAM block while accessing another single-access RAM block at the same time.

All external reads take at least one machine cycle while all external writes take at least two machine cycles. However, if an external write is immediately followed or preceded by an external read cycle, then the external write requires three cycles. If the wait state generator or the READY pin is used to add *m* (*m* > 0) wait states to an external access, then external reads require *m* + 1 cycles, and external write accesses require *m* + 2 cycles.

The instruction-cycle timings are based on the following assumptions:

- ❑ At least the next four instructions are fetched from the same memory section (internal or external) that was used to fetch the current instruction (except in the case of PC discontinuity instructions, such as B, CALL, etc.)
- ❑ In the single-execution mode, there is no pipeline conflict between the current instruction and the instructions immediately preceding or following that instruction. The only exception is the conflict between the fetch phase of the pipeline and the memory read/write (if any) access of the instruction under consideration. See Section 5.2, *Pipeline Operation*, on page 5-7 for more information about pipeline operations.
- ❑ In the repeat execution mode, all conflicts caused by the pipelined execution of an instruction are considered.

8.2.9 Examples

Example code is included for each instruction. The effect of the code on memory and/or registers is summarized. Consider this example of the ADD instruction:

```
ADD *+,0,AR0
```

	Before Instruction		After Instruction
ARP	4	ARP	0
AR4	0302h	AR4	0303h
Data Memory 302h	2h	Data Memory 302h	2h
ACC	2h	ACC	04h
C	X	C	0

Here are the facts and events represented in this example:

- ❑ The auxiliary register pointer (ARP) points to the current auxiliary register. Because ARP = 4, the current auxiliary register is AR4.
- ❑ When the addition takes place, the CPU follows AR4 to data-memory address 0302h. The content of that address, 2h, is added to the content of the accumulator, also 2h. The result (4h) is placed in the accumulator. (Because the second operand of the instruction specifies a left shift of 0, the data-memory value is not shifted before being added to the accumulator value.)

- The instruction specifies an increment of 1 for the contents of the current auxiliary register (*+); therefore, after the addition is performed, the content of AR4 is incremented to 0303h.
- The instruction also specifies that AR0 is the next auxiliary register; therefore, after the instruction ARP = 0.
- Because no carry is generated during the addition, the carry bit (C) is cleared to 0.

8.3 Instruction Descriptions

This section contains detailed information on the instruction set for the 'C24x. A summary of the instruction set is shown in Section 8.1 on page 8-2. The instructions are presented alphabetically, and the description for each instruction presents the following categories of information:

- Syntax
- Operands
- Opcode
- Execution
- Status Bits
- Description
- Words
- Cycles
- Examples

For a description of how to use each of these categories, see Section 8.2 on page 8-12.

Syntax **ABS**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	0	0	0	0

Execution Increment PC, then ...
$|(ACC)| \rightarrow ACC; 0 \rightarrow C$

Status Bits

Affected by	*Affects*
OVM	C and OV

This instruction is not affected by SXM

Description If the contents of the accumulator are greater than or equal to 0, the accumulator is unchanged by the execution of ABS. If the contents of the accumulator are less than 0, the accumulator is replaced by its 2s-complement value. The carry bit (C) on the 'C24x is always reset to 0 by the execution of this instruction.

Note that 8000 0000h is a special case. When the overflow mode is not set (OVM = 0), the ABS of 8000 0000h is 8000 0000h. When the overflow mode is set (OVM = 1), the ABS of 8000 0000h is 7FFF FFFFh. In either case, the OV status bit is set.

Words 1

Cycles

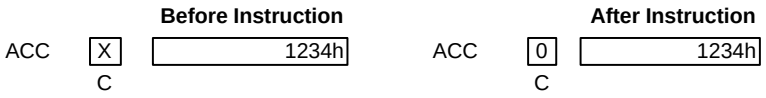
Cycles for a Single ABS Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

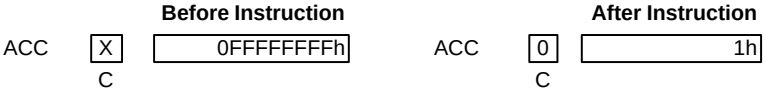
Cycles for a Repeat (RPT) Execution of an ABS Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

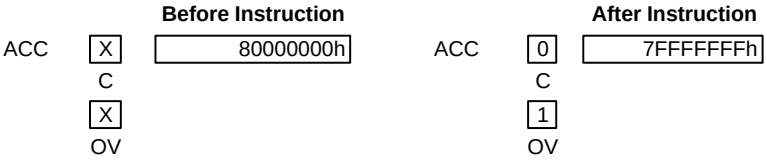
Example 1 ABS



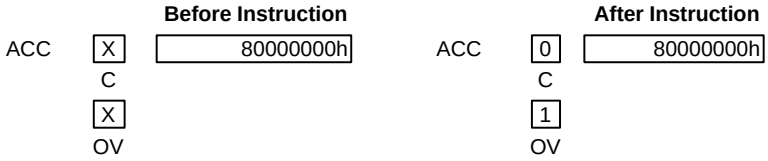
Example 2 ABS



Example 3 ABS ;(OVM = 1)



Example 4 ABS ;(OVM = 0)



Syntax

ADD *dma* [, *shift*]	Direct addressing
ADD *dma*, **16**	Direct with left shift of 16
ADD *ind* [, *shift* [, **AR***n*]]	Indirect addressing
ADD *ind*, **16** [, **AR***n*]	Indirect with left shift of 16
ADD #*k*	Short immediate addressing
ADD #*lk* [, *shift*]	Long immediate addressing

Operands

dma:	7 LSBs of the data-memory address
shift:	Left shift value from 0 to 15 (defaults to 0)
n:	Value from 0 to 7 designating the next auxiliary register
k:	8-bit short immediate value
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

ADD *dma* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	shift				0	dma						

ADD *dma*, **16**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	0	1	0	dma						

ADD *ind* [, *shift* [, **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	shift				1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

ADD *ind*, **16** [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode* Format (page 7-12).

ADD #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	0	0	0	k							

ADD #*lk* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	1	0	0	1	shift			
lk															

Execution

Increment PC, then ...

Event	*Addressing mode*
(ACC) + ((data-memory address) $\times$ 2^{shift}) $\rightarrow$ ACC	Direct or indirect
(ACC) + ((data-memory address) $\times$ 2^{16}) $\rightarrow$ ACC	Direct or indirect (shift of 16)
(ACC) + k $\rightarrow$ ACC	Short immediate
(ACC) + lk $\times$ 2^{shift} $\rightarrow$ ACC	Long immediate

Status Bits

Affected by	*Affects*	*Addressing mode*
SXM and OVM	C and OV	Direct or indirect
OVM	C and OV	Short immediate
SXM and OVM	C and OV	Long immediate

Description

The content of the addressed data memory location or an immediate constant is left-shifted and added to the accumulator. During shifting, low-order bits are zero filled. High-order bits are sign extended if SXM = 1 and zero filled if SXM = 0. The result is stored in the accumulator. When short immediate addressing is used, the addition is unaffected by SXM and is not repeatable.

If you are using indirect addressing and update the ARP, you must specify a shift operand. However, if you do not want a shift to occur, enter a 0 for this operand. For example:

```
ADD *+,0,AR2
```

Normally, the carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry. However, when adding with a shift of 16, the carry bit is set if a carry is generated but otherwise, the carry bit is unaffected. This allows the accumulator to generate the proper single carry when adding a 32-bit number to the accumulator.

Words

Words	*Addressing mode*
1	Direct, indirect, or short immediate
2	Long immediate

Cycles

Cycles for a Single ADD Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an ADD Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Cycles for a Single ADD Instruction (Using Short Immediate Addressing)

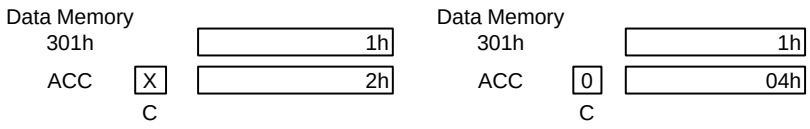
ROM	**DARAM**	**SARAM**	**External**
1	1	1	1+p

Cycles for a Single ADD Instruction (Using Long Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+2p

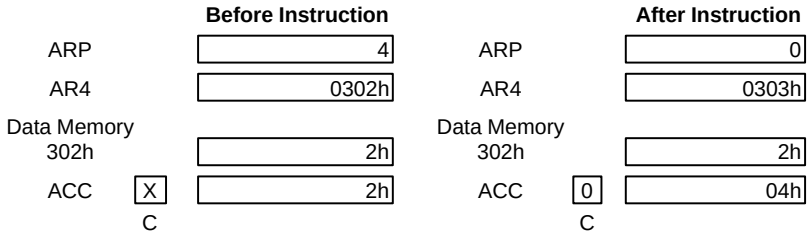
Example 1

```
ADD     1,1         ;(DP = 6)
```



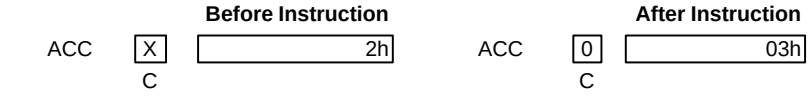
Example 2

```
ADD     *+,0,AR0
```



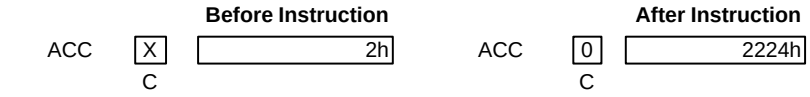
Example 3

```
ADD     #1h         ;Add short immediate
```



Example 4

```
ADD     #1111h,1  ;Add long immediate with shift of 1
```



Syntax

ADDC *dma* Direct addressing
ADDC *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

ADDC *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	0	0	0	dma						

ADDC *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) + (data-memory address) + (C) → ACC

Status Bits

Affected by	*Affects*
OVM	C and OV

This instruction is not affected by SXM.

Description

The contents of the addressed data-memory location and the value of the carry bit are added to the accumulator with sign extension suppressed. The carry bit is then affected in the normal manner: the carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry.

The ADDC instruction can be used in performing multiple-precision arithmetic.

Words

1

Cycles

Cycles for a Single ADDC Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

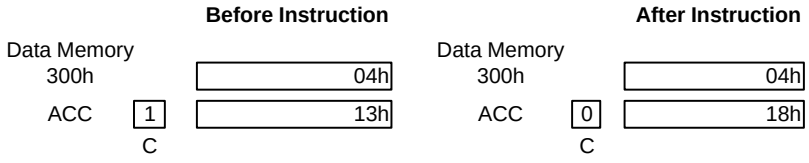
Cycles for a Repeat (RPT) Execution of an ADDC Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

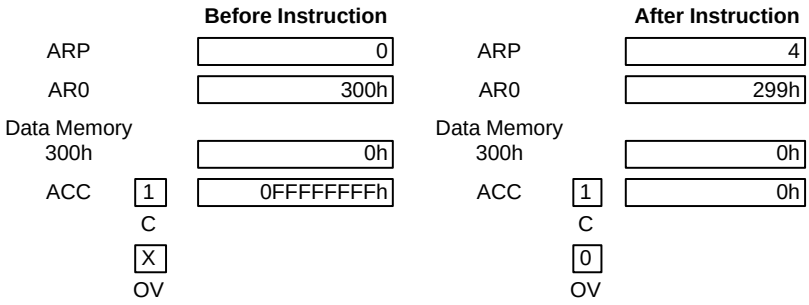
Example 1

```
ADDC    DAT300    ;(DP = 6: addresses 0300h–037Fh;
                  ;DAT300 is a label for 300h)
```



Example 2

```
ADDC    *–,AR4    ;(OVM = 0)
```



Syntax

ADDS *dma* Direct addressing
ADDS *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

ADDS *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	1	0	0	dma						

ADDS *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) + (data-memory address) → ACC

Status Bits

Affected by	*Affects*
OVM	C and OV

This instruction is not affected by SXM.

Description

The contents of the specified data-memory location are added to the accumulator with sign extension suppressed. The data is treated as an unsigned 16-bit number, regardless of SXM. The accumulator contents are treated as a signed number. Note that ADDS produces the same results as an ADD instruction with SXM = 0 and a shift count of 0.

The carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry.

Words

1

Cycles

Cycles for a Single ADDS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an ADDS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
ADDS    0          ;(DP = 6: addresses 0300h–037Fh)
```

	Before Instruction		After Instruction	
Data Memory 300h		0F006h		0F006h
ACC	X	00000003h	0	0000F009h
	C		C	

Example 2

```
ADDS    *
```

	Before Instruction		After Instruction	
ARP		0		0
AR0		0300h		0300h
Data Memory 300h		0FFFFh		0FFFFh
ACC	X	7FFF0000h	0	7FFFFFFFh
	C		C	

Syntax

ADDT *dma* Direct addressing
ADDT *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

ADDT *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	1	1	0	dma						

ADDT *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) + [(data-memory address) $\times$ $2^{(TREG(3:0))}$] $\rightarrow$ (ACC)

Status Bits

Affected by	*Affects*
SXM or OVM	C and OV

Description

The data-memory value is left shifted and added to the accumulator, and the result replaces the accumulator contents. The left shift is defined by the four LSBs of the TREG, resulting in shift options from 0 to 15 bits. Sign extension on the data-memory value is controlled by SXM. The carry bit (C) is set when a carry is generated out of the MSB of the accumulator; if no carry is generated, the carry bit is cleared.

Words

1

Cycles

Cycles for a Single ADDT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block.

Cycles for a Repeat (RPT) Execution of an ADDT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
ADDT    127     ;(DP = 4: addresses 0200h–027Fh,
                ;SXM = 0)
```

	Before Instruction		After Instruction	
Data Memory 027Fh	09h		09h	
TREG	0FF94h		0FF94h	
ACC	0F715h	X	0F7A5h	0
		C		C

Example 2

```
ADDT    *–,AR4  ;(SXM = 0)
```

	Before Instruction		After Instruction	
ARP	0		4	
AR0	027Fh		027Eh	
Data Memory 027Fh	09h		09h	
TREG	0FF94h		0FF94h	
ACC	0F715h	X	0F7A5h	0
		C		C

Syntax **ADRK** #*k* Short immediate addressing

Operands k: 8-bit short immediate value

Opcode

ADRK #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	0	0	k							

Execution

Increment PC, then ...
(current AR) + 8-bit positive constant → current AR

Status Bits None

Description The 8-bit immediate value is added, right justified, to the current auxiliary register (the one specified by the current ARP value) and the result replaces the auxiliary register contents. The addition takes place in the ARAU, with the immediate value treated as an 8-bit positive integer. All arithmetic operations on the auxiliary registers are unsigned.

Words 1

Cycles

Cycles for a Single ADRK Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Example

```
ADRK    #80h
```

	Before Instruction		After Instruction
ARP	5	ARP	5
AR5	4321h	AR5	43A1h

Syntax

AND *dma*	Direct addressing
AND *ind* [, **AR***n*]	Indirect addressing
AND #*lk* [, *shift*]	Long immediate addressing
AND #*lk*, **16**	Long immediate with left shift of 16

Operands

dma:	7 LSBs of the data-memory address
shift:	Left shift value from 0 to 15 (defaults to 0)
n:	Value from 0 to 7 designating the next auxiliary register
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

AND *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	1	0	0	dma						

AND *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode* Format (page 7-12).

AND #*lk* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	1	0	1	1	shift			
lk															

AND #*lk*, **16**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	1	0	0	0	0	0	0	1
lk															

Execution

Increment PC, then ...

Event(s)	*Addressing mode*
(ACC(15:0)) AND (data-memory address) $\rightarrow$ ACC(15:0) 0 $\rightarrow$ ACC(31:16)	Direct or indirect
(ACC(31:0)) AND lk $\times$ 2^{shift} $\rightarrow$ ACC	Long immediate
(ACC(31:0)) AND lk $\times$ 2^{16} $\rightarrow$ ACC	Long immediate with left shift of 16

Status Bits None

This instruction is not affected by SXM.

Description

If direct or indirect addressing is used, the low word of the accumulator is ANDed with a data-memory value, and the result is placed in the low word position in the accumulator. The high word of the accumulator is zeroed. If immediate addressing is used, the long-immediate constant can be shifted. During the shift, low-order and high-order bits not filled by the shifted value are zeroed. The resulting value is ANDed with the accumulator contents.

Words

Words	*Addressing mode*
1	Direct or indirect
2	Long immediate

Cycles

Cycles for a Single AND Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an AND Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Cycles for a Single AND Instruction (Using Long Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+2p

Example 1

```
AND     16          ;(DP = 4: addresses 0200h–027Fh)
```

	Before Instruction		After Instruction
Data Memory 0210h	00FFh	Data Memory 0210h	00FFh
ACC	12345678h	ACC	00000078h

Example 2

```
AND     *
```

	Before Instruction		After Instruction
ARP	0	ARP	0
AR0	0301h	AR0	0301h
Data Memory 0301h	0FF00h	Data Memory 0301h	0FF00h
ACC	12345678h	ACC	00005600h

Example 3

```
AND     #00FFh,4
```

	Before Instruction		After Instruction
ACC	12345678h	ACC	00000670h

Syntax **APAC**

Operands None

Opcode

APAC

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	0	1	0	0

Execution Increment PC, then ...
(ACC) + shifted (PREG) → ACC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

This instruction is not affected by SXM.

Description The contents of PREG are shifted as defined by the PM status bits of the ST1 register (see Table 8–7) and added to the contents of the accumulator. The result is placed in the accumulator. APAC is not affected by the SXM bit of the status register. PREG is always sign extended. The task of the APAC instruction is also performed as a subtask of the LTA, LTD, MAC, MACD, MPYA, and SQRA instructions.

Table 8–7. Product Shift Modes

PM Bits		
Bit 1	**Bit 0**	**Resulting Shift**
0	0	No shift
0	1	Left shift of 1 bit
1	0	Left shift of 4 bits
1	1	Right shift of 6 bits

Words 1

Cycles

Cycles for a Single APAC Instruction

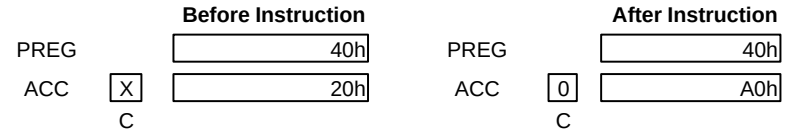
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an APAC Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
APAC        ;(PM = 01)
```



Syntax **B** *pma* [, *ind* [, **AR***n*]] Indirect addressing

Operands

pma: 16-bit program-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

B *pma* [, *ind* [, **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	0	1	1	ARU			N	NAR		
pma															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

pma → PC
Modify (current AR) and (ARP) as specified.

Status Bits None

Description The current auxiliary register and ARP contents are modified as specified, and control is passed to the designated program-memory address (pma). The pma can be either a symbolic or numeric address.

Words 2

Cycles

Cycles for a Single B Instruction

ROM	DARAM	SARAM	External
4	4	4	4+4p

Note: When this instruction reaches the execute phase of the pipeline, two additional instruction words have entered the pipeline. When the PC discontinuity is taken, these two instruction words are discarded.

Example

```
B       191,*+,AR1
```

The value 191 is loaded into the program counter, and the program continues to execute from that location. The current auxiliary register is incremented by 1, and ARP is set to point to auxiliary register 1 (AR1).

Syntax **BACC**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	1	0	0	0	0	0

Execution ACC(15:0) → PC

Status Bits None

Description Control is passed to the 16-bit address residing in the lower half (16 LSBs) of the accumulator.

Words 1

Cycles

Cycles for a Single BACC Instruction

ROM	DARAM	SARAM	External
4	4	4	4+3p

Note: When this instruction reaches the execute phase of the pipeline, two additional instruction words have entered the pipeline. When the PC discontinuity is taken, these two instruction words are discarded.

Example

```
BACC      ;(ACC contains the value 191)
```

The value 191 is loaded into the program counter, and the program continues to execute from that location.

Syntax

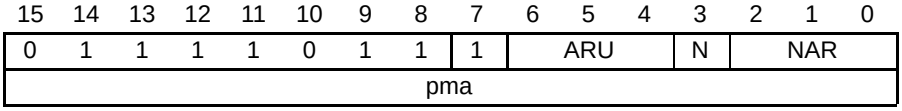
BANZ *pma* [, *ind* [, **AR***n*]] Indirect addressing

Operands

pma: 16-bit program-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

BANZ *pma* [, *ind* [,**AR***n*]]



Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

If (current AR) ≠ 0
Then pma → PC
Else (PC) + 2 → PC
Modify (current AR) and (ARP) as specified

Status Bits

None

Description

Control is passed to the designated program-memory address (pma) if the contents of the current auxiliary register are not 0. Otherwise, control passes to the next instruction.The default modification to the current AR is a decrement by 1. N loop iterations can be executed by initializing an auxiliary register (as a loop counter) to N–1 prior to loop entry. The pma can be either a symbolic or a numeric address.

Words

2

Cycles

Cycles for a Single BANZ Instruction

Condition	ROM	DARAM	SARAM	External
True	4	4	4	4+4p
False	2	2	2	2+2p

Note: The ’C24x performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example 1

```
BANZ      PGM0      ;(PGM0 labels program address 0)
```

	Before Instruction		After Instruction
ARP	0	ARP	0
AR0	5h	AR0	4h

Because the content of AR0 is not 0, program address 0 is loaded into the program counter (PC) and the program continues executing from that location. The default auxiliary register operation is a decrement of the current auxiliary register content; thus, AR0 contains 4h at the end of the execution.

or

	Before Instruction		After Instruction
ARP	0	ARP	0
AR0	0h	AR0	FFFFh

Because the content of AR0 is 0, the branch is not executed; instead, the PC is incremented by 2, and execution continues with the instruction following the BANZ instruction. Because of the default decrement, AR0 is decremented by 1, becoming –1.

Example 2

```
          MAR *,AR0       ;Set ARP to point to AR0.
          LAR AR1,#3      ;Load AR1 with 3.
          LAR AR0,#60h    ;Load AR0 with 60h.
PGM191    ADD *+,AR1      ;Loop: While AR1 not zero,
          BANZ PGM191,AR0 ;add data referenced by AR0
                          ;to accumulator and increment
                          ;AR0 value.
```

The contents of data-memory locations 60h–63h are added to the accumulator.

Syntax

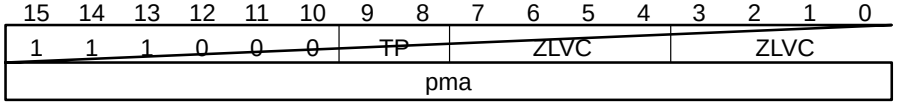
BCND *pma*, *cond* 1 [,*cond* 2] [,...]

Operands

pma: 16-bit program-memory address

cond	*Condition*
EQ	ACC = 0
NEQ	ACC ≠ 0
LT	ACC < 0
LEQ	ACC ≤ 0
GT	ACC > 0
GEQ	ACC ≥ 0
NC	C = 0
C	C = 1
NOV	OV = 0
OV	$\overline{\text{OV}}$ = 1
BIO	$\overline{\text{BIO}}$ low
NTC	TC = 0
TC	TC = 1
UNC	Unconditionally

Opcode



Note: The TP and ZLVC fields are defined on pages 8-3 and 8-4.

Execution

If *cond* 1 AND *cond* 2 AND ...
Then pma → PC
Else increment PC

Status Bits

None

Description

A branch is taken to the specified program-memory address (pma) if the specified conditions are met. Not all combinations of conditions are meaningful. For example, testing for LT and GT is contradictory. In addition, testing $\overline{\text{BIO}}$ is mutually exclusive to testing TC.

Words

2

Cycles

Cycles for a Single BCND Instruction

Condition	ROM	DARAM	SARAM	External
True	4	4	4	4+4p
False	2	2	2	2+2p

Note: The 'C24x performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example

```
BCND        PGM191,LEQ,C
```

If the accumulator contents are less than or equal to 0 and the carry bit is set, program address 191 is loaded into the program counter, and the program continues to execute from that location. If these conditions do not hold, execution continues from location PC + 2.

Syntax

BIT *dma*, *bit code* Direct addressing
BIT *ind*, *bit code* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
bit code: Value from 0 to 15 indicating which bit to test (see Figure 8–1)
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

BIT *dma*, *bit code*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	0	bit code				0	dma						

BIT *ind*, *bit code* [,**AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	0	bit code				1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data bit number (15 – bit code)) → TC

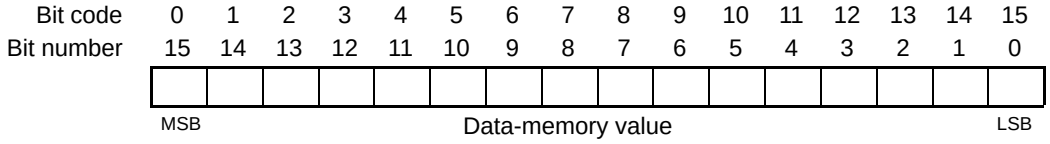
Status Bits

Affects
TC

Description

The BIT instruction copies the specified bit of the data-memory value to the TC bit of status register ST1. Note that the BITT, CMPR, LST #1, and NORM instructions also affect the TC bit in ST1. A bit code value is specified that corresponds to a certain bit number of the data-memory value, as shown in Figure 8–1.

Figure 8–1. Bit Numbers and Their Corresponding Bit Codes for BIT Instruction



Words

1

Cycles

Cycles for a Single BIT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a BIT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
BIT       0h,15     ;(DP = 6). Test LSB at 300h
```

	Before Instruction	After Instruction
Data Memory 300h	4DC8h	4DC8h
TC	0	0

Example 2

```
BIT       *,0,AR1   ;Test MSB at 310h, then set ARP = 1
```

	Before Instruction	After Instruction
ARP	0	1
AR0	310h	310h
Data Memory 310h	8000h	8000h
TC	0	1

Syntax

BITT *dma* Direct addressing
BITT *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

BITT *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	1	1	0	dma						

BITT *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data bit number (15 –TREG(3:0))) → TC

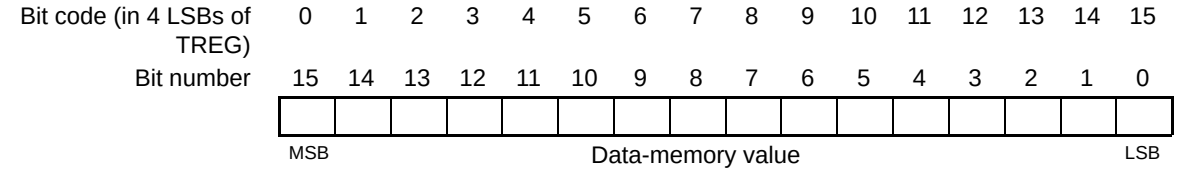
Status Bits

Affects
TC

Description

The BITT instruction copies the specified bit of the data-memory value to the TC bit of status register ST1. Note that the BITT, CMPR, LST #1, and NORM instructions also affect the TC bit in status register ST1. The bit number is specified by a bit code value contained in the four LSBs of the TREG, as shown in Figure 8–2.

Figure 8–2. Bit Numbers and Their Corresponding Bit Codes for BITT Instruction



Words

1

Cycles

Cycles for a Single BITT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an BITT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
BITT    00h     ;(DP = 6) Test bit 14 of data
                ;at 300h
```

	Before Instruction		After Instruction
Data Memory 300h	4DC8h	Data Memory 300h	4DC8h
TREG	1h	TREG	1h
TC	0	TC	1

Example 2

```
BITT    *       ;Test bit 1 of data at 310h
```

	Before Instruction		After Instruction
ARP	1	ARP	1
AR1	310h	AR1	310h
Data Memory 310h	8000h	Data Memory 310h	8000h
TREG	0Eh	TREG	0Eh
TC	0	TC	0

Syntax

General syntax: **BLDD** *source, destination*

BLDD #*lk*, *dma*	Direct with long immediate source
BLDD #*lk*, *ind* [, **AR***n*]	Indirect with long immediate source
BLDD *dma*, #*lk*	Direct with long immediate destination
BLDD *ind*, #*lk* [, **AR***n*]	Indirect with long immediate destination

Operands

dma:	7 LSBs of the data-memory address
n:	Value from 0 to 7 designating the next auxiliary register
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

BLDD #*lk*, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	0	0	0	0	dma						
lk															

BLDD #*lk*, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	0	0	0	1	ARU			N	NAR		
lk															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

BLDD *dma*, #*lk*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	0	0	1	0	dma						
lk															

BLDD *ind*, #*lk* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	0	0	1	1	ARU			N	NAR		
lk															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(PC) → MSTACK
lk → PC
(source) → destination
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC

While (repeat counter) ≠ 0:
    (source) → destination
    For indirect, modify (current AR) and (ARP) as specified
    (PC) + 1 → PC
    (repeat counter) –1 → repeat counter

(MSTACK) → PC

Status Bits

None

Description

The word in data memory pointed to by *source* is copied to a data-memory space pointed at by *destination*. The word of the source and/or destination space can be pointed at with a long-immediate value or by a data-memory address. Note that not all source/destination combinations of pointer types are valid.

Note:

BLDD does not work with core memory-mapped registers such as IMR, IFR, and GREG.

RPT can be used with the BLDD instruction to move consecutive words in data memory. The number of words to be moved is one greater than the number contained in the repeat counter (RPTC) at the beginning of the instruction. When the BLDD instruction is repeated, the source (destination) address specified by the long immediate constant is stored to the PC. Because the PC is incremented by 1 during each repetition, it is possible to access a series of source (destination) addresses. If you use indirect addressing to specify the destination (source) address, a new destination (source) address can be accessed during each repetition. If you use the direct addressing mode, the specified destination (source) address is a constant; it is not modified during each repetition.

The source and destination blocks do not have to be entirely on chip or off chip. Interrupts are inhibited during a BLDD operation used with the RPT instruction. When used with RPT, BLDD becomes a single-cycle instruction once the RPT pipeline is started.

Words 2

Cycles

Cycles for a Single BLDD Instruction

Operand	ROM	DARAM	SARAM	External
Source: DARAM Destination: DARAM	3	3	3	3+2p
Source: SARAM Destination: DARAM	3	3	3	3+2p
Source: External Destination: DARAM	3+d_{src}	3+d_{src}	3+d_{src}	3+d_{src}+2p
Source: DARAM Destination: SARAM	3	3	3 4†	3+2p
Source: SARAM Destination: SARAM	3	3	3 4†	3+2p
Source: External Destination: SARAM	3+d_{src}	3+d_{src}	3+d_{src} 4+d_{src}†	3+d_{src}+2p
Source: DARAM Destination: External	4+d_{dst}	4+d_{dst}	4+d_{dst}	6+d_{dst}+2p
Source: SARAM Destination: External	4+d_{dst}	4+d_{dst}	4+d_{dst}	6+d_{dst}+2p
Source: External Destination: External	4+d_{src}+d_{dst}	4+d_{src}+d_{dst}	4+d_{src}+d_{dst}	6+d_{src}+d_{dst}+2p

† If the destination operand and the code are in the same SARAM block.

Cycles for a Repeat (RPT) Execution of a BLDD Instruction

Operand	ROM	DARAM	SARAM	External
Source: DARAM Destination: DARAM	n+2	n+2	n+2	n+2+2p
Source: SARAM Destination: DARAM	n+2	n+2	n+2	n+2+2p
Source: External Destination: DARAM	n+2+nd_{src}	n+2+nd_{src}	n+2+nd_{src}	n+2+nd_{src}+2p
Source: DARAM Destination: SARAM	n+2	n+2	n+2 n+4†	n+2+2p
Source: SARAM Destination: SARAM	n+2 2n‡	n+2 2n‡	n+2 2n‡ n+4† 2n+2§	n+2+2p 2n+2p‡
Source: External Destination: SARAM	n+2+nd_{src}	n+2+nd_{src}	n+2+nd_{src} n+4+nd_{src}†	n+2+nd_{src}+2p
Source: DARAM Destination: External	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+2+nd_{dst}+2p
Source: SARAM Destination: External	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+2+nd_{dst}+2p
Source: External Destination: External	4n+nd_{src}+nd_{dst}‡	4n+nd_{src}+nd_{dst}	4n+nd_{src}+nd_{dst}	4n+2+nd_{src}+nd_{dst}+2p

† If the destination operand and the code are in the same SARAM block
‡ If both the source and the destination operands are in the same SARAM block
§ If both operands and the code are in the same SARAM block

Example 1

```
BLDD    #300h,20h ;(DP = 6)
```

	Before Instruction		After Instruction
Data Memory 300h	0h	Data Memory 300h	0h
320h	0Fh	320h	0h

Example 2

```
BLDD    *+,#321h,AR3
```

	Before Instruction		After Instruction
ARP	2	ARP	3
AR2	301h	AR2	302h
Data Memory 301h	01h	Data Memory 301h	01h
321h	0Fh	321h	01h

Syntax

General syntax: **BLPD** *source*, *destination*

BLPD #*pma*, *dma* — Direct with long immediate source

BLPD #*pma*, *ind* [, **AR***n*] — Indirect with long immediate source

Operands

pma: 16-bit program-memory address
dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

BLPD *#pma, dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	0	1	0	dma						
pma															

BLPD *#pma, ind* [*,* **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	0	1	1	ARU			N	NAR		
pma															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(PC) → MSTACK
pma → PC
(source) → destination
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC

While (repeat counter) ≠ 0:
    (source) → destination
    For indirect, modify (current AR) and (ARP) as specified
    (PC) + 1 → PC
    (repeat counter) –1 → repeat counter

(MSTACK) → PC

Status Bits

None

Description

A word in program memory pointed to by the *source* is copied to data-memory space pointed to by *destination*. The first word of the source space is pointed to by a long-immediate value. The data-memory destination space is pointed to by a data-memory address or auxiliary register pointer. Not all source/destination combinations of pointer types are valid.

RPT can be used with the BLPD instruction to move consecutive words. The number of words to be moved is one greater than the number contained in the repeat counter (RPTC) at the beginning of the instruction. When the BLPD instruction is repeated, the source (program-memory) address specified by the long immediate constant is stored to the PC. Because the PC is incremented by 1 during each repetition, it is possible to access a series of program-memory addresses. If you use indirect addressing to specify the destination (data-memory) address, a new data-memory address can be accessed during each repetition. If you use the direct addressing mode, the specified data-memory address is a constant; it is not modified during each repetition.

The source and destination blocks do not have to be entirely on chip or off chip. Interrupts are inhibited during a repeated BLPD instruction. When used with RPT, BLPD becomes a single-cycle instruction once the RPT pipeline is started.

Words

2

Cycles

Cycles for a Single BLPD Instruction

Operand	ROM	DARAM	SARAM	External
Source: DARAM/ROM Destination: DARAM	3	3	3	$3+2p_{code}$
Source: SARAM Destination: DARAM	3	3	3	$3+2p_{code}$
Source: External Destination: DARAM	$3+p_{src}$	$3+p_{src}$	$3+p_{src}$	$3+p_{src}+2p_{code}$
Source: DARAM/ROM Destination: SARAM	3	3	3 4†	$3+2p_{code}$
Source: SARAM Destination: SARAM	3	3	3 4†	$3+2p_{code}$
Source: External Destination: SARAM	$3+p_{src}$	$3+p_{src}$	$3+p_{src}$ $4+p_{src}$†	$3+p_{src}+2p_{code}$
Source: DARAM/ROM Destination: External	$4+d_{dst}$	$4+d_{dst}$	$4+d_{dst}$	$6+d_{dst}+2p_{code}$

† If the destination operand and the code are in the same SARAM block

Cycles for a Single BLPD Instruction (Continued)

Operand	ROM	DARAM	SARAM	External
Source: SARAM Destination: External	$4+d_{dst}$	$4+d_{dst}$	$4+d_{dst}$	$6+d_{dst}+2p_{code}$
Source: External Destination: External	$4+p_{src}+d_{dst}$	$4+p_{src}+d_{dst}$	$4+p_{src}+d_{dst}$	$6+p_{src}+d_{dst}+2p_{code}$

† If the destination operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a BLPD Instruction

Operand	ROM	DARAM	SARAM	External
Source: DARAM/ROM Destination: DARAM	n+2	n+2	n+2	$n+2+2p_{code}$
Source: SARAM Destination: DARAM	n+2	n+2	n+2	$n+2+2p_{code}$
Source: External Destination: DARAM	$n+2+np_{src}$	$n+2+np_{src}$	$n+2+np_{src}$	$n+2+np_{src}+2p_{code}$
Source: DARAM/ROM Destination: SARAM	n+2	n+2	n+2 n+4†	$n+2+2p_{code}$
Source: SARAM Destination: SARAM	n+2 2n‡	n+2 2n‡	n+2 2n‡ n+4† 2n+2§	$n+2+2p_{code}$ $2n+2p_{code}$‡
Source: External Destination: SARAM	$n+2+np_{src}$†	$n+2+np_{src}$	$n+2+np_{src}$ $n+4+np_{src}$†	$n+2+np_{src}+2p_{code}$
Source: DARAM/ROM Destination: External	$2n+2+nd_{dst}$	$2n+2+nd_{dst}$	$2n+2+nd_{dst}$	$2n+2+nd_{dst}+2p_{code}$
Source: SARAM Destination: External	$2n+2+nd_{dst}$	$2n+2+nd_{dst}$	$2n+2+nd_{dst}$	$2n+2+nd_{dst}+2p_{code}$
Source: External Destination: External	$4n+np_{src}+nd_{dst}$‡	$4n+np_{src}+nd_{dst}$	$4n+np_{src}+nd_{dst}$	$4n+2+np_{src}+nd_{dst}+2p_{code}$

† If the destination operand and the code are in the same SARAM block
‡ If both the source and the destination operands are in the same SARAM block
§ If both operands and the code are in the same SARAM block

Example 1

```
BLPD    #800h,00h ;(DP=6)
```

	Before Instruction	After Instruction
Program Memory 800h	0Fh	0Fh
Data Memory 300h	0h	0Fh

Example 2

```
BLPD    #800h,*,AR7
```

	Before Instruction	After Instruction
ARP	0	7
AR0	310h	310h
Program Memory 800h	1111h	1111h
Data Memory 310h	0100h	1111h

Syntax **CALA**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	1	1	0	0	0	0

Execution

PC + 1 → TOS
ACC(15:0) → PC

Status Bits None

Description The current program counter (PC) is incremented and pushed onto the top of the stack (TOS). Then, the contents of the lower half of the accumulator are loaded into the PC. Execution continues at this address.

The CALA instruction is used to perform computed subroutine calls.

Words 1

Cycles

Cycles for a Single CALA Instruction

ROM	DARAM	SARAM	External
4	4	4	4+3p

Note: When this instruction reaches the execute phase of the pipeline, two additional instruction words have entered the pipeline. When the PC discontinuity is taken, these two instruction words are discarded.

Example CALA

	Before Instruction		After Instruction
PC	25h	PC	83h
ACC	83h	ACC	83h
TOS	100h	TOS	26h

Syntax **CALL** *pma* [, *ind* [, **AR***n*]] Indirect addressing

Operands

pma: 16-bit program-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

CALL *pma* [*, ind* [*,* **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	1	0	1	ARU			N	NAR		
pma															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

PC + 2 → TOS
pma → PC
Modify (current AR) and (ARP) as specified.

Status Bits None

Description The current program counter (PC) is incremented and pushed onto the top of the stack (TOS). Then, the contents of the pma, either a symbolic or numeric address, are loaded into the PC. Execution continues at this address. The current auxiliary register and ARP contents are modified as specified.

Words 2

Cycles

Cycles for a Single CALL Instruction

ROM	DARAM	SARAM	External
4	4	4	4+4p†

Note: When this instruction reaches the execute phase of the pipeline, two additional instruction words have entered the pipeline. When the PC discontinuity is taken, these two instruction words are discarded.

Example

```
CALL    191,*+,AR0
```

	Before Instruction		After Instruction
ARP	1	ARP	0
AR1	05h	AR1	06h
PC	30h	PC	0BFh
TOS	100h	TOS	32h

Program address 0BFh (191) is loaded into the program counter, and the program continues executing from that location.

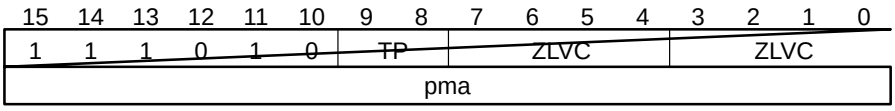
Syntax **CC** *pma*, *cond* 1 [,*cond* 2] [,...]

Operands

pma: 16-bit program-memory address

cond	*Condition*
EQ	ACC = 0
NEQ	ACC ≠ 0
LT	ACC < 0
LEQ	ACC ≤ 0
GT	ACC > 0
GEQ	ACC ≥ 0
NC	C = 0
C	C = 1
NOV	OV = 0
OV	OV = 1
BIO	BIO low
NTC	TC = 0
TC	TC = 1
UNC	Unconditionally

Opcode



Note: The TP and ZLVC fields are defined on pages 8-3 and 8-4.

Execution

If *cond* 1 AND *cond* 2 AND ...
Then
PC + 2 → TOS
pma → PC
Else
Increment PC

Status Bits None

Description Control is passed to the specified program-memory address (pma) if the specified conditions are met. Not all combinations of conditions are meaningful. For example, testing for LT and GT is contradictory. In addition, testing BIO is mutually exclusive to testing TC. The CC instruction operates like the CALL instruction if all conditions are true.

Words 2

Cycles

Cycles for a Single CC Instruction

Condition	ROM	DARAM	SARAM	External
True	4	4	4	4+4p†
False	2	2	2	2+2p

† The processor performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example

```
CC          PGM191,LEQ,C
```

If the accumulator contents are less than or equal to 0 and the carry bit is set, 0BFh (191) is loaded into the program counter, and the program continues to execute from that location. If the conditions are not met, execution continues at the instruction following the CC instruction.

Syntax **CLRC** *control bit*

Operands

control bit: Select one of the following control bits:

C	Carry bit of status register ST1
CNF	RAM configuration control bit of status register ST1
INTM	Interrupt mode bit of status register ST0
OVM	Overflow mode bit of status register ST0
SXM	Sign-extension mode bit of status register ST1
TC	Test/control flag bit of status register ST1
XF	XF pin status bit of status register ST1

Opcode

CLRC C

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	1	1	1	0

CLRC CNF

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	1	0	0

CLRC INTM

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	0	0	0

CLRC OVM

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	0	1	0

CLRC SXM

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	1	1	0

CLRC TC

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	1	0	1	0

CLRC XF

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	1	1	0	0

Execution

Increment PC, then ...
0 → control bit

Status Bits None

Description The specified control bit is cleared to 0. Note that the LST instruction can also be used to load ST0 and ST1. See subsection 3.5, *Status Registers ST0 and ST1*, on page 3-15 for more information about these control bits.

Words 1

Cycles

Cycles for a Single CLRC Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of a CLRC Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
CLRC TC   ;(TC is bit 11 of ST1)
```

	Before Instruction		After Instruction
ST1	x9xxh	ST1	x1xxh

Syntax **CMPL**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	0	0	0	1

Execution Increment PC, then ...
$\overline{(ACC)} \rightarrow ACC$

Status Bits None

Description The contents of the accumulator are replaced with its logical inversion (1s complement). The carry bit is unaffected.

Words 1

Cycles

Cycles for a Single CMPL Instruction

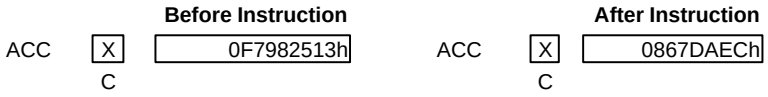
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an CMPL Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
CMPL
```



Syntax **CMPR** *CM*

Operands CM: Value from 0 to 3

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	0	1	0	0	0	1	CM	

Execution Increment PC, then ...
Compare (current AR) to (AR0) and place the result in the TC bit of status register ST1.

Status Bits *Affects*
TC

This instruction is not affected by SXM. It does not affect SXM.

Description The CMPR instruction performs a comparison specified by the value of CM:

If CM = 00, test whether (current AR) = (AR0)
If CM = 01, test whether (current AR) < (AR0)
If CM = 10, test whether (current AR) > (AR0)
If CM = 11, test whether (current AR) ≠ (AR0)

If the condition is true, the TC bit is set to 1. If the condition is false, the TC bit is cleared to 0.

Note that the auxiliary register values are treated as unsigned integers in the comparisons.

Words 1

Cycles

Cycles for a Single CMPR Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an CMPR Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
CMPR    2       ;(current AR) > (AR0)?
```

	Before Instruction		After Instruction
ARP	4	ARP	4
AR0	0FFFFh	AR0	0FFFFh
AR4	7FFFh	AR4	7FFFh
TC	1	TC	0

Syntax

DMOV *dma* Direct addressing
DMOV *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

DMOV *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	1	1	0	dma						

DMOV *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → data-memory address + 1

Status Bits

Affected by
CNF

Description

The contents of the specified data-memory address are copied into the contents of the next higher address. When data is copied from the addressed location to the next higher location, the contents of the addressed location remain unaltered.

DMOV works only within on-chip data RAM blocks. It works within any configurable RAM block if that block is configured as data memory. In addition, the data move function is continuous across block boundaries. The data move function cannot be performed on external data memory. If the instruction specifies an external memory address, DMOV reads the specified memory location but performs *no* operations.

The data move function is useful in implementing the z^{-1} delay encountered in digital signal processing. The DMOV function is a subtask of the LTD and MACD instructions (see the LTD and MACD instructions for more information).

Words

1

Cycles

Cycles for a Single DMOV Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 3†	1+p
External‡	2+2d	2+2d	2+2d	5+2d+p

† If the operand and the code are in the same SARAM block
‡ If used on external memory, DMOV reads the specified memory location but performs no operations.

Cycles for a Repeat (RPT) Execution of a DMOV Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	2n–2	2n–2	2n–2, 2n+1†	2n–2+p
External‡	4n–2+2nd	4n–2+2nd	4n–2+2nd	4n+1+2nd+p

† If the operand and the code are in the same SARAM block
‡ If used on external memory, DMOV reads the specified memory location but performs no operations.

Example 1

```
DMOV    DAT8    ;(DP = 6)
```

	Before Instruction		After Instruction
Data Memory 308h	43h	Data Memory 308h	43h
Data Memory 309h	2h	Data Memory 309h	43h

Example 2

```
DMOV    *,AR1
```

	Before Instruction		After Instruction
ARP	0	ARP	1
AR0	30Ah	AR0	30Ah
Data Memory 30Ah	40h	Data Memory 30Ah	40h
Data Memory 30Bh	41h	Data Memory 30Bh	40h

Syntax **IDLE**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	1	0	0	0	1	0

Execution Increment PC, then wait for unmasked or nonmaskable hardware interrupt.

Status Bits *Affected by*
INTM

Description The IDLE instruction forces the program being executed to halt until the CPU receives a request from an unmasked hardware interrupt (external or internal), $\overline{\text{NMI}}$, or reset. Execution of the IDLE instruction causes the 'C24x to enter a power-down mode. The PC is incremented once before the 'C24x enters power down; it is not incremented during the idle state. On-chip peripherals remain active; thus, their interrupts are among those that can wake the processor.

The idle state is exited by an unmasked interrupt even if INTM is 1. (INTM, the interrupt mode bit of status register ST0, normally disables maskable interrupts when it is set to 1.) When the idle state is exited by an unmasked interrupt, the CPU's next action, however, depends on INTM:

- If INTM is 0, the program branches to the corresponding interrupt service routine.
- If INTM is 1, the program continues executing at the instruction following the IDLE.

$\overline{\text{NMI}}$ and reset are not maskable; therefore, if the idle state is exited by $\overline{\text{NMI}}$ or reset, the corresponding interrupt service routine is executed, regardless of INTM.

Words 1

Cycles

Cycles for a Single IDLE Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Example

```
IDLE    ;The processor idles until a hardware reset,
        ;a hardware NMI, or an unmasked interrupt
        ;occurs.
```

Syntax

IN *dma*, *PA* Direct addressing
IN *ind*, *PA* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
PA: 16-bit I/O port or I/O-mapped register address
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

IN *dma , PA*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	1	1	0	dma						
PA															

IN *ind ,PA* [,**AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	1	1	1	ARU			N	NAR		
PA															

Note: ARU, N, and NAR are defined in Section 7.3, *Indirect Addressing Mode* (page 7–12).

Execution

Increment PC, then ...
PA → address bus lines A15–A0
Data bus lines D15–D0 → data-memory address
(PA) → data-memory address

Status Bits

None

Description

The IN instruction reads a 16-bit value from an I/O location into the specified data-memory location. The $\overline{IS}$ line goes low to indicate an I/O access. The $\overline{STRB}$, $\overline{RD}$, and READY timings are the same as for an external data-memory read.

The repeat (RPT) instruction can be used with the IN instruction to read in consecutive words from I/O space to data space.

Words

2

Cycles

Cycles for a Single IN Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Destination: DARAM	$2+io_{src}$	$2+io_{src}$	$2+io_{src}$	$3+io_{src}+2p_{code}$
Destination: SARAM	$2+io_{src}$	$2+io_{src}$	$2+io_{src}$ $3+io_{src}$ †	$3+io_{src}+2p_{code}$
Destination: External	$3+d_{dst}+io_{src}$	$3+d_{dst}+io_{src}$	$3+d_{dst}+io_{src}$	$6+d_{dst}+io_{src}+2p_{code}$

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an IN Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Destination: DARAM	$2n+nio_{src}$	$2n+nio_{src}$	$2n+nio_{src}$	$2n+1+nio_{src}+2p_{code}$
Destination: SARAM	$2n+nio_{src}$	$2n+nio_{src}$	$2n+nio_{src}$ $2n+2+nio_{src}$ †	$2n+1+nio_{src}+2p_{code}$
Destination: External	$4n-1+nd_{dst}+nio_{src}$	$4n-1+nd_{dst}+nio_{src}$	$4n-1+nd_{dst}+nio_{src}$	$4n+2+nd_{dst}+nio_{src}+2p_{code}$

† If the operand and the code are in the same SARAM block

Example 1

```
IN    7,1000h      ;Read in word from peripheral on
                   ;port address 1000h. Store word in
                   ;data memory location 307h (DP=6).
```

Example 2

```
IN    *,5h         ;Read in word from peripheral on
                   ;port address 5h. Store word in
                   ;data memory location specified by
                   ;current auxiliary register.
```

Syntax **INTR** *K*

Operands K: Value from 0 to 31 that indicates the interrupt vector location to branch to

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	1	K				

Execution (PC) + 1 → stack
corresponding interrupt vector location → PC

Status Bits *Affects*
INTM

This instruction is not affected by INTM.

Description The processor has locations for 32 interrupt vectors; each location is represented by a value K from 0 to 31. The INTR instruction is a software interrupt that transfers program control to the program-memory address specified by K. The vector at that address then leads to the corresponding interrupt service routine. Thus, the instruction allows any one of the interrupt service routines to be executed from your software. For a list of interrupts and their corresponding K values, see Table 6–2 on page 6-10. During execution of the instruction, the value PC + 1 (the return address) is pushed onto the stack. Neither the INTM bit nor the interrupt masks affect the INTR instruction. An INTR for the external interrupts looks exactly like an external interrupt (an interrupt acknowledge is generated, and maskable interrupts are globally disabled by setting INTM = 1).

Words 1

Cycles

Cycles for a Single INTR Instruction

ROM	DARAM	SARAM	External
4	4	4	4+3p†

† The processor performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example

```
INTR    3       ;PC + 1 is pushed onto the stack.
                ;Then control is passed to program
                ;memory location 6h.
```

Syntax

LACC *dma* [, *shift*]	Direct addressing
LACC *dma,* **16**	Direct with left shift of 16
LACC *ind* [, *shift* [, **AR***n*]]	Indirect addressing
LACC *ind*, **16**[, **AR***n*]	Indirect with left shift of 16
LACC #*lk* [, *shift*]	Long immediate addressing

Operands

dma:	7 LSBs of the data-memory address
shift:	Left shift value from 0 to 15 (defaults to 0)
n:	Value from 0 to 7 designating the next auxiliary register
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LACC *dma* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	shift				0	dma						

LACC *dma*, **16**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	1	0	0	dma						

LACC *ind* [, *shift*[, **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	shift				1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

LACC *ind*, **16**[, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

LACC #*lk* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	1	0	0	0	shift			
lk															

Execution

Increment PC, then ...

Event	*Addressing mode*
(data-memory address) $\times 2^{shift} \rightarrow$ ACC	Direct or indirect
(data-memory address) $\times 2^{16} \rightarrow$ ACC	Direct or indirect (shift of 16)
lk $\times 2^{shift} \rightarrow$ ACC	Long immediate

Status Bits

Affected by
SXM

Description

The contents of the specified data-memory address or a 16-bit constant are left shifted and loaded into the accumulator. During shifting, low-order bits are zero filled. High-order bits are sign extended if SXM = 1 and zeroed if SXM = 0.

Words

Words	*Addressing mode*
1	Direct or indirect
2	Long immediate

Cycles

Cycles for a Single LACC Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LACC Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

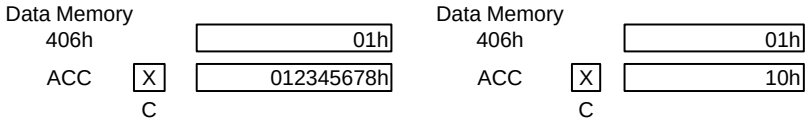
† If the operand and the code are in the same SARAM block

Cycles for a Single LACC Instruction (Using Immediate Addressing)

ROM	DARAM	SARAM	External
2	2	2	2+2p

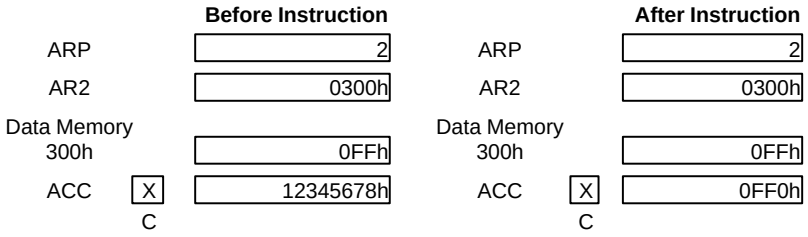
Example 1

```
LACC    6,4         ;(DP = 8: addresses 0400h–047Fh,
                    ;SXM = 0)
```



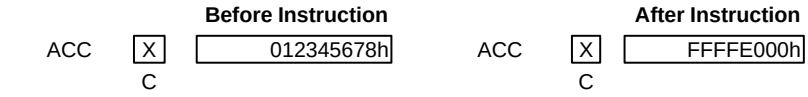
Example 2

```
LACC    *,4         ;(SXM = 0)
```



Example 3

```
LACC    #0F000h,1 ;(SXM = 1)
```



Syntax

LACL *dma* Direct addressing
LACL *ind* [, **AR***n*] Indirect addressing
LACL #*k* Short immediate

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
k: 8-bit short immediate value
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LACL *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	0	1	0	dma						

LACL *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

LACL #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	0	0	1	k							

Execution

Increment PC, then ...

Events	*Addressing mode*
0 → ACC(31:16) (data-memory address) → ACC(15:0)	Direct or indirect
0 → ACC(31:8) k → ACC(7:0)	Short immediate

Status Bits

This instruction is not affected by SXM.

Description

The contents of the addressed data-memory location or a zero-extended 8-bit constant are loaded into the 16 low-order bits of the accumulator. The upper half of the accumulator is zeroed. The data is treated as an unsigned 16-bit number rather than a 2s-complement number. There is no sign extension of the operand with this instruction, regardless of the state of SXM.

Words

1

Cycles

Cycles for a Single LACL Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LACL Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Cycles for a Single LACL Instruction (Using Immediate Addressing)

ROM	DARAM	SARAM	External
1	1	1	1+p

Example 1

```
LACL    1          ;(DP = 6: addresses 0300h–037Fh)
```

	Before Instruction		After Instruction	
Data Memory 301h		0h	Data Memory 301h	0h
ACC	X	7FFFFFFFh	ACC	X 0h
	C			C

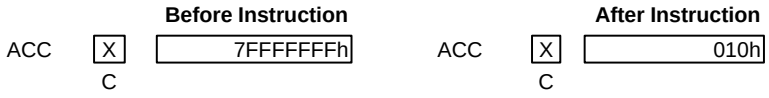
Example 2

```
LACL    *–,AR4
```

	Before Instruction		After Instruction	
ARP		0	ARP	4
AR0		401h	AR0	400h
Data Memory 401h		00FFh	Data Memory 401h	00FFh
ACC	X	7FFFFFFFh	ACC	X 0FFh
	C			C

Example 3

```
LACL    #10h
```



Syntax

LACT *dma* Direct addressing
LACT *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LACT *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	1	1	0	dma						

LACT *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) $\times$ $2^{(TREG(3:0))}$ $\rightarrow$ ACC

If SXM = 1:
Then (data-memory address) is sign extended.
If SXM = 0:
Then (data-memory address) is not sign extended.

Status Bits

Affected by
SXM

Description

The LACT instruction loads the accumulator with a data-memory value that has been left shifted. The left shift is specified by the four LSBs of the TREG, resulting in shift options from 0 to 15 bits. Using the four LSBs of the TREG as a shift code provides a dynamic shift mechanism. During shifting, the high-order bits are sign extended if SXM = 1 and zeroed if SXM = 0.

LACT may be used to denormalize a floating-point number if the actual exponent is placed in the four LSBs of the TREG register and the mantissa is referenced by the data-memory address. This method of denormalization can be used only when the magnitude of the exponent is four bits or less.

Words

1

Cycles

Cycles for a Single LACT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LACT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
LACT    1          ;(DP = 6: addresses 0300h–037Fh,
                   ;SXM = 0)
```

	Before Instruction	After Instruction
Data Memory 301h	1376h	1376h
TREG	14h	14h
ACC	98F7EC83h	13760h
X		
C		

Example 2

```
LACT    *–,AR3     ;(SXM = 1)
```

	Before Instruction	After Instruction
ARP	1	3
AR1	310h	30Fh
Data Memory 310h	0FF00h	0FF00h
TREG	11h	11h
ACC	098F7EC83h	0FFFFFE00h
X		
C		

Syntax

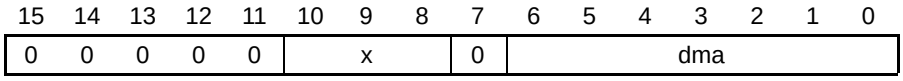
LAR AR*x*, *dma*	Direct addressing
LAR AR*x*, *ind* [, **AR***n*]	Indirect addressing
LAR AR*x*, *#k*	Short immediate addressing
LAR AR*x*, *#lk*	Long immediate addressing

Operands

x:	Value from 0 to 7 designating the auxiliary register to be loaded
dma:	7 LSBs of the data-memory address
k:	8-bit short immediate value
lk:	16-bit long immediate value
n:	Value from 0 to 7 designating the next auxiliary register
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

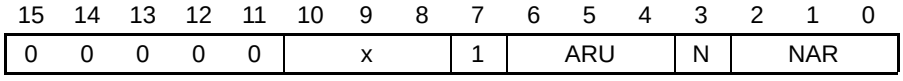
Opcode

LAR AR*x*, *dma*



15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	x			0	dma						

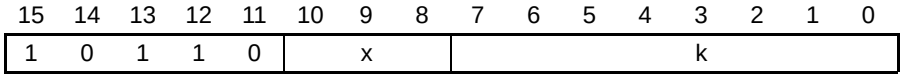
LAR AR*x*, *ind* [, **AR***n*]



15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	x			1	ARU			N	NAR		

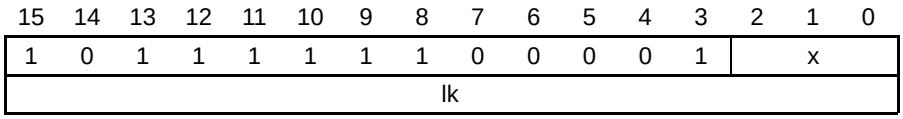
Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

LAR AR*x*, *#k*



15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	0	x			k							

LAR AR*x*, *#lk*



15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	0	0	0	0	1	x		
lk															

Execution

Increment PC, then ...

Event	*Addressing mode*
(data-memory address) → ARx	Direct or indirect
k → ARx	Short immediate
lk → ARx	Long immediate

Status Bits

None

Description

The contents of the specified data-memory address or an 8-bit or 16-bit constant are loaded into the specified auxiliary register (ARx). The specified constant is acted upon like an unsigned integer, regardless of the value of SXM.

The LAR and SAR (store auxiliary register) instructions can be used to load and store the auxiliary registers during subroutine calls and interrupts. If an auxiliary register is not being used for indirect addressing, LAR and SAR enable the register to be used as an additional storage register, especially for swapping values between data-memory locations without affecting the contents of the accumulator.

Words

Words	*Addressing mode*
1	Direct, indirect or short immediate
2	Long immediate

Cycles

Cycles for a Single LAR Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2	2	2	2+p_{code}
SARAM	2	2	2, 3†	2+p_{code}
External	2+d_{src}	2+d_{src}	2+d_{src}	3+d_{src}+p_{code}

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LAR Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2n	2n	2n	2n+p_{code}
SARAM	2n	2n	2n, 2n+1†	2n+p_{code}
External	2n+nd_{src}	2n+nd_{src}	2n+nd_{src}	2n+1+n$d_{src}p_{code}$

† If the operand and the code are in the same SARAM block

Cycles for a Single LAR Instruction (Using Short Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+p_{code}

Cycles for a Single LAR Instruction (Using Long Immediate Addressing)

ROM	DARAM	SARAM	External
2	2	2	2+2p

Example 1

```
LAR     AR0,16    ;(DP = 6: addresses 0300h–037Fh)
```

	Before Instruction		After Instruction
Data Memory 310h	18h	Data Memory 310h	18h
AR0	6h	AR0	18h

Example 2

```
LAR     AR4,*–
```

	Before Instruction		After Instruction
ARP	4	ARP	4
Data Memory 300h	32h	Data Memory 300h	32h
AR4	300h	AR4	32h

Note:

LAR in the indirect addressing mode ignores any AR modifications if the AR specified by the instruction is the same as that pointed to by the ARP. Therefore, in Example 2, AR4 is not decremented after the LAR instruction.

Example 3

```
LAR     AR4,#01h
```

	Before Instruction		After Instruction
AR4	0FF09h	AR4	01h

Example 4

```
LAR     AR6,#3FFFh
```

	Before Instruction		After Instruction
AR6	0h	AR6	3FFFh

Syntax

LDP *dma* Direct addressing
LDP *ind* [, **AR***n*] Indirect addressing
LDP #*k* Short immediate addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
k: 9-bit short immediate value
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LDP *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	0	1	0	dma						

LDP *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

LDP #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	0	k								

Execution

Increment PC, then ...

Event	*Addressing mode*
Nine LSBs of (data-memory address) → DP	Direct or indirect
k → DP	Short immediate

Status Bits

Affects
DP

Description

The nine LSBs of the contents of the addressed data-memory location or a 9-bit immediate value is loaded into the data page pointer (DP) of status register ST0. The DP can also be loaded by the LST instruction.

In direct addressing, the 9-bit DP and the 7-bit value specified in the instruction (dma) are concatenated to form the 16-bit data-memory address accessed by the instruction. The DP provides the 9 MSBs, and dma provides the 7 LSBs.

Words

1

Cycles

Cycles for a Single LDP Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2	2	2	2+p_{code}
SARAM	2	2	2, 3†	2+p_{code}
External	2+d_{src}	2+d_{src}	2+d_{src}	3+d_{src}+p_{code}

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LDP Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2n	2n	2n	2n+p_{code}
SARAM	2n	2n	2n, 2n+1†	2n+p_{code}
External	2n+nd_{src}	2n+nd_{src}	2n+nd_{src}	2n+1+$nd_{src}p_{code}$

† If the operand and the code are in the same SARAM block

Cycles for a Single LDP Instruction (Using Short Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+p_{code}

Example 1

```
LDP     127     ;(DP = 511: addresses FF80h–FFFFh)
```

	Before Instruction		After Instruction
Data Memory FFFFh	FEDCh	Data Memory FFFFh	FEDCh
DP	1FFh	DP	0DCh

Example 2

```
LDP     #0h
```

	Before Instruction		After Instruction
DP	1FFh	DP	0h

Example 3

```
LDP     *,AR5
```

	Before Instruction		After Instruction
ARP	4	ARP	5
AR4	300h	AR4	300h
Data Memory 300h	06h	Data Memory 300h	06h
DP	1FFh	DP	06h

Syntax

LPH *dma* Direct addressing
LPH *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LPH *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	0	1	0	dma						

LPH *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → PREG (31:16)

Status Bits

None

Description

The 16 high-order bits of the PREG are loaded with the content of the specified data-memory address. The low-order PREG bits are unaffected.

The LPH instruction can be used for restoring the high-order bits of the PREG after interrupts and subroutine calls.

Words

1

Cycles

Cycles for a Single LPH Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LPH Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
LPH     DAT0    ;(DP = 4)
```

	Before Instruction		After Instruction
Data Memory 200h	0F79Ch	Data Memory 200h	0F79Ch
PREG	30079844h	PREG	0F79C9844h

Example 2

```
LPH     *,AR6
```

	Before Instruction		After Instruction
ARP	5	ARP	6
AR5	200h	AR5	200h
Data Memory 200h	0F79Ch	Data Memory 200h	0F79Ch
PREG	30079844h	PREG	0F79C9844h

Syntax

LST #*m*, *dma* Direct addressing
LST #*m*, *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
m: Select one of the following:
0 Indicates that ST0 will be loaded
1 Indicates that ST1 will be loaded
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LST #0, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	1	0	0	dma						

LST #0, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

LST #1, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	1	1	0	dma						

LST #1, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → status register STm

For details about the differences between an LST #0 operation and an LST #1 operation, see Figure 8–3, Figure 8–4, and the description paragraph.

Figure 8–3. LST #0 Operation

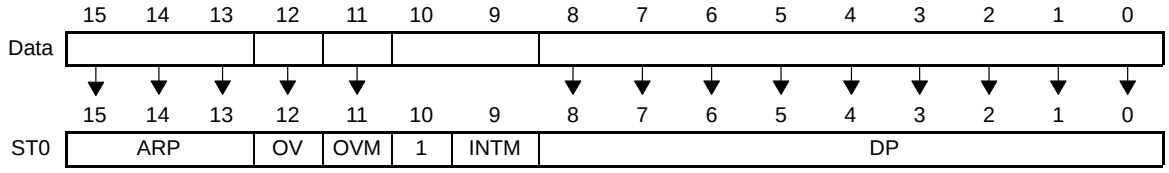
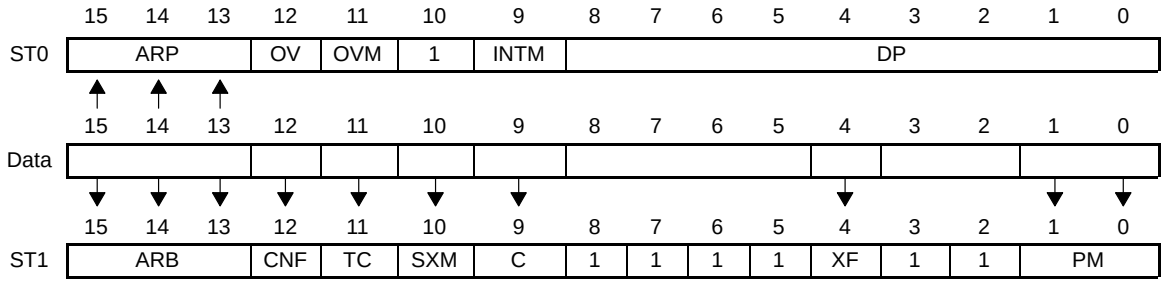


Figure 8–4. LST #1 Operation



Status Bits

Affects
ARB, ARP, OV, OVM, DP, CNF, TC, SXM, C, XF, and PM

This instruction does not affect INTM.

Description

The specified status register (ST0 or ST1) is loaded with the addressed data-memory value. Note the following points:

- The LST #0 operation does not affect the ARB field in the ST1 register, even though a new ARP is loaded.
- During the LST #1 operation, the value loaded into ARB is also loaded into ARP.
- If a next AR value is specified as an operand in the indirect addressing mode, this operand is ignored. ARP is loaded with the three MSBs of the value contained in the addressed data-memory location.
- Reserved bit values in the status registers are always read as 1s. Writes to these bits have no effect.

The LST instruction can be used for restoring the status registers after subroutine calls and interrupts.

Words

1

Cycles

Cycles for a Single LST Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2	2	2	$2+p_{code}$
SARAM	2	2	2, 3†	$2+p_{code}$
External	$2+d_{src}$	$2+d_{src}$	$2+d_{src}$	$3+d_{src}+p_{code}$

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LST Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2n	2n	2n	$2n+p_{code}$
SARAM	2n	2n	2n, 2n+1†	$2n+p_{code}$
External	$2n+nd_{src}$	$2n+nd_{src}$	$2n+nd_{src}$	$2n+1+nd_{src}+p_{code}$

† If the operand and the code are in the same SARAM block

Example 1

```
MAR   *,AR0
LST   #0,*,AR1  ;The data memory word addressed by the
                ;contents of auxiliary register AR0 is
                ;loaded into status register ST0,except
                ;for the INTM bit. Note that even
                ;though a next ARP value is specified,
                ;that value is ignored. Also note that
                ;the old ARP is not loaded into the
                ;ARB.
```

Example 2

```
LST       #0,60h    ;(DP = 0)
```

	Before Instruction		**After Instruction**
Data Memory 60h	2404h	Data Memory 60h	2404h
ST0	6E00h	ST0	2604h
ST1	05ECh	ST1	05ECh

Example 3

```
LST     #0,*-,AR1
```

	Before Instruction		After Instruction
ARP	4	ARP	7
AR4	3FFh	AR4	3FEh
Data Memory 3FFh	EE04h	Data Memory 3FFh	EE04h
ST0	EE00h	ST0	EE04h
ST1	F7ECh	ST1	F7ECh

Example 4

```
LST     #1,00h    ;(DP = 6)
                  ;Note that the ARB is loaded with
                  ;the new ARP value.
```

	Before Instruction		After Instruction
Data Memory 300h	E1BCh	Data Memory 300h	E1BCh
ST0	0406h	ST0	E406h
ST1	09ECh	ST1	E1FCh

Syntax

LT *dma* Direct addressing
LT *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LT *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	1	1	0	dma						

LT *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → TREG

Status Bits

None

Description

TREG is loaded with the contents of the specified data-memory address. The LT instruction may be used to load TREG in preparation for multiplication. See also the LTA, LTD, LTP, LTS, MPY, MPYA, MPYS, and MPYU instructions.

Words

1

Cycles

Cycles for a Single LT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
LT    24          ;(DP = 8: addresses 0400h–047Fh)
```

	Before Instruction	After Instruction
Data Memory 418h	62h	62h
TREG	3h	62h

Example 2

```
LT    *,AR3
```

	Before Instruction	After Instruction
ARP	2	3
AR2	418h	418h
Data Memory 418h	62h	62h
TREG	3h	62h

Syntax

LTA *dma* Direct addressing
LTA *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LTA *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	0	0	0	dma						

LTA *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → TREG
(ACC) + shifted (PREG) → ACC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

TREG is loaded with the contents of the specified data-memory address. The contents of the product register, shifted as defined by the PM status bits, are added to the accumulator, and the result is placed in the accumulator.

The carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry.

The function of the LTA instruction is a subtask of the LTD instruction.

Words

1

Cycles

Cycles for a Single LTA Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

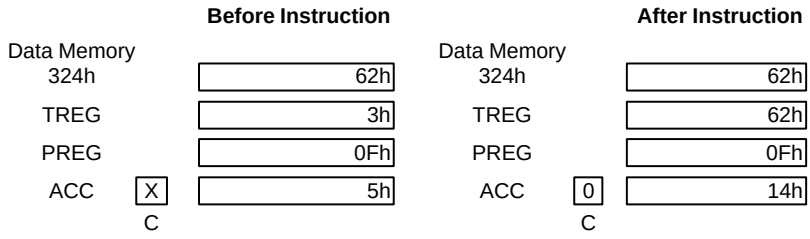
Cycles for a Repeat (RPT) Execution of an LTA Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

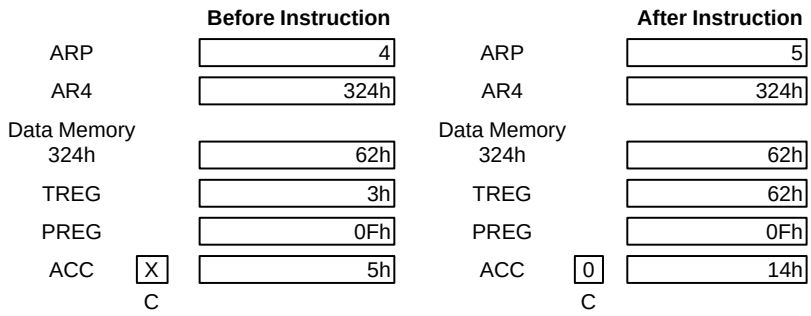
Example 1

```
LTA     36      ;(DP = 6: addresses 0300h–037Fh,
                ;PM =0: no shift of product)
```



Example 2

```
LTA     *,AR5   ;(PM = 0)
```



Syntax

LTD *dma* Direct addressing
LTD *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LTD *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	1	0	0	dma						

LTD *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → TREG
(data-memory address) → data-memory address + 1
(ACC) + shifted (PREG) → ACC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

TREG is loaded with the contents of the specified data-memory address. The contents of the PREG, shifted as defined by the PM status bits, are added to the accumulator, and the result is placed in the accumulator. The contents of the specified data-memory address are also copied to the next higher data-memory address.

This instruction is valid for all blocks of on-chip RAM configured as data memory. The data move function is continuous across the boundaries of contiguous blocks of memory but cannot be used with external data memory or memory-mapped registers. The data move function is described under the instruction DMOV.

Note:

If LTD is used with external data memory, its function is identical to that of LTA; that is, the previous product is accumulated, and the TREG is loaded from external data memory, but *the data move does not occur.*

The carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry.

Words 1

Cycles

Cycles for a Single LTD Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External‡**
DARAM	1	1	1	1+p
SARAM	1	1	1, 3†	1+p
External	2+2d	2+2d	2+2d	5+2d+p

† If the operand and the code are in the same SARAM block
‡ If the LTD instruction is used with external memory, the data move does not occur. (The previous product is accumulated, and the TREG is loaded.)

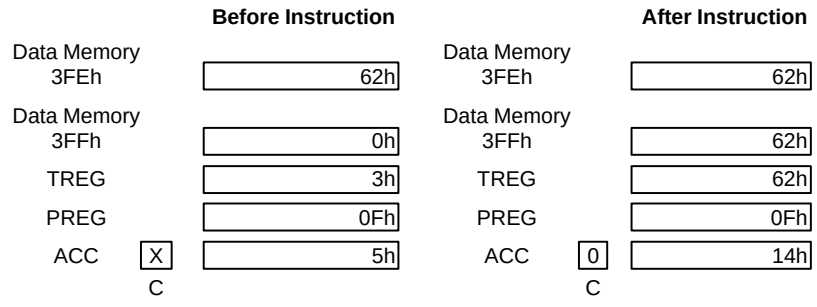
Cycles for a Repeat (RPT) Execution of an LTD Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External‡**
DARAM	n	n	n	n+p
SARAM	2n–2	2n–2	2n–2, 2n+1†	2n–2+p
External	4n–2+2nd	4n–2+2nd	4n–2+2nd	4n+1+2nd+p

† If the operand and the code are in the same SARAM block
‡ If the LTD instruction is used with external memory, the data move does not occur. (The previous product is accumulated, and the TREG is loaded.)

Example 1

```
LTD     126     ;(DP = 7: addresses 0380h–03FFh,
                ;PM = 0: no shift of product).
```



Example 2

```
LTD     *,AR3     ;(PM = 0)
```

	Before Instruction	After Instruction
ARP	1	3
AR1	3FEh	3FEh
Data Memory 3FEh	62h	62h
Data Memory 3FFh	0h	62h
TREG	3h	62h
PREG	0Fh	0Fh
ACC	5h	14h
C	X	0

Note: The data move function for LTD can occur only within on-chip data memory RAM blocks.

Syntax

LTP *dma* Direct addressing
LTP *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LTP *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	0	1	0	dma						

LTP *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → TREG
shifted (PREG) → ACC

Status Bits

Affected by
PM

Description

The TREG is loaded with the content of the addressed data-memory location, and the PREG value is stored in the accumulator. The shift at the output of the PREG is controlled by the PM status bits.

Words

1

Cycles

Cycles for a Single LTP Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

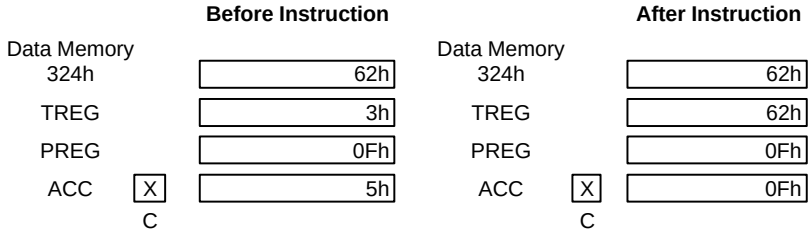
Cycles for a Repeat (RPT) Execution of an LTP Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

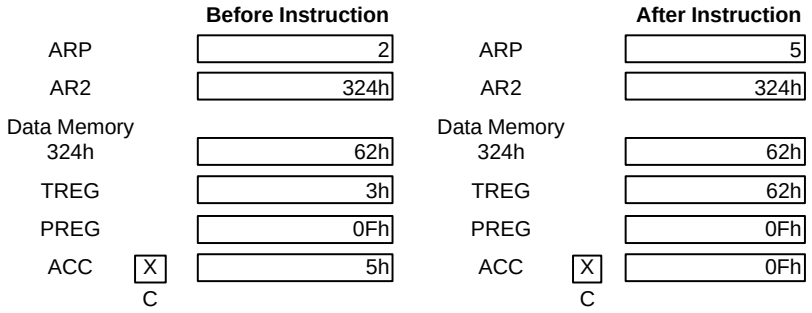
Example 1

```
LTP     36      ;(DP = 6: addresses 0300h–037Fh,
                ;PM = 0: no shift of product)
```



Example 2

```
LTP     *,AR5   ;(PM = 0)
```



Syntax

LTS *dma* Direct addressing
LTS *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

LTS *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	0	0	0	dma						

LTS *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → TREG
ACC – shifted (PREG) → ACC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

TREG is loaded with the contents of the addressed data-memory location. The contents of the product register, shifted as defined by the contents of the PM status bits, are subtracted from the accumulator. The result is placed in the accumulator.

The carry bit is cleared (C = 0) if the result of the subtraction generates a borrow and is set (C = 1) if it does not generate a borrow.

Words

1

Cycles for a Single LTS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an LTS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
LTS     DAT36     ;(DP = 6: addresses 0300h–037Fh,
                  ;PM = 0: no shift of product)
```

		Before Instruction		After Instruction
Data Memory 324h		62h	Data Memory 324h	62h
TREG		3h	TREG	62h
PREG		0Fh	PREG	0Fh
ACC	X	05h	ACC	0 0FFFFFFF6h
	C			C

Example 2

```
LTS     *,AR2     ;(PM = 0)
```

		Before Instruction		After Instruction
ARP		1	ARP	2
AR1		324h	AR1	324h
324h		62h	324h	62h
TREG		3h	TREG	62h
PREG		0Fh	PREG	0Fh
ACC	X	05h	ACC	0 0FFFFFFF6h
	C			C

Syntax

MAC *pma*, *dma* Direct addressing
MAC *pma*, *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
pma: 16-bit program-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MAC *pma*, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	0	1	0	0	dma						
pma															

MAC *pma*, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	0	1	0	1	ARU			N	NAR		
pma															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then . . .
(PC) → MSTACK
pma → PC
(ACC) + shifted (PREG) → ACC
(data-memory address) → TREG
(data-memory address) × (pma) → PREG
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC

While (repeat counter) ≠ 0:
(ACC) + shifted (PREG) → ACC
(data-memory address) → TREG
(data-memory address) × (pma) → PREG
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC
(repeat counter) – 1 → repeat counter

(MSTACK) → PC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

The MAC instruction:

- Adds the previous product, shifted as defined by the PM status bits, to the accumulator. The carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry.
- Loads the TREG with the content of the specified data-memory address.
- Multiplies the data-memory value in the TREG by the contents of the specified program-memory address.

The data and program memory locations on the 'C24x may be any non-reserved on-chip or off-chip memory locations. If the program memory is block B0 of on-chip RAM, the CNF bit must be set to 1.

When the MAC instruction is repeated, the program-memory address contained in the PC is incremented by 1 during each repetition. This makes it possible to access a series of operands in program memory. If you use indirect addressing to specify the data-memory address, a new data-memory address can be accessed during each repetition. If you use the direct addressing mode, the specified data-memory address is a constant; it is not modified during each repetition.

MAC is useful for long sum-of-products operations because, when repeated, it becomes a single-cycle instruction once the RPT pipeline is started.

Words

2

Cycles

Cycles for a Single MAC Instruction

Operand	ROM	DARAM	SARAM	External
Operand 1: DARAM/ ROM Operand 2: DARAM	3	3	3	$3+2p_{code}$
Operand 1: SARAM Operand 2: DARAM	3	3	3	$3+2p_{code}$
Operand 1: External Operand 2: DARAM	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}+2p_{code}$
Operand 1: DARAM/ ROM Operand 2: SARAM	3	3	3	$3+2p_{code}$
Operand 1: SARAM Operand 2: SARAM	3 4†	3 4†	3 4†	$3+2p_{code}$ $4+2p_{code}$†
Operand 1: External Operand 2: SARAM	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}+2p_{code}$
Operand 1: DARAM/ ROM Operand 2: External	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}+2p_{code}$
Operand 1: SARAM Operand 2: External	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}+2p_{code}$
Operand 1: External Operand 2: External	$4+p_{op1}+d_{op2}$	$4+p_{op1}+d_{op2}$	$4+p_{op1}+d_{op2}$	$4+p_{op1}+d_{op2}+2p_{code}$

† If both operands are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an MAC Instruction

Operand	ROM	DARAM	SARAM	External
Operand 1: DARAM/ ROM Operand 2: DARAM	n+2	n+2	n+2	$n+2+2p_{code}$
Operand 1: SARAM Operand 2: DARAM	n+2	n+2	n+2	$n+2+2p_{code}$
Operand 1: External Operand 2: DARAM	$n+2+np_{op1}$	$n+2+np_{op1}$	$n+2+np_{op1}$	$n+2+np_{op1}+2p_{code}$

† If both operands are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an MAC Instruction (Continued)

Operand	ROM	DARAM	SARAM	External
Operand 1: DARAM/ ROM Operand 2: SARAM	n+2	n+2	n+2	n+2+2p_{code}
Operand 1: SARAM Operand 2: SARAM	n+2 2n+2†	n+2 2n+2†	n+2 2n+2†	n+2+2p_{code} 2n+2†
Operand 1: External Operand 2: SARAM	n+2+np_{op1}	n+2+np_{op1}	n+2+np_{op1}	n+2+np_{op1}+2p_{code}
Operand 1: DARAM/ ROM Operand 2: External	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}+2p_{code}
Operand 1: SARAM Operand 2: External	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}+2p_{code}
Operand 1: External Operand 2: External	2n+2+np_{op1}+ nd_{op2}	2n+2+np_{op1}+nd_{op2}	2n+2+np_{op1}+nd_{op2}	2n+2+np_{op1}+nd_{op2}+ 2p_{code}

† If both operands are in the same SARAM block

Example 1

```
MAC     0FF00h,02h      ;(DP = 6, PM = 0, CNF = 1)
```

	Before Instruction	After Instruction
Data Memory 302h	23h	23h
Program Memory FF00h	4h	4h
TREG	45h	23h
PREG	458972h	08Ch
ACC	723EC41h	76975B3h
C	X	0

Example 2

```
MAC     0FF00h,*,AR5    ;(PM = 0, CNF = 1)
```

	Before Instruction	After Instruction
ARP	4	5
AR4	302h	302h
Data Memory 302h	23h	23h
Program Memory FF00h	4h	4h
TREG	45h	23h
PREG	458972h	8Ch
ACC	723EC41h	76975B3h
C	X	0

Syntax

MACD *pma*, *dma* Direct addressing
MACD *pma*, *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
pma: 16-bit program-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MACD *pma, dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	0	1	1	0	dma						
pma															

MACD *pma, ind* [*,* **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	0	1	1	1	ARU			N	NAR		
pma															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then . . .
(PC) → MSTACK
pma → PC
(ACC) + shifted (PREG) → ACC
(data-memory address) → TREG
(data-memory address) × (pma) → PREG
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC
(data-memory address) → data-memory address + 1

While (repeat counter) ≠ 0:
(ACC) + shifted (PREG) → ACC
(data-memory address) → TREG
(data-memory address) × (pma) → PREG
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC
(data-memory address) → data-memory address + 1
(repeat counter) – 1 → repeat counter

(MSTACK) → PC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

The MACD instruction:

- Adds the previous product, shifted as defined by the PM status bits, to the accumulator. The carry bit is set (C = 1) if the result of the addition generates a carry and is cleared (C = 0) if it does not generate a carry.
- Loads the TREG with the content of the specified data-memory address.
- Multiplies the data-memory value in the TREG by the contents of the specified program-memory address.
- Copies the contents of the specified data-memory address to the next higher data-memory address.

The data- and program-memory locations on the ’C24x may be any nonreserved, on-chip or off-chip memory locations. If the program memory is block B0 of on-chip RAM, the CNF bit must be set to 1. If MACD addresses one of the memory-mapped registers or external memory as a data-memory location, the effect of the instruction is that of a MAC instruction; the data move does not occur (see the DMOV instruction description on page 8-65).

When the MACD instruction is repeated, the program-memory address contained in the PC is incremented by 1 during each repetition. This makes it possible to access a series of operands in program memory. If you use indirect addressing to specify the data-memory address, a new data-memory address can be accessed during each repetition. If you use the direct addressing mode, the specified data-memory address is a constant; it will not be modified during each repetition.

MACD functions in the same manner as MAC, with the addition of a data move for on-chip RAM blocks. This feature makes MACD useful for applications such as convolution and transversal filtering. When used with RPT, MACD becomes a single-cycle instruction once the RPT pipeline is started.

Words

2

Cycles

Cycles for a Single MACD Instruction

Operand	ROM	DARAM	SARAM	External
Operand 1: DARAM/ ROM Operand 2: DARAM	3	3	3	$3+2p_{code}$
Operand 1: SARAM Operand 2: DARAM	3	3	3	$3+2p_{code}$
Operand 1: External Operand 2: DARAM	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}+2p_{code}$
Operand 1: DARAM/ ROM Operand 2: SARAM	3	3	3	$3+2p_{code}$
Operand 1: SARAM Operand 2: SARAM	3	3	3 4† 5‡	$3+2p_{code}$ $4+2p_{code}$†
Operand 1: External Operand 2: SARAM	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}$	$3+p_{op1}+2p_{code}$
Operand 1: DARAM/ ROM Operand 2: External§	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}+2p_{code}$
Operand 1: SARAM Operand 2: External§	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}$	$3+d_{op2}+2p_{code}$
Operand 1: External Operand 2: External§	$4+p_{op1}+d_{op2}$	$4+p_{op1}+d_{op2}$	$4+p_{op1}+d_{op2}$	$4+p_{op1}+d_{op2}+2p_{code}$

† If both operands are in the same SARAM block
‡ If both operands and code are in the same SARAM block
§ Data move operation is not performed when operand2 is in external data memory.

Cycles for a Repeat (RPT) Execution of an MACD Instruction

Operand	ROM	DARAM	SARAM	External
Operand 1: DARAM/ ROM Operand 2: DARAM	n+2	n+2	n+2	n+2+2p_{code}
Operand 1: SARAM Operand 2: DARAM	n+2	n+2	n+2	n+2+2p_{code}
Operand 1: External Operand 2: DARAM	n+2+np_{op1}	n+2+np_{op1}	n+2+np_{op1}	n+2+np_{op1}+2p_{code}
Operand 1: DARAM/ ROM Operand 2: SARAM	2n	2n	2n 2n+2†	2n+2p_{code}
Operand 1: SARAM Operand 2: SARAM	2n 3n‡	2n 3n‡	2n 2n+2† 3n‡ 3n+2§	2n+2p_{code} 3n‡
Operand 1: External Operand 2: SARAM	2n+np_{op1}	2n+np_{op1}	2n+np_{op1} 2n+2+np_{op1}†	2n+np_{op1}+2p_{code}
Operand 1: DARAM/ ROM Operand 2: External¶	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}+2p_{code}
Operand 1: SARAM Operand 2: External¶	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}	n+2+nd_{op2}+2p_{code}
Operand 1: External Operand 2: External¶	2n+2+np_{op1}+ nd_{op2}	2n+2+np_{op1}+nd_{op2}	2n+2+np_{op1}+nd_{op2}	2n+2+np_{op1}+nd_{op2}+ 2p_{code}

† If operand 2 and code are in the same SARAM block
‡ If both operands are in the same SARAM block
§ If both operands and code are in the same SARAM block
¶ Data move operation is not performed when operand2 is in external data memory.

Example 1

```
MACD  0FF00h,08h    ;(DP = 6: addresses 0300h–037Fh,
                    ;PM = 0: no shift of product,
                    ;CNF = 1: RAM B0 configured to
                    ;program memory).
```

	Before Instruction	After Instruction
Data Memory 308h	23h	23h
Data Memory 309h	18h	23h
Program Memory FF00h	4h	4h
TREG	45h	23h
PREG	458972h	8Ch
ACC	723EC41h	76975B3h
C	X	0

Example 2

```
MACD      0FF00h,*,AR6        ;(PM = 0, CNF = 1)
```

	Before Instruction	After Instruction
ARP	5	6
AR5	308h	308h
Data Memory 308h	23h	23h
Data Memory 309h	18h	23h
Program Memory FF00h	4h	4h
TREG	45h	23h
PREG	458972h	8Ch
ACC	723EC41h	76975B3h
C	X	0

Note: The data move function for MACD can occur only within on-chip data memory RAM blocks.

Syntax

MAR *dma* Direct addressing
MAR *ind* [, **AR***n*] Indirect addressing

Operands

n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MAR *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	1	1	0	dma						

MAR *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Event(s)	*Addressing mode*
Increment PC	Direct
Increment PC Modify (current AR) and (ARP) as specified	Indirect

Status Bits

Affects	*Addressing mode*
None	Direct
ARP and ARB	Indirect

Description

In the direct addressing mode, the MAR instruction acts as a NOP instruction.

In the indirect addressing mode, an auxiliary register value and the ARP value can be modified; however, the memory being referenced is not used. When MAR modifies the ARP value, the old ARP value is copied to the ARB field of ST1. Any operation that MAR performs with indirect addressing can also be performed with any instruction that supports indirect addressing. In addition, the ARP can also be loaded by an LST instruction.

The LARP instruction from the ’C25 instruction set is a subset of MAR. For example, MAR *, AR4 performs the same function as LARP 4, which loads the ARP with 4.

For loading an auxiliary register, see the description for the LAR instruction on page 8-79. For storing an auxiliary register value to data memory, see the SAR instruction on page 8-151.

Words 1

Cycles

Cycles for a Single MAR Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an MAR Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example 1

```
MAR     *,AR1     ;Load the ARP with 1.
```

	Before Instruction		After Instruction
ARP	0	ARP	1
ARB	7	ARB	0

Example 2

```
MAR     *+,AR5    ;Increment current auxiliary
                  ;register (AR1) and load ARP
                  ;with 5.
```

	Before Instruction		After Instruction
AR1	34h	AR1	35h
ARP	1	ARP	5
ARB	0	ARB	1

Syntax

MPY *dma*	Direct addressing
MPY *ind* [, **AR***n*]	Indirect addressing
MPY #*k*	Short immediate addressing

Operands

dma:	7 LSBs of the data-memory address
n:	Value from 0 to 7 designating the next auxiliary register
k:	13-bit short immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MPY *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	1	0	0	0	dma						

MPY *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	1	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

MPY #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	k												

Execution

Increment PC, then ...

Event	*Addressing mode*
(TREG) × (data-memory address) → PREG	Direct or indirect
(TREG) × k → PREG	Short immediate

Status Bits

None

Description

The contents of TREG are multiplied by the contents of the addressed data memory location. The result is placed in the product register (PREG). With short immediate addressing, TREG is multiplied by a signed 13-bit constant. The short-immediate value is right justified and sign extended before the multiplication, regardless of SXM.

Words

1

Cycles

Cycles for a Single MPY Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an MPY Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Cycles for a Single MPY Instruction (Using Short Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
1	1	1	1+p

Example 1

```
MPY     DAT13     ;(DP = 8)
```

	Before Instruction		**After Instruction**
Data Memory 40Dh	7h	Data Memory 40Dh	7h
TREG	6h	TREG	6h
PREG	36h	PREG	2Ah

Example 2 `MPY *,AR2`

	Before Instruction		After Instruction
ARP	1	ARP	2
AR1	40Dh	AR1	40Dh
Data Memory 40Dh	7h	Data Memory 40Dh	7h
TREG	6h	TREG	6h
PREG	36h	PREG	2Ah

Example 3 `MPY #031h`

	Before Instruction		After Instruction
TREG	2h	TREG	2h
PREG	36h	PREG	62h

Syntax

MPYA *dma* Direct addressing
MPYA *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MPYA *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	0	0	0	dma						

MPYA *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) + shifted (PREG) → ACC
(TREG) × (data-memory address) → PREG

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

The contents of TREG are multiplied by the contents of the addressed data memory location. The result is placed in the product register (PREG). The previous product, shifted as defined by the PM status bits, is also added to the accumulator.

Words

1

Cycles

Cycles for a Single MPYA Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an MPYA Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
MPYA    DAT13    ;(DP = 6, PM = 0)
```

	Before Instruction	After Instruction
Data Memory 30Dh	7h	7h
TREG	6h	6h
PREG	36h	2Ah
ACC	54h	8Ah
C	X	0

Example 2

```
MPYA    *,AR4    ;(PM = 0)
```

	Before Instruction	After Instruction
ARP	3	4
AR3	30Dh	30Dh
Data Memory 30Dh	7h	7h
TREG	6h	6h
PREG	36h	2Ah
ACC	54h	8Ah
C	X	0

Syntax

MPYS *dma* Direct addressing
MPYS *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MPYS *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	0	1	0	dma						

MPYS *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) – shifted (PREG) → ACC
(TREG) × (data-memory address) → PREG

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

Description

The contents of TREG are multiplied by the contents of the addressed data memory location. The result is placed in the product register (PREG). The previous product, shifted as defined by the PM status bits, is also subtracted from the accumulator, and the result is placed in the accumulator.

Words

1

Cycles

Cycles for a Single MPYS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an MPYS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
MPYS    DAT13    ;(DP = 6, PM = 0)
```

	Before Instruction		After Instruction	
Data Memory 30Dh		7h		7h
TREG		6h		6h
PREG		36h		2Ah
ACC	X	54h	1	1Eh
	C		C	

Example 2

```
MPYS    *,AR5    ;(PM = 0)
```

	Before Instruction		After Instruction	
ARP		4		5
AR4		30Dh		30Dh
Data Memory 30Dh		7h		7h
TREG		6h		6h
PREG		36h		2Ah
ACC	X	54h	1	1Eh
	C		C	

Syntax

MPYU *dma* Direct addressing
MPYU *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

MPYU *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	1	0	1	0	dma						

MPYU *ind* [,**AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	1	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
Unsigned (TREG) $\times$ unsigned (data-memory address) $\rightarrow$ PREG

Status Bits

None

This instruction is not affected by SXM.

Description

The unsigned contents of TREG are multiplied by the unsigned contents of the addressed data-memory location. The result is placed in the product register (PREG). The multiplier acts as a signed 17 $\times$ 17-bit multiplier for this instruction, with the MSB of both operands forced to 0.

When another instruction passes the resulting PREG value to data memory or to the CALU, the value passes first through the product shifter at the output of the PREG. This shifter always invokes sign extension on the PREG value when PM = 3 (right-shift-by-6 mode). Therefore, this shift mode should not be used if unsigned products are desired.

The MPYU instruction is particularly useful for computing multiple-precision products, such as when multiplying two 32-bit numbers to yield a 64-bit product.

Words

1

Cycles

Cycles for a Single MPYU Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an MPYU Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
MPYU    16      ;(DP = 4: addresses 0200h–027Fh)
```

	Before Instruction		After Instruction
Data Memory 210h	0FFFFh	Data Memory 210h	0FFFFh
TREG	0FFFFh	TREG	0FFFFh
PREG	1h	PREG	0FFFE0001h

Example 2

```
MPYU    *,AR6
```

	Before Instruction		After Instruction
ARP	5	ARP	6
AR5	210h	AR5	210h
Data Memory 210h	0FFFFh	Data Memory 210h	0FFFFh
TREG	0FFFFh	TREG	0FFFFh
PREG	1h	PREG	0FFFE0001h

Syntax **NEG**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	0	0	1	0

Execution Increment PC, then ...
(ACC) × –1 → ACC

Status Bits

Affected by	*Affects*
OVM	C and OV

Description The content of the accumulator is replaced with its arithmetic complement (2s complement). The OV bit is set when taking the NEG of 8000 0000h. If OVM = 1, the accumulator content is replaced with 7FFF FFFFh. If OVM = 0, the result is 8000 0000h. The carry bit (C) is cleared to 0 by this instruction for all nonzero values of the accumulator, and is set to 1 if the accumulator equals 0.

Words 1

Cycles

Cycles for a Single NEG Instruction

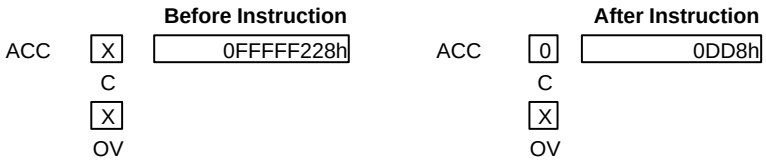
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an NEG Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

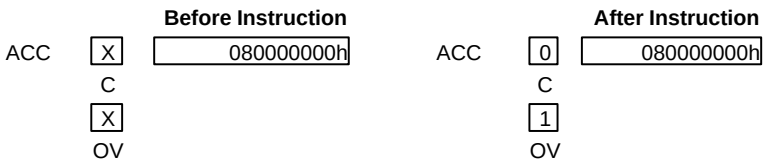
Example 1

```
NEG       ;(OVM = X) Convert –3544 to +3544
```



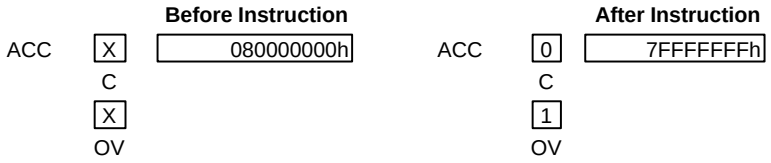
Example 2

```
NEG       ;(OVM = 0)
```



Example 3

```
NEG        ;(OVM = 1)
```



Syntax **NMI**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	1	0	0	1	0

Execution

(PC) + 1 → stack
24h → PC
1 → INTM

Status Bits

Affects
INTM

This instruction is not affected by INTM.

Description The NMI instruction forces the program counter to the nonmaskable interrupt vector located at 24h. This instruction has the same effect as the hardware nonmaskable interrupt $\overline{\text{NMI}}$.

Words 1

Cycles

Cycles for a Single NMI Instruction

ROM	DARAM	SARAM	External
4	4	4	4+3p†

† The 'C24x performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example

```
NMI   ;PC + 1 is pushed onto the stack, and then
      ;control is passed to program memory location
      ;24h.
```

Syntax **NOP**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	1	1	0	0	0	0	0	0	0	0

Execution Increment PC

Status Bits None

Description No operation is performed. The NOP instruction affects only the PC. The NOP instruction is useful to create pipeline and execution delays.

Words 1

Cycles

Cycles for a Single NOP Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an NOP Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
NOP       ;No operation is performed.
```

Syntax **NORM** *ind* Indirect addressing

Operands ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

NORM *ind*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	0	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...

If (ACC) = 0:
  Then TC $\rightarrow$ 1;
Else, if (ACC(31)) XOR (ACC(30)) = 0:
  Then TC $\rightarrow$ 0,
    (ACC) $\times$ 2 $\rightarrow$ ACC
    Modify (current AR) as specified;
  Else TC $\rightarrow$ 1.

Status Bits

Affects
TC

Description

The NORM instruction normalizes a signed number that is contained in the accumulator. Normalizing a fixed-point number separates it into a mantissa and an exponent. By finding the magnitude of the sign-extended number. An exclusive-OR operation is performed on accumulator bits 31 and 30 to determine if bit 30 is part of the magnitude or part of the sign extension. If they are the same, they are both sign bits, and the accumulator is left shifted to eliminate the extra sign bit.

The current AR is modified as specified to generate the magnitude of the exponent. It is assumed that the current AR is initialized before normalization begins. The default modification of the current AR is an increment.

Multiple executions of the NORM instruction may be required to completely normalize a 32-bit number in the accumulator. Although using NORM with RPT does not cause execution of NORM to fall out of the repeat loop automatically when the normalization is complete, no operation is performed for the remainder of the repeat loop. NORM functions on both positive and negative 2s-complement numbers.

Notes:

For the NORM instruction, the auxiliary register operations are executed during the fourth phase of the pipeline, the execution phase. For other instructions, the auxiliary register operations take place in the second phase of the pipeline, in the decode phase. Therefore:

1) **The auxiliary register values should not be modified by the two instruction words following NORM.** If the auxiliary register used in the NORM instruction is to be affected by either of the next two instruction words, the auxiliary register value is modified by the other instructions *before* it is modified by the NORM instruction.

2) **The value in the auxiliary register pointer (ARP) should not be modified by the two instruction words following NORM.** If either of the next two instruction words specify a change in the ARP value, the ARP value is changed *before* NORM is executed; the ARP does not point to the correct auxiliary register when NORM is executed.

Words 1

Cycles

Cycles for a Single NORM Instruction

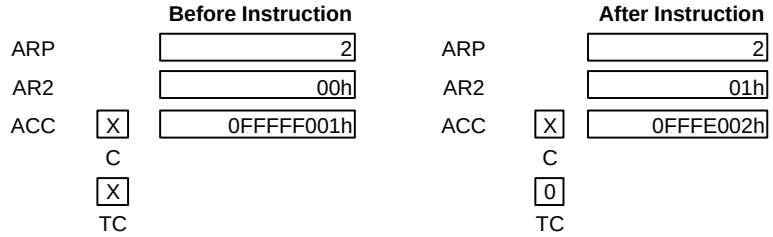
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of a NORM Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example 1

```
NORM    *+
```



Example 2 31-Bit Normalization:

```
        MAR    *,AR1      ;Use AR1 to store the exponent.
        LAR    AR1,#0h    ;Clear out exponent counter.
LOOP    NORM   *+         ;One bit is normalized.
        BCND   LOOP,NTC   ;If TC = 0, magnitude not found yet.
```

Example 3 15-Bit Normalization:

```
MAR    *,AR1      ;Use AR1 to store the exponent.
LAR    AR1,#0Fh   ;Initialize exponent counter.
RPT    #14        ;15-bit normalization specified (yielding
                  ;a 4-bit exponent and 16-bit mantissa).
NORM   *–         ;NORM automatically stops shifting when first
                  ;significant magnitude bit is found,
                  ;performing NOPs for the remainder of the
                  ;repeat loops.
```

The method used in Example 2 normalizes a 32-bit number and yields a 5-bit exponent magnitude. The method used in Example 3 normalizes a 16-bit number and yields a 4-bit magnitude. If the number requires only a small amount of normalization, the Example 2 method may be preferable to the Example 3 method because the loop in Example 2 runs only until normalization is complete. Example 3 always executes all 15 cycles of the repeat loop. Specifically, Example 2 is more efficient if the number requires three or fewer shifts. If the number requires six or more shifts, Example 3 is more efficient.

Syntax

OR *dma*	Direct addressing
OR *ind* [, **AR***n*]	Indirect addressing
OR #*lk* [, *shift*]	Long immediate addressing
OR #*lk,* **16**	Long immediate with left shift of 16

Operands

dma:	7 LSBs of the data-memory address
shift:	Left shift value from 0 to 15 (defaults to 0)
n:	Value from 0 to 7 designating the next auxiliary register
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

OR *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	0	1	0	dma						

OR *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

OR #*lk* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	1	1	0	0	shift			
lk															

OR #*lk* [*,* **16**]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	1	0	0	0	0	0	1	0
lk															

Execution

Increment PC, then ...

Event(s)	*Addressing mode*
(ACC(15:0)) OR (data-memory address) → ACC(15:0) (ACC(31:16)) → ACC(31:16)	Direct or indirect
(ACC) OR lk × 2^{shift} → ACC	Long immediate
(ACC) OR lk × 2^{16} → ACC	Long immediate with left shift of 16

Status Bits

None
This instruction is not affected by SXM.

Description

An OR operation is performed on the contents of the accumulator and the contents of the addressed data-memory location or a long-immediate value. The long-immediate value may be shifted before the OR operation. The result remains in the accumulator. All bit positions unoccupied by the data operand are zero filled, regardless of the value of the SXM status bit. Thus, the high word of the accumulator is unaffected by this instruction if direct or indirect addressing is used, or if immediate addressing is used with a shift of 0. Zeros are shifted into the least significant bits of the operand if immediate addressing is used with a nonzero shift count.

Words

Words	*Addressing mode*
1	Direct or indirect
2	Long immediate

Cycles

Cycles for a Single OR Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an OR Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

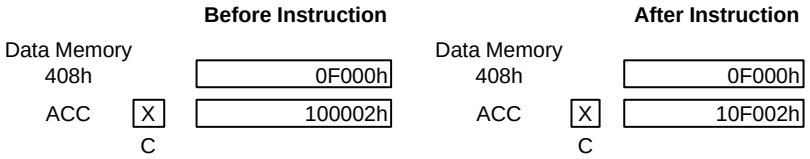
† If the operand and the code are in the same SARAM block

Cycles for a Single OR Instruction (Using Long Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+2p

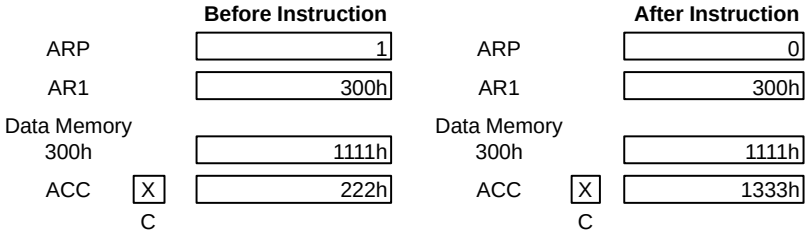
Example 1

```
OR DAT8    ;(DP = 8)
```



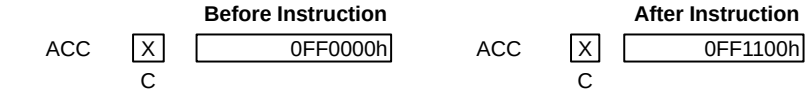
Example 2

```
OR *,AR0
```



Example 3

```
OR #08111h,8
```



Syntax

OUT *dma*, *PA* Direct addressing
OUT *ind*, *PA* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
PA: 16-bit I/O address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

OUT *dma*, *PA*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	0	0	0	dma						
PA															

OUT *ind*, *PA* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	1	0	0	1	ARU			N	NAR		
PA															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
PA → address bus A15–A0
(data-memory address) → data bus D15–D0
(data-memory address) → PA

Status Bits

None

Description

The OUT instruction writes a 16-bit value from a data-memory location to the specified I/O location. The $\overline{IS}$ line goes low to indicate an I/O access. The $\overline{STRB}$, R/$\overline{W}$, and READY timings are the same as for an external data-memory write.

RPT can be used with the OUT instruction to write consecutive words from data memory to I/O space.

Words

2

Cycles

Cycles for a Single OUT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM	$3+io_{dst}$	$3+io_{dst}$	$3+io_{dst}$	$5+io_{dst}+2p_{code}$
Source: SARAM	$3+io_{dst}$	$3+io_{dst}$	$3+io_{dst}$ $4+io_{dst}$†	$5+io_{dst}+2p_{code}$
Source: External	$3+d_{src}+io_{dst}$	$3+d_{src}+io_{dst}$	$3+d_{src}+io_{dst}$	$6+d_{src}+io_{dst}+2p_{code}$

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an OUT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Destination: DARAM	$3n+nio_{dst}$	$3n+nio_{dst}$	$3n+nio_{dst}$	$3n+3+nio_{dst}+2p_{code}$
Destination: SARAM	$3n+nio_{dst}$	$3n+nio_{dst}$	$3n+nio_{dst}$ $3n+1+nio_{dst}$†	$3n+3+nio_{dst}+2p_{code}$
Destination: External	$5n-2+nd_{src}+nio_{dst}$	$5n-2+nd_{src}+nio_{dst}$	$5n-2+nd_{src}+nio_{dst}$	$5n+1+nd_{src}+nio_{dst}+2p_{code}$

† If the operand and the code are in the same SARAM block

Example 1

```
OUT   DAT0,100h     ;(DP = 4) Write data word stored in
                    ;data memory location 200h to
                    ;peripheral at I/O port address
                    ;100h.
```

Example 2

```
OUT   *,100h        ;Write data word referenced by
                    ;current auxiliary register to
                    ;peripheral at I/O port address
                    ;100h.
```

Syntax **PAC**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	0	0	1	1

Execution Increment PC, then ...
shifted (PREG) → ACC

Status Bits *Affected by*
PM

Description The content of PREG, shifted as specified by the PM status bits, is loaded into the accumulator.

Words 1

Cycles

Cycles for a Single PAC Instruction

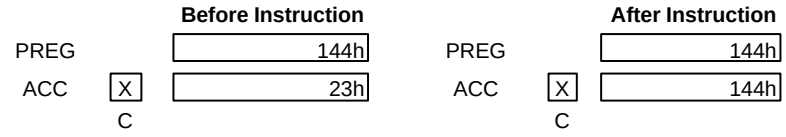
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of a PAC Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
PAC       ;(PM = 0: no shift of product)
```



Syntax **POP**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	1	1	0	0	1	0

Execution

Increment PC, then ...
(TOS) $\rightarrow$ ACC(15:0)
0 $\rightarrow$ ACC(31:16)
Pop stack one level

Status Bits None

Description

The content of the top of the stack (TOS) is copied to the low accumulator, and then the stack values move up one level. The upper half of the accumulator is set to all 0s.

The hardware stack functions as a last-in, first-out stack with eight locations. Any time a pop occurs, every stack value is copied to the next higher stack location, and the top value is removed from the stack. After a pop, the bottom two stack words have the same value. Because each stack value is copied, if more than seven stack pops (using the POP, POPD, RETC, or RET instructions) occur before any pushes occur, all levels of the stack contain the same value. No provision exists to check stack underflow.

Words 1

Cycles

Cycles for a Single POP Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of a POP Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example POP

		Before Instruction			After Instruction
ACC	X	82h	ACC	X	45h
	C			C	
Stack		45h	Stack		16h
		16h			7h
		7h			33h
		33h			42h
		42h			56h
		56h			37h
		37h			61h
		61h			61h

Syntax

POPD *dma* Direct addressing
POPD *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

POPD *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	1	0	0	dma						

POPD *ind* [,**AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(TOS) → data-memory address
Pop stack one level

Status Bits

None

Description

The value from the top of the stack is transferred into the data-memory location specified by the instruction. In the lower seven locations of the stack, the values are copied up one level. The stack operation is explained in the description for the POP instruction. No provision exists to check stack underflow.

Words

1

Cycles

Cycles for a Single POPD Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a POPD Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
POPD    DAT10    ;(DP = 8)
```

	Before Instruction		After Instruction
Data Memory 40Ah	55h	Data Memory 40Ah	92h
Stack	92h	Stack	72h
	72h		8h
	8h		44h
	44h		81h
	81h		75h
	75h		32h
	32h		0AAh
	0AAh		0AAh

Example 2

```
POPD    *+,AR1
```

	Before Instruction		After Instruction
ARP	0	ARP	1
AR0	300h	AR0	301h
Data Memory 300h	55h	Data Memory 300h	92h
Stack	92h	Stack	72h
	72h		8h
	8h		44h
	44h		81h
	81h		75h
	75h		32h
	32h		0AAh
	0AAh		0AAh

Syntax

PSHD *dma* Direct addressing
PSHD *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

PSHD *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	1	0	0	dma						

PSHD *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → TOS
Push all stack locations down one level

Status Bits

None

Description

The value from the data-memory location specified by the instruction is transferred to the top of the stack. In the lower seven locations of the stack, the values are also copied one level down, as explained in the description for the PUSH instruction. The value in the lowest stack location is lost.

Words

1

Cycles

Cycles for a Single PSHD Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a PSHD Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
PSHD    127     ;(DP = 3: addresses 0180–01FFh)
```

	Before Instruction		After Instruction
Data Memory 1FFh	65h	Data Memory 1FFh	65h
Stack	2h	Stack	65h
	33h		2h
	78h		33h
	99h		78h
	42h		99h
	50h		42h
	0h		50h
	0h		0h

Example 2

```
PSHD    *,AR1
```

	Before Instruction		After Instruction
ARP	0	ARP	1
AR0	1FFh	AR0	1FFh
Data Memory 1FFh	12h	Data Memory 1FFh	12h
Stack	2h	Stack	12h
	33h		2h
	78h		33h
	99h		78h
	42h		99h
	50h		42h
	0h		50h
	0h		0h

Syntax **PUSH**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	1	1	1	1	0	0

Execution

Increment PC, then...
Push all stack locations down one level
ACC(15:0) → TOS

Status Bits None

Description

The stack values move down one level. Then, the content of the lower half of the accumulator is copied onto the top of the hardware stack.

The hardware stack operates as a last-in, first-out stack with eight locations. If more than eight pushes (due to a CALA, CALL, CC, PSHD, PUSH, TRAP, INTR, or NMI instruction) occur before a pop, the first data values written are lost with each succeeding push.

Words 1

Cycles

Cycles for a Single PUSH Instruction

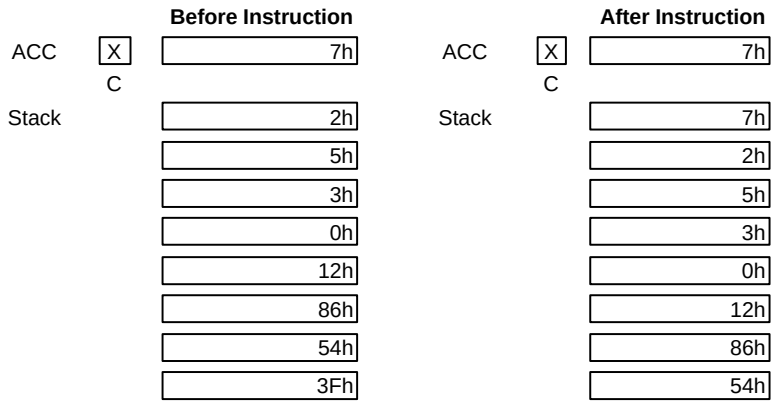
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of a PUSH Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

PUSH



Syntax **RET**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	1	1	1	1	0	0	0	0	0	0	0	0

Execution (TOS) → PC
Pop stack one level.

Status Bits None

Description The contents of the top stack register are copied into the program counter. The remaining stack values are then copied up one level. RET concludes subroutines and interrupt service routines to return program control to the calling or interrupted program sequence.

Words 1

Cycles

Cycles for a Single RET Instruction

ROM	DARAM	SARAM	External
4	4	4	4+3p

Note: When this instruction reaches the execute phase of the pipeline, two additional instruction words have entered the pipeline. When the PC discontinuity is taken, these two instruction words are discarded.

Example RET

	Before Instruction		After Instruction
PC	96h	PC	37h
Stack	37h	Stack	45h
	45h		75h
	75h		21h
	21h		3Fh
	3Fh		45h
	45h		6Eh
	6Eh		6Eh
	6Eh		6Eh

Syntax **RETC** *cond* 1 [, *cond* 2] [,...]

Operands

cond	*Condition*
EQ	ACC = 0
NEQ	ACC ≠ 0
LT	ACC < 0
LEQ	ACC ≤ 0
GT	ACC > 0
GEQ	ACC ≥ 0
NC	C = 0
C	C =1
NOV	OV = 0
OV	$\overline{\text{OV}}$ = 1
BIO	$\overline{\text{BIO}}$ low
NTC	TC = 0
TC	TC = 1
UNC	Unconditionally

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	1	1	TP		ZLVC				ZLVC			

Note: The TP and ZLVC fields are defined on pages 8-3 and 8-4.

Execution

If *cond* 1 AND *cond* 2 AND ...
    (TOS) → PC
    Pop stack one level
Else, continue

Status Bits None

Description If the specified condition or conditions are met, a standard return is executed (see the description for the RET instruction on page 8-141). Note that not all combinations of conditions are meaningful. For example, testing for LT and GT is contradictory. In addition, testing $\overline{\text{BIO}}$ is mutually exclusive to testing TC.

Words 1

Cycles

Cycles for a Single RETC Instruction

Condition	ROM	DARAM	SARAM	External
True	4	4	4	4+4p
False	2	2	2	2+2p

Note: The processor performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example

```
RETC      GEQ,NOV   ;A return is executed if the
                    ;accumulator content is positive
                    ;or zero and if the OV (overflow)
                    ;-bit is zero.
```

Syntax **ROL**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	1	1	0	0

Execution

Increment PC, then ...
C $\rightarrow$ ACC(0)
(ACC(31)) $\rightarrow$ C
(ACC(30:0)) $\rightarrow$ ACC(31:1)

Status Bits

Affects
C

This instruction is not affected by SXM.

Description The ROL instruction rotates the accumulator left one bit. The value of the carry bit is shifted into the LSB, then the MSB is shifted into the carry bit.

Words 1

Cycles

Cycles for a Single ROL Instruction

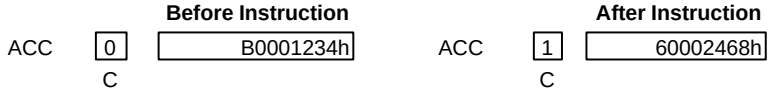
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an ROL Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
ROL
```



Syntax **ROR**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	1	1	0	1

Execution

Increment PC, then ...
C → ACC(31)
(ACC(0)) → C
(ACC(31:1)) → ACC(30:0)

Status Bits

Affects
C

This instruction is not affected by SXM.

Description The ROR instruction rotates the accumulator right one bit. The value of the carry bit is shifted into the MSB of the accumulator, then the LSB of the accumulator is shifted into the carry bit.

Words 1

Cycles

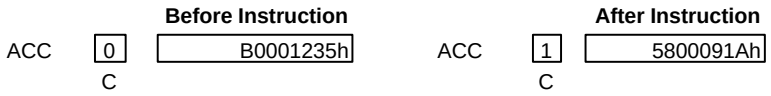
Cycles for a Single ROR Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an ROR Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example R0R



Syntax

RPT *dma*	Direct addressing
RPT *ind* [, **AR***n*]	Indirect addressing
RPT #*k*	Short immediate

Operands

dma:	7 LSBs of the data-memory address
n:	Value from 0 to 7 designating the next auxiliary register
k:	8-bit short immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

RPT *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	0	1	1	0	dma						

RPT *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	0	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

RPT #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	0	1	1	k							

Execution

Increment PC, then ...

Event	*Addressing mode*
(data-memory address) → RPTC	Direct or indirect
k → RPTC	Short immediate

Status Bits

None

Description

The repeat counter (RPTC) is loaded with the content of the addressed data-memory location if direct or indirect addressing is used; it is loaded with an 8-bit immediate value if short immediate addressing is used. The instruction following the RPT is repeated *n* times, where *n* is the initial value of the RPTC plus 1. Since the RPTC cannot be saved during a context switch, repeat loops are regarded as multicycle instructions and are not interruptible. The RPTC is cleared to 0 on a device reset.

RPT is especially useful for block moves, multiply/accumulates, and normalization. The repeat instruction itself is not repeatable.

Words 1

Cycles

Cycles for a Single RPT Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Single RPT Instruction (Using Short Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
1	1	1	1+p

Example 1

```
RPT   DAT127    ;(DP = 31: addresses 0F80h–0FFFh)
                ;Repeat next instruction 13 times.
```

	Before Instruction		**After Instruction**
Data Memory 0FFFh	0Ch	Data Memory 0FFFh	0Ch
RPTC	0h	RPTC	0Ch

Example 2

```
RPT   *,AR1     ;Repeat next instruction 4096 times.
```

	Before Instruction		**After Instruction**
ARP	0	ARP	1
AR0	300h	AR0	300h
Data Memory 300h	0FFFh	Data Memory 300h	0FFFh
RPTC	0h	RPTC	0FFFh

Example 3

```
RPT   #1        ;Repeat next instruction two times.
```

	Before Instruction		**After Instruction**
RPTC	0h	RPTC	1h

Syntax

SACH *dma* [, *shift2*] Direct addressing
SACH *ind* [, *shift2* [, **AR***n*]] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
shift2: Left shift value from 0 to 7 (defaults to 0)
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SACH *dma* [, *shift2*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	1	shift2			0	dma						

SACH *ind* [, *shift2*[, **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	1	shift2			1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
16 MSBs of ((ACC) $\times$ 2^{shift2}) $\rightarrow$ data-memory address

Status Bits

This instruction is not affected by SXM

Description

The SACH instruction copies the entire accumulator into the output shifter, where it left shifts the entire 32-bit number from 0 to 7 bits. It then copies the upper 16 bits of the shifted value into data memory. During the shift, the low-order bits are filled with 0s, and the high-order bits are lost. The accumulator itself remains unaffected.

Words

1

Cycles

Cycles for a Single SACH Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SACH Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
SACH    DAT10,1   ;(DP = 4: addresses 0200h–027Fh,
                  ;left shift of 1)
```

	Before Instruction		After Instruction	
ACC	4208001h		ACC	4208001h
C	X		C	X
Data Memory 20Ah	0h		Data Memory 20Ah	0841h

Example 2

```
SACH    *+,0,AR2  ;(No shift)
```

	Before Instruction		After Instruction	
ARP	1		ARP	2
AR1	300h		AR1	301h
ACC	4208001h		ACC	4208001h
C	X		C	X
Data Memory 300h	0h		Data Memory 300h	0420h

Syntax

SACL *dma* [, *shift2*] Direct addressing
SACL ind [, *shift2* [, **AR***n*]] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
shift2: Left shift value from 0 to 7 (defaults to 0)
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SACL *dma* [, *shift2*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	0	shift2			0	dma						

SACL *ind* [, *shift2* [, **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	0	shift2			1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
16 LSBs of ((ACC) $\times$ 2^{shift2}) $\rightarrow$ data-memory address

Status Bits

This instruction is not affected by SXM.

Description

The SACL instruction copies the entire accumulator into the output shifter, where it left shifts the entire 32-bit number from 0 to 7 bits. It then copies the lower 16 bits of the shifted value into data memory. During the shift, the low-order bits are filled with 0s, and the high-order bits are lost. The accumulator itself remains unaffected.

Words

1

Cycles

Cycles for a Single SACL Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block.

Cycles for a Repeat (RPT) Execution of an SACL Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block.

Example 1

```
SACL    DAT11,1   ;(DP = 4: addresses 0200h–027Fh,
                  ;left shift of 1)
```

	Before Instruction		After Instruction	
ACC	X	7C63 8421	ACC	X 7C63 8421h
	C			C
Data Memory 20Bh		05h	Data Memory 20Bh	0842h

Example 2

```
SACL    *,0,AR7   ;(No shift)
```

	Before Instruction		After Instruction	
ARP		6	ARP	7
AR6		300h	AR6	300h
ACC	X	00FF 8421h	ACC	X 00FF 8421h
	C			C
Data Memory 300h		05h	Data Memory 300h	8421h

Syntax

SAR AR*x*, *dma* Direct addressing
SAR AR*x,* ind [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
x: Value from 0 to 7 designating the auxiliary register value to be stored
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SAR AR*x*, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	0	x			0	dma						

SAR AR*x*, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	0	x			0	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ARx) → data-memory address

Status Bits

None

Description

The content of the designated auxiliary register (ARx) is stored in the specified data-memory location. When the content of the designated auxiliary register is also modified by the instruction (in indirect addressing mode), SAR copies the auxiliary register value to data memory before it increments or decrements the contents of the auxiliary register.

Words

1

Cycles

Cycles for a Single SAR Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SAR Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
SAR     AR0,DAT30 ;(DP = 6: addresses 0300h–037Fh)
```

	Before Instruction		After Instruction
AR0	37h	AR0	37h
Data Memory 31Eh	18h	Data Memory 31Eh	37h

Example 2

```
SAR     AR0,*+
```

	Before Instruction		After Instruction
ARP	0	ARP	0
AR0	401h	AR0	402h
Data Memory 401h	0h	Data Memory 401h	401h

Syntax **SBRK** #*k* Short immediate addressing

Operands k: 8-bit positive short immediate value

Opcode

SBRK #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	1	0	0	k							

Execution

Increment PC, then ...
(current AR) – k → current AR

Note that k is an 8-bit positive constant.

Status Bits None

Description The 8-bit immediate value is subtracted, right justified, from the content of the current auxiliary register (the one pointed to by the ARP) and the result replaces the contents of the auxiliary register. The subtraction takes place in the auxiliary register arithmetic unit (ARAU), with the immediate value treated as an 8-bit positive integer. All arithmetic operations on the auxiliary registers are unsigned.

Words 1

Cycles

Cycles for a Single SBRK Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Example

```
SBRK    #0FFh
```

	Before Instruction		After Instruction
ARP	7	ARP	7
AR7	0h	AR7	FF01h

Syntax **SETC** *control bit*

Operands

control bit: Select one of the following control bits:

C	Carry bit of status register ST1
CNF	RAM configuration control bit of status register ST1
INTM	Interrupt mode bit of status register ST0
OVM	Overflow mode bit of status register ST0
SXM	Sign-extension mode bit of status register ST1
TC	Test/control flag bit of status register ST1
XF	XF pin status bit of status register ST1

Opcode

SETC C

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	1	1	1	1

SETC CNF

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	1	0	1

SETC INTM

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	0	0	1

SETC OVM

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	0	1	1

SETC SXM

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	0	1	1	1

SETC TC

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	1	0	1	1

SETC XF

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	0	1	1	0	1

Execution

Increment PC, then ...
1 → control bit

Status Bits None

Description The specified control bit is set to 1. Note that LST may also be used to load ST0 and ST1. See Section 3.5, *Status Registers ST0 and ST1*, on page 3-15 for more information on each control bit.

Words 1

Cycles

Cycles for a Single SETC Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an SETC Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
SETC    TC      ;TC is bit 11 of ST1
```

	Before Instruction		After Instruction
ST1	x1xxh	ST1	x9xxh

Syntax **SFL**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	1	0	0	1

Execution

Increment PC, then ...
(ACC(31)) $\rightarrow$ C
(ACC(30:0)) $\rightarrow$ ACC(31:1)
0 $\rightarrow$ ACC(0)

Status Bits

Affects
C

This instruction is not affected by SXM.

Description The SFL instruction shifts the entire accumulator left one bit. The LSB is filled with a 0, and the MSB is shifted into the carry bit (C). SFL, unlike SFR, is unaffected by SXM.

Words 1

Cycles

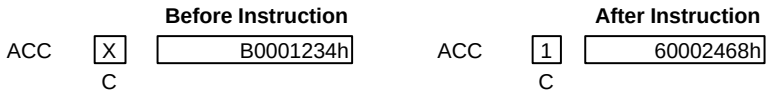
Cycles for a Single SFL Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an SFL Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example SFL



Syntax **SFR**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	1	0	1	0

Execution

Increment PC, then ...
If SXM = 0
  Then 0 → ACC(31).
If SXM = 1
  Then (ACC(31)) → ACC(31)

(ACC(31:1)) → ACC(30:0)
(ACC(0)) → C

Status Bits

Affected by	*Affects*
SXM	C

Description

The SFR instruction shifts the accumulator right one bit.

- If SXM = 1, the instruction produces an arithmetic right shift. The sign bit (MSB) is unchanged and is also copied into bit 30. Bit 0 is shifted into the carry bit (C).
- If SXM = 0, the instruction produces a logic right shift. All of the accumulator bits are shifted right by one bit. The LSB is shifted into the carry bit, and the MSB is filled with a 0.

Words 1

Cycles

Cycles for a Single SFR Instruction

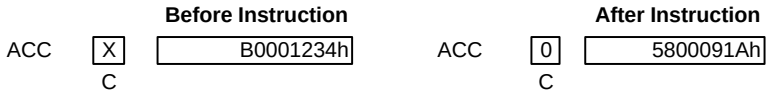
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an SFR Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

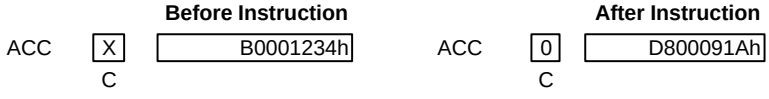
Example 1

```
SFR     ;(SXM = 0: no sign extension)
```



Example 2

```
SFR     ;(SXM = 1: sign extend)
```



Syntax **SPAC**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	0	0	0	0	1	0	1

Execution Increment PC, then ...
(ACC) – shifted (PREG) → ACC

Status Bits

Affected by	*Affects*
PM and OVM	C and OV

This instruction is not affected by SXM.

Description The content of PREG, shifted as defined by the PM status bits, is subtracted from the content of the accumulator. The result is stored in the accumulator. SPAC is not affected by SXM, and the PREG value is always sign extended.

The function of the SPAC instruction is a subtask of the LTS, MPYS, and SQRS instructions.

Words 1

Cycles

Cycles for a Single SPAC Instruction

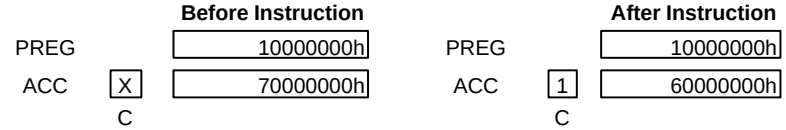
ROM	DARAM	SARAM	External
1	1	1	1+p

Cycles for a Repeat (RPT) Execution of an SPAC Instruction

ROM	DARAM	SARAM	External
n	n	n	n+p

Example

```
SPAC        ;(PM = 0)
```



Syntax

SPH *dma* Direct addressing
SPH *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SPH *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	0	1	0	dma						

SPH *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
16 MSBs of shifted (PREG) → data-memory address

Status Bits

Affected by
PM

Description

The 16 high-order bits of the PREG, shifted as specified by the PM bits, are stored in data memory. First, the 32-bit PREG value is copied into the product shifter, where it is shifted as specified by the PM bits. If the right-shift-by-6 mode is selected, the high-order bits are sign extended and the low-order bits are lost. If a left shift is selected, the high-order bits are lost and the low-order bits are zero filled. If PM = 00, no shift occurs. Then the 16 MSBs of the shifted value are stored in data memory. Neither the PREG value nor the accumulator value is modified by this instruction.

Words

1

Cycles

Cycles for a Single SPH Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SPH Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
SPH     DAT3    ;(DP = 4: addresses 0200h–027Fh,
                ;PM = 0: no shift)
```

	Before Instruction		After Instruction
PREG	FE079844h	PREG	FE079844h
Data Memory 203h	4567h	Data Memory 203h	FE07h

Example 2

```
SPH     *,AR7   ;(PM = 2: left shift of four)
```

	Before Instruction		After Instruction
ARP	6	ARP	7
AR6	203h	AR6	203h
PREG	FE079844h	PREG	FE079844h
Data Memory 203h	4567h	Data Memory 203h	E079h

Syntax

SPL *dma* Direct addressing
SPL *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SPL *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	0	0	0	dma						

SPL *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
16 LSBs of shifted (PREG) → data-memory address

Status Bits

Affected by
PM

Description

The 16 low-order bits of the PREG, shifted as specified by the PM bits, are stored in data memory. First, the 32-bit PREG value is copied into the product shifter, where it is shifted as specified by the PM bits. If the right-shift-by-6 mode is selected, the high-order bits are sign extended and the low-order bits are lost. If a left shift is selected, the high-order bits are lost and the low-order bits are zero filled. If PM = 00, no shift occurs. Then the 16 LSBs of the shifted value are stored in data memory. Neither the PREG value nor the accumulator value is modified by this instruction.

Words

1

Cycles

Cycles for a Single SPL Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SPL Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
SPL     DAT5    ;(DP = 4: addresses 0200h–027Fh,
                ;PM = 2: left shift of four)
```

	Before Instruction		After Instruction
PREG	0FE079844h	PREG	0FE079844h
Data Memory 205h	4567h	Data Memory 205h	08440h

Example 2

```
SPL     *,AR3   ;(PM = 0: no shift)
```

	Before Instruction		After Instruction
ARP	2	ARP	3
AR2	205h	AR2	205h
PREG	0FE079844h	PREG	0FE079844h
Data Memory 205h	4567h	Data Memory 205h	09844h

Syntax

SPLK #*lk*, *dma* Direct addressing
SPLK #*lk*, *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
lk: 16-bit long immediate value
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SPLK *#lk, dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	1	0	0	dma						
lk															

SPLK *#lk, ind* [*,* **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	1	0	1	ARU			N	NAR		
lk															

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
lk → data-memory address

Status Bits

None

Description

The SPLK instruction allows a full 16-bit pattern to be written into any data memory location.

Words

2

Cycles

Cycles for a Single SPLK Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	2	2	2	2+2p
SARAM	2	2	2, 3†	2+2p
External	3+d	3+d	3+d	5+d+2p

† If the operand and the code are in the same SARAM block

Example 1

```
SPLK    #7FFFh,DAT3          ;(DP = 6)
```

	Before Instruction		After Instruction
Data Memory 303h	FE07h	Data Memory 303h	7FFFh

Example 2

```
SPLK    #1111h,*+,AR4
```

	Before Instruction		After Instruction
ARP	0	ARP	4
AR0	300h	AR0	301h
Data Memory 300h	07h	Data Memory 300h	1111h

Syntax **SPM** *constant*

Operands constant: Value from 0 to 3 that determines the product shift mode

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	0	0	0	0	0	0	constant	

Execution Increment PC, then ...
constant → product shift mode (PM) bits

Status Bits *Affects*
PM

This instruction is not affected by SXM.

Description The two LSBs of the instruction word are copied into the product shift mode (PM) bits of status register ST1 (bits 1 and 0 of ST1). The PM bits control the mode of the shifter at the output of the PREG. This shifter can shift the PREG output either one or four bits to the left or six bits to the right. The possible PM bit combinations and their meanings are shown in Table 8–8. When an instruction accesses the PREG value, the value first passes through the shifter, where it is shifted by the specified amount.

Table 8–8. Product Shift Modes

PM Field	Specified Product Shift
00	No shift of PREG output
01	PREG output to be left shifted 1 place
10	PREG output to be left shifted 4 places
11	PREG output to be right shifted 6 places and sign extended

The left shifts allow the product to be justified for fractional arithmetic. The right-shift-by-six mode allows up to 128 multiply accumulate processes without the possibility of overflow occurring. PM may also be loaded by an LST #1 instruction.

Words 1

Cycles

Cycles for a Single SPM Instruction

ROM	DARAM	SARAM	External
1	1	1	1+p

Example

```
SPM   3     ;Product register shift mode 3 (PM = 11)
            ;is selected causing all subsequent
            ;transfers from the product register (PREG)
            ;to be shifted to the right six places.
```

Syntax

SQRA *dma* Direct addressing
SQRA *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SQRA *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	1	0	0	dma						

SQRA *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) + shifted (PREG) → ACC
(data-memory address) → TREG
(TREG) × (data-memory address) → PREG

Status Bits

Affected by	*Affects*
OVM and PM	OV and C

Description

The content of the PREG, shifted as defined by the PM status bits, is added to the accumulator. Then the addressed data-memory value is loaded into the TREG, squared, and stored in the PREG.

Words

1

Cycles

Cycles for a Single SQRA Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SQRA Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
SQRA    DAT30    ;(DP = 6: addresses 0300h–037Fh,
                 ;PM = 0: no shift of product)
```

	Before Instruction		After Instruction	
Data Memory			Data Memory	
31Eh		0Fh	31Eh	0Fh
TREG		3h	TREG	0Fh
PREG		12Ch	PREG	0E1h
ACC	X	1F4h	ACC	0	320h
	C			C	

Example 2

```
SQRA    *,AR4    ;(PM = 0)
```

	Before Instruction		After Instruction	
ARP		3	ARP	4
AR3		31Eh	AR3	31Eh
Data Memory			Data Memory	
31Eh		0Fh	31Eh	0Fh
TREG		3h	TREG	0Fh
PREG		12Ch	PREG	0E1h
ACC	X	1F4h	ACC	0	320h
	C			C	

Syntax

SQRS *dma* Direct addressing
SQRS *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SQRS *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	1	1	0	dma						

SQRS *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	0	0	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) – shifted (PREG) → ACC
(data-memory address) → TREG
(TREG) × (data-memory address) → PREG

Status Bits

Affected by	*Affects*
OVM and PM	OV and C

Description

The content of the PREG, shifted as defined by the PM status bits, is subtracted from the accumulator. Then the addressed data-memory value is loaded into the TREG, squared, and stored in the PREG.

Words

1

Cycles

Cycles for a Single SQRS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SQRS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
SQRS    DAT9    ;(DP = 6: addresses 0300h–037Fh,
                ;PM = 0: no shift of product)
```

	Before Instruction	After Instruction
Data Memory 309h	08h	08h
TREG	1124h	08h
PREG	190h	40h
ACC	1450h	12C0h
C	X	1

Example 2

```
SQRS    *,AR5   ;(PM = 0)
```

	Before Instruction	After Instruction
ARP	3	5
AR3	309h	309h
Data Memory 309h	08h	08h
TREG	1124h	08h
PREG	190h	40h
ACC	1450h	12C0h
C	X	1

Syntax

SST #*m*, *dma* Direct addressing
SST #*m*, *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
m: Select one of the following:
0 Indicates that ST0 will be stored
1 Indicates that ST1 will be stored
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SST #0, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	1	0	0	dma						

SST #0, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

SST #1, *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	1	1	0	dma						

SST #1, *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(status register STm) → data-memory address

Status Bits

None

Description

Status register ST0 or ST1 (whichever is specified) is stored in data memory.

In direct addressing mode, the specified status register is always stored in page 0, regardless of the value of the data page pointer (DP) in ST0. Although the processor automatically accesses page 0, the DP is not physically modified; this allows the DP value to be stored unchanged when ST0 is stored. The specific storage location within page 0 is given in the instruction.

In indirect addressing mode, the storage address is obtained from the auxiliary register selected; thus, the specified status register contents can be stored to an address on any page in data memory.

Status registers ST0 and ST1 are defined in Section 3.5, *Status Registers ST0 and ST1*, on page 3-15.

Words 1

Cycles

Cycles for a Single SST Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	2+d	2+d	2+d	4+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SST Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+2†	n+p
External	2n+nd	2n+nd	2n+nd	2n+2+nd+p

† If the operand and the code are in the same SARAM block

Example 1

```
SST     #0,96     (Direct addressing: data page 0
                  ;accessed automatically)
```

	Before Instruction		After Instruction
ST0	0A408h	ST0	0A408h
Data Memory 60h	0Ah	Data Memory 60h	0A408h

Example 2

```
SST     #1,*,AR7  (Indirect addressing)
```

	Before Instruction		After Instruction
ARP	0	ARP	7
AR0	300h	AR0	300h
ST1	2580h	ST1	2580h
Data Memory 300h	0h	Data Memory 300h	2580h

Syntax

SUB *dma* [, *shift*]	Direct addressing
SUB *dma*,**16**	Direct with left shift of 16
SUB *ind* [,*shift* [, **AR***n*]]	Indirect addressing
SUB *ind*,**16**[, **AR***n*]	Indirect with left shift of 16
SUB #*k*	Short immediate
SUB #*lk* [,*shift*]	Long immediate

Operands

dma:	7 LSBs of the data-memory address
shift:	Left shift value from 0 to 15 (defaults to 0)
n:	Value from 0 to 7 designating the next auxiliary register
k:	8-bit short immediate value
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SUB *dma* [,*shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	shift				0	dma						

SUB *dma,* **16**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	0	1	0	dma						

SUB *ind* [, *shift* [, **AR***n*]]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	shift				1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

SUB *ind*,**16** [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	0	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

SUB #*k*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	0	1	0	k							

SUB #*lk* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	1	0	1	0	shift			
lk															

Execution

Increment PC, then ...

Event	*Addressing mode*
(ACC) – ((data-memory address) $\times$ 2^{shift}) $\rightarrow$ ACC	Direct or indirect
(ACC) – ((data-memory address) $\times$ 2^{16}) $\rightarrow$ ACC	Direct or indirect (shift of 16)
(ACC) – k $\rightarrow$ ACC	Short immediate
(ACC) – lk $\times$ 2^{shift} $\rightarrow$ ACC	Long immediate

Status Bits

Affected by	*Affects*	*Addressing mode*
OVM and SXM	OV and C	Direct or indirect
OVM	OV and C	Short immediate
OVM and SXM	OV and C	Long immediate

Description

In direct, indirect, and long immediate addressing, the content of the addressed data-memory location or a 16-bit constant are left shifted and subtracted from the accumulator. During shifting, low-order bits are zero filled. High-order bits are sign extended if SXM = 1 and zero filled if SXM = 0. The result is then stored in the accumulator.

If short immediate addressing is used, an 8-bit positive constant is subtracted from the accumulator. In this case, no shift value may be specified, the subtraction is unaffected by SXM, and the instruction is not repeatable.

Normally, the carry bit is cleared (C = 0) if the result of the subtraction generates a borrow and is set (C = 1) if it does not generate a borrow. However, if a 16-bit shift is specified with the subtraction, the instruction clears the carry bit if a borrow is generated but does not affect the carry bit otherwise.

Words

Words	*Addressing mode*
1	Direct, indirect or short immediate
2	Long immediate

Cycles

Cycles for a Single SUB Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block.

Cycles for a Repeat (RPT) Execution of an SUB Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block.

Cycles for a Single SUB Instruction (Using Short Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
1	1	1	1+p

Cycles for a Single SUB Instruction (Using Long Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+2p

Example 1

```
SUB     DAT80     ;(DP = 8: addresses 0400h–047Fh,
                  ;SXM=0: sign-extension suppressed)
```

	Before Instruction		After Instruction	
Data Memory 450h		11h		11h
ACC	X	24h	1	13h
	C		C	

Example 2

```
SUB     *–,1,AR0  ;(Left shift by 1, SXM = 0)
```

	Before Instruction		After Instruction	
ARP		7		0
AR7		301h		300h
Data Memory 301h		04h		04h
ACC	X	09h	1	01h
	C		C	

Example 3

```
SUB     #8h       ;(SXM = 1: sign-extension mode)
```

	Before Instruction		After Instruction	
ACC	X	07h	0	FFFFFFFFh
	C		C	

Example 4

```
SUB     #0FFFh,4  ;(Left shift by four, SXM = 0)
```

	Before Instruction		After Instruction	
ACC	X	0FFFFh	1	0Fh
	C		C	

Syntax

SUBB *dma* Direct addressing
SUBB *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SUBB *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	0	0	0	dma						

SUBB *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) – (data-memory address) – (logical inversion of C) → ACC

Status Bits

Affected by	*Affects*
OVM	OV and C

This instruction is not affected by SXM.

Description

The content of the addressed data-memory location and the logical inversion of the carry bit is subtracted from the accumulator with sign extension suppressed. The carry bit is then affected in the normal manner: the carry bit is cleared (C = 0) if the result of the subtraction generates a borrow and is set (C = 1) if it does not generate a borrow.

The SUBB instruction can be used in performing multiple-precision arithmetic.

Words

1

Cycles

Cycles for a Single SUBB Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

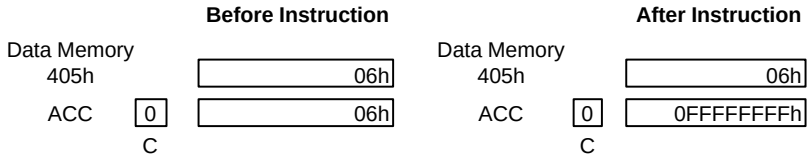
Cycles for a Repeat (RPT) Execution of an SUBB Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

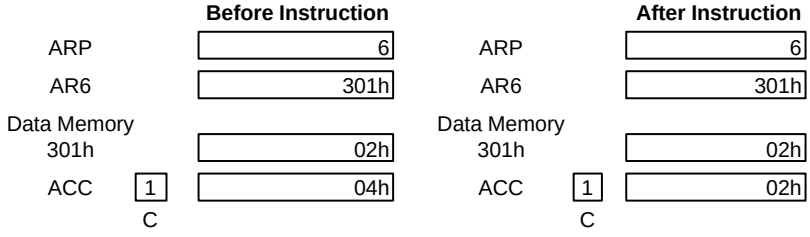
Example 1

```
SUBB    DAT5    ;(DP = 8: addresses 0400h–047Fh)
```



Example 2

```
SUBB    *
```



In the first example, C is originally zeroed, presumably from the result of a previous subtract instruction that performed a borrow. The effective operation performed was 6 – 6 – (1) = –1, generating another borrow (resetting carry) in the process. In the second example, no borrow was previously generated (C = 1), and the result from the subtract instruction does not generate a borrow.

Syntax

SUBC *dma* Direct addressing
SUBC *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SUBC *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	0	1	0	0	dma						

SUBC *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	0	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

For (ACC) ≥ 0 and (data-memory address) ≥ 0:

Increment PC, then ...
(ACC) – [(data-memory address) × 2^{15}] → ALU output

If ALU output ≥ 0
Then (ALU output) × 2 + 1 → ACC
Else (ACC) × 2 → ACC

Status Bits

Affects
OV and C

Description

The SUBC instruction performs conditional subtraction, which can be used for division as follows: Place a positive 16-bit dividend in the low accumulator and clear the high accumulator. Place a 16-bit positive divisor in data memory. Execute SUBC 16 times. After completion of the last SUBC, the quotient of the division is in the lower-order 16 bits of the accumulator, and the remainder is in the higher-order 16 bits of the accumulator. For negative accumulator and/or data-memory values, SUBC cannot be used for division.

If the 16-bit dividend contains fewer than 16 significant bits, the dividend may be placed in the accumulator and left shifted by the number of leading non-significant 0s. The number of executions of SUBC is reduced from 16 by that number. One leading 0 is always significant.

SUBC operations performed as previously stated are not affected by the sign-extension mode bit (SXM).

SUBC affects OV but is not affected by OVM; therefore, the accumulator does not saturate upon positive or negative overflows when executing this instruction. The carry bit is affected in the normal manner during this instruction: the carry bit is cleared (C = 0) if the result of the subtraction generates a borrow and is set (C = 1) if it does not generate a borrow.

Words 1

Cycles

Cycles for a Single SUBC Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

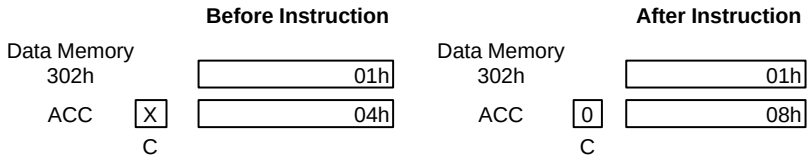
Cycles for a Repeat (RPT) Execution of an SUBC Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

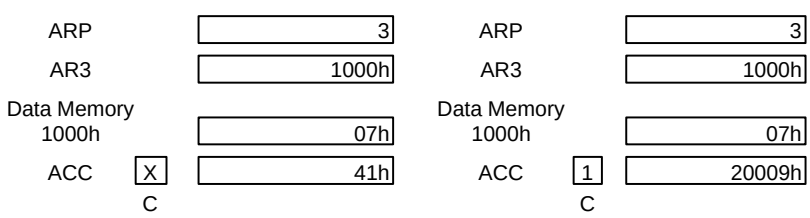
Example 1

```
SUBC    DAT2    ;(DP = 6)
```



Example 2

```
RPT     #15
SUBC    *
```



Syntax

SUBS *dma* Direct addressing
SUBS *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SUBS *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	1	0	0	dma						

SUBS *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) – (data-memory address) → ACC

Status Bits

Affected by	*Affects*
OVM	OV and C

This instruction is not affected by SXM.

Description

The content of the specified data-memory location is subtracted from the accumulator with sign extension suppressed. The data is treated as a 16-bit unsigned number, regardless of SXM. The accumulator behaves as a signed number. SUBS produces the same results as a SUB instruction with SXM = 0 and a shift count of 0.

The carry bit is cleared (C = 0) if the result of the subtraction generates a borrow and is set (C = 1) if it does not generate a borrow.

Words

1

Cycles

Cycles for a Single SUBS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an SUBS Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
SUBS    DAT2        ;(DP = 16, SXM = 1)
```

	Before Instruction		After Instruction	
Data Memory 802h		0F003h		0F003h
ACC	X	0F105h	1	102h
C				

Example 2

```
SUBS    *           ;(SXM = 1)
```

	Before Instruction		After Instruction	
ARP		0		0
AR0		310h		310h
Data Memory 310h		0F003h		0F003h
ACC	X	0FFFF105h	1	0FFF0102h
C				

Syntax

SUBT *dma* Direct addressing
SUBT *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

SUBT *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	1	1	0	dma						

SUBT *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(ACC) – [(data-memory address) $\times$ $2^{(TREG(3:0))}$] $\rightarrow$ (ACC)

If SXM = 1
Then (data-memory address) is sign-extended.
If SXM = 0
Then (data-memory address) is not sign-extended.

Status Bits

Affected by	*Affects*
OVM and SXM	OV and C

Description

The data-memory value is left shifted and subtracted from the accumulator. The left shift is defined by the four LSBs of TREG, resulting in shift options from 0 to 15 bits. The result replaces the accumulator contents. Sign extension on the data-memory value is controlled by the SXM status bit.

The carry bit is cleared (C = 0) if the result of the subtraction generates a borrow and is set (C = 1) if it does not generate a borrow.

Words

1

Cycles

Cycles for a Single SUBT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block.

Cycles for a Repeat (RPT) Execution of an SUBT Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block.

Example 1

```
SUBT    DAT127    ;(DP = 5: addresses 0280h–02FFh)
```

	Before Instruction		After Instruction	
Data Memory 2FFh	06h		06h	
TREG	08h		08h	
ACC	0FDA5h	X	0F7A5h	1
		C		C

Example 2

```
SUBT    *
```

	Before Instruction		After Instruction	
ARP	1		1	
AR1	800h		800h	
Data Memory 800h	01h		01h	
TREG	08h		08h	
ACC	0h	X	FFFFFF00h	0
		C		C

Syntax

TBLR *dma* Direct addressing
TBLR *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

TBLR *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	1	0	0	dma						

TBLR *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	1	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...

(PC) → MSTACK
(ACC(15:0)) → PC
(pma) → data-memory address
For indirect, modify (current AR) and (ARP) as specified,
(PC) + 1 → PC

While (repeat counter) ≠ 0
  (pma) → data-memory address
  For indirect, modify (current AR) and (ARP) as specified,
  (PC) + 1 → PC
  (repeat counter) –1 → repeat counter.

(MSTACK) → PC

Status Bits

None

Description

The TBLR instruction transfers a word from a location in program memory to a data-memory location specified by the instruction. The program-memory address is defined by the low-order 16 bits of the accumulator. For this operation, a read from program memory is performed, followed by a write to data memory. When repeated with the repeat (RPT) instruction, TBLR effectively becomes a single-cycle instruction, and the program counter that was loaded with (ACC(15:0)) is incremented once each cycle.

Words

1

Cycles

Cycles for a Single TBLR Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM/ROM Destination: DARAM	3	3	3	$3+p_{code}$
Source: SARAM Destination: DARAM	3	3	3	$3+p_{code}$
Source: External Destination: DARAM	$3+p_{src}$	$3+p_{src}$	$3+p_{src}$	$3+p_{src}+p_{code}$
Source: DARAM/ROM Destination: SARAM	3	3	3 4†	$3+p_{code}$
Source: SARAM Destination: SARAM	3	3	3 4†	$3+p_{code}$
Source: External Destination: SARAM	$3+p_{src}$	$3+p_{src}$	$3+p_{src}$ $4+p_{src}$†	$3+p_{src}+p_{code}$
Source: DARAM/ROM Destination: External	$4+d_{dst}$	$4+d_{dst}$	$4+d_{dst}$	$6+d_{dst}+p_{code}$
Source: SARAM Destination: External	$4+d_{dst}$	$4+d_{dst}$	$4+d_{dst}$	$6+d_{dst}+p_{code}$
Source: External Destination: External	$4+p_{src}+d_{dst}$	$4+p_{src}+d_{dst}$	$4+p_{src}+d_{dst}$	$6+p_{src}+d_{dst}+p_{code}$

† If the destination operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a TBLR Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM/ROM Destination: DARAM	n+2	n+2	n+2	$n+2+p_{code}$
Source: SARAM Destination: DARAM	n+2	n+2	n+2	$n+2+p_{code}$
Source: External Destination: DARAM	$n+2+np_{src}$	$n+2+np_{src}$	$n+2+np_{src}$	$n+2+np_{src}+p_{code}$

† If the destination operand and the code are in the same SARAM block
‡ If both the source and the destination operands are in the same SARAM block
§ If both operands and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a TBLR Instruction (Continued)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM/ROM Destination: SARAM	n+2	n+2	n+2 n+4†	n+2+p_{code}
Source: SARAM Destination: SARAM	n+2 2n‡	n+2 2n‡	n+2 2n‡ 2n+2§	n+2+p_{code} 2n‡
Source: External Destination: SARAM	n+2+np_{src}	n+2+np_{src}	n+2+np_{src} n+4+np_{src}†	n+2+np_{src}+p_{code}
Source: DARAM/ROM Destination: External	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+4+nd_{dst}+p_{code}
Source: SARAM Destination: External	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+2+nd_{dst}	2n+4+nd_{dst}+p_{code}
Source: External Destination: External	4n+np_{src}+nd_{dst}	4n+np_{src}+nd_{dst}	4n+np_{src}+nd_{dst}	4n+2+np_{src}+nd_{dst}+ p_{code}

† If the destination operand and the code are in the same SARAM block
‡ If both the source and the destination operands are in the same SARAM block
§ If both operands and the code are in the same SARAM block

Example 1

```
TBLR    DAT6    ;(DP = 4: addresses 0200h–027Fh)
```

	Before Instruction		After Instruction
ACC	23h	ACC	23h
Program Memory 23h	306h	Program Memory 23h	306h
Data Memory 206h	75h	Data Memory 206h	306h

Example 2

```
TBLR    *,AR7
```

	Before Instruction		After Instruction
ARP	0	ARP	7
AR0	300h	AR0	300h
ACC	24h	ACC	24h
Program Memory 24h	307h	Program Memory 24h	307h
Data Memory 300h	75h	Data Memory 300h	307h

Syntax

TBLW *dma* Direct addressing
TBLW *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

TBLW *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	1	1	0	dma						

TBLW *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	1	1	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(PC+1) → MSTACK
(ACC(15:0)) → PC+1
(data-memory address) → pma,
For indirect, modify (current AR) and (ARP) as specified
(PC) + 1 → PC

While (repeat counter) ≠ 0
  (data-memory address) → pma,
  For indirect, modify (current AR) and (ARP) as specified
  (PC) + 1 → PC
  (repeat counter) –1 → repeat counter.

(MSTACK) → PC+1

Status Bits

None

Description

The TBLW instruction transfers a word in data memory to program memory. The data-memory address is specified by the instruction, and the program-memory address is specified by the lower 16 bits of the accumulator. A read from data memory is followed by a write to program memory to complete the instruction. When repeated with the repeat (RPT) instruction, TBLW effectively becomes a single-cycle instruction, and the program counter that was loaded with (ACC(15:0)) is incremented once each cycle.

Words

1

Cycles

Cycles for a Single TBLW Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM/ROM Destination: DARAM	3	3	3	$3+p_{code}$
Source: SARAM Destination: DARAM	3	3	3	$3+p_{code}$
Source: External Destination: DARAM	$3+d_{src}$	$3+d_{src}$	$3+d_{src}$	$3+d_{src}+p_{code}$
Source: DARAM/ROM Destination: SARAM	3	3	3 4†	$3+p_{code}$
Source: SARAM Destination: SARAM	3	3	3 4†	$3+p_{code}$
Source: External Destination: SARAM	$3+d_{src}$	$3+d_{src}$	$3+d_{src}$ $4+d_{src}$†	$3+d_{src}+p_{code}$
Source: DARAM/ROM Destination: External	$4+p_{dst}$	$4+p_{dst}$	$4+p_{dst}$	$5+p_{dst}+p_{code}$
Source: SARAM Destination: External	$4+p_{dst}$	$4+p_{dst}$	$4+p_{dst}$	$5+p_{dst}+p_{code}$
Source: External Destination: External	$4+d_{src}+p_{dst}$	$4+d_{src}+p_{dst}$	$4+d_{src}+p_{dst}$	$5+d_{src}+p_{dst}+p_{code}$

† If the destination operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a TBLW Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM/ROM Destination: DARAM	n+2	n+2	n+2	$n+2+p_{code}$
Source: SARAM Destination: DARAM	n+2	n+2	n+2	$n+2+p_{code}$
Source: External Destination: DARAM	$n+2+nd_{src}$	$n+2+nd_{src}$	$n+2+nd_{src}$	$n+2+nd_{src}+p_{code}$

† If the destination operand and the code are in the same SARAM block
‡ If both the source and the destination operands are in the same SARAM block
§ If both operands and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a TBLW Instruction (Continued)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
Source: DARAM/ROM Destination: SARAM	n+2	n+2	n+2 n+3†	n+2+p_{code}
Source: SARAM Destination: SARAM	n+2 2n‡	n+2 2n‡	n+2 2n‡ 2n+1§	n+2+p_{code} 2n‡
Source: External Destination: SARAM	n+2+nd_{src}	n+2+nd_{src}	n+2+nd_{src} n+3+nd_{src}†	n+2+nd_{src}+p_{code}
Source: DARAM/ROM Destination: External	2n+2+np_{dst}	2n+2+np_{dst}	2n+2+np_{dst}	2n+3+np_{dst}+p_{code}
Source: SARAM Destination: External	2n+2+np_{dst}	2n+2+np_{dst}	2n+2+np_{dst}	2n+3+np_{dst}+p_{code}
Source: External Destination: External	4n+nd_{src}+np_{dst}	4n+nd_{src}+np_{dst}	4n+nd_{src}+np_{dst}	4n+1+nd_{src}+np_{dst}+ p_{code}

† If the destination operand and the code are in the same SARAM block
‡ If both the source and the destination operands are in the same SARAM block
§ If both operands and the code are in the same SARAM block

Example 1

```
TBLW    DAT5    ;(DP = 32: addresses 1000h–107Fh)
```

	Before Instruction		After Instruction
ACC	257h	ACC	257h
Data Memory 1005h	4339h	Data Memory 1005h	4339h
Program Memory 257h	306h	Program Memory 257h	4399h

Example 2

```
TBLW    *
```

	Before Instruction		After Instruction
ARP	6	ARP	6
AR6	1006h	AR6	1006h
ACC	258h	ACC	258h
Data Memory 1006h	4340h	Data Memory 1006h	4340h
Program Memory 258h	307h	Program Memory 258h	4340h

Syntax **TRAP**

Operands None

Opcode

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	0	1	0	1	0	0	0	1

Execution (PC) + 1 → stack
22h → PC

Status Bits Not affected by INTM; does not affect INTM.

Description The TRAP instruction is a software interrupt that transfers program control to program-memory location 22h and pushes the program counter (PC) plus 1 onto the hardware stack. The instruction at location 22h may contain a branch instruction to transfer control to the TRAP routine. Putting (PC + 1) onto the stack enables a return instruction to pop the return address (which points to the instruction after TRAP) from the stack. The TRAP instruction is not maskable.

Words 1

Cycles

Cycles for a Single TRAP Instruction

ROM	DARAM	SARAM	External
4	4	4	4+3p†

† The processor performs speculative fetching by reading two additional instruction words. If the PC discontinuity is taken, these two instruction words are discarded.

Example

```
TRAP    ;PC + 1 is pushed onto the stack, and then
        ;control is passed to program memory location
        ;22h.
```

Syntax

XOR *dma*	Direct addressing
XOR *ind* [, **AR***n*]	Indirect addressing
XOR #*lk*, [, *shift*]	Long immediate addressing
XOR #*lk*,**16**	Long immediate with left shift of 16

Operands

dma:	7 LSBs of the data-memory address
shift:	Left shift value from 0 to 15 (defaults to 0)
n:	Value from 0 to 7 designating the next auxiliary register
lk:	16-bit long immediate value
ind:	Select one of the following seven options: * *+ *– *0+ *0– *BR0+ *BR0–

Opcode

XOR *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	0	0	0	dma						

XOR *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

XOR #*lk* [, *shift*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	1	1	1	0	1	shift			
lk															

XOR #*lk,* **16**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	1	0	1	0	0	0	0	0	1	1
lk															

Execution

Increment PC, then ...

Event(s)	*Addressing mode*
(ACC(15:0)) XOR (data-memory address) → ACC(15:0) (ACC(31:16)) → ACC(31:16)	Direct or indirect
(ACC(31:0)) XOR lk × 2^{shift} → ACC(31:0)	Long immediate
(ACC(31:0)) XOR lk × 2^{16} → ACC(31:0)	Long immediate with left shift of 16

Status Bits None

Description With direct or indirect addressing, the low half of the accumulator value is exclusive ORed with the content of the addressed data memory location, and the result replaces the low half of the accumulator value; the upper half of the accumulator value is unaffected. With immediate addressing, the long immediate constant is shifted and zero filled on both ends and exclusive ORed with the entire content of the accumulator. The carry bit (C) is unaffected by XOR.

Words

Words	*Addressing mode*
1	Direct or indirect
2	Long immediate

Cycles

Cycles for a Single XOR Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of an XOR Instruction (Using Direct and Indirect Addressing)

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Cycles for a Single XOR Instruction (Using Long Immediate Addressing)

ROM	**DARAM**	**SARAM**	**External**
2	2	2	2+2p

Example 1

```
XOR   DAT127        ;(DP = 511: addresses FF80h–FFFFh)
```

	Before Instruction	After Instruction
Data Memory 0FFFFh	0F0F0h	0F0F0h
ACC	12345678h	1234A688h
C	X	X

Example 2

```
XOR   *+,AR0
```

	Before Instruction	After Instruction
ARP	7	0
AR7	300h	301h
Data Memory 300h	0FFFFh	0FFFFh
ACC	1234F0F0h	12340F0Fh
C	X	X

Example 3

```
XOR   #0F0F0h,4     ;(First shift data value left by
                    ;four)
```

	Before Instruction	After Instruction
ACC	11111010h	111E1F10h
C	X	X

Syntax

ZALR *dma* Direct addressing
ZALR *ind* [, **AR***n*] Indirect addressing

Operands

dma: 7 LSBs of the data-memory address
n: Value from 0 to 7 designating the next auxiliary register
ind: Select one of the following seven options:
* *+ *– *0+ *0– *BR0+ *BR0–

Opcode

ZALR *dma*

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	0	0	0	dma						

ZALR *ind* [, **AR***n*]

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	0	0	1	ARU			N	NAR		

Note: ARU, N, and NAR are defined in subsection 7.3.4, *Indirect Addressing Opcode Format* (page 7-12).

Execution

Increment PC, then ...
(data-memory address) → ACC(31:16)
8000h → ACC(15:0)

Status Bits

None

Description

To load a data-memory value into the high-order half of the accumulator, the ZALR instruction rounds the value by adding 1/2 LSB; that is, the 15 low bits (bits 14–0) of the accumulator are cleared to 0, and bit 15 of the accumulator is set to 1.

Words

1

Cycles

Cycles for a Single ZALR Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	1	1	1	1+p
SARAM	1	1	1, 2†	1+p
External	1+d	1+d	1+d	2+d+p

† If the operand and the code are in the same SARAM block

Cycles for a Repeat (RPT) Execution of a ZALR Instruction

	Program			
Operand	**ROM**	**DARAM**	**SARAM**	**External**
DARAM	n	n	n	n+p
SARAM	n	n	n, n+1†	n+p
External	n+nd	n+nd	n+nd	n+1+p+nd

† If the operand and the code are in the same SARAM block

Example 1

```
ZALR    DAT3    ;(DP = 32: addresses 1000h–107Fh)
```

	Before Instruction	After Instruction
Data Memory 1003h	3F01h	3F01h
ACC	77FFFFh	3F018000h
C	X	X

Example 2

```
ZALR    *–,AR4
```

	Before Instruction	After Instruction
ARP	7	4
AR7	0FF00h	0FEFFh
Data Memory 0FF00h	0E0E0h	0E0E0h
ACC	107777h	0E0E08000h
C	X	X

Appendix A

TMS320C1x/C2x/C24x/C5x Instruction Set Comparison

Note:

The instruction set for the TMS320C24x is identical to that of the TMS320C2xx. All references to 'C2xx devices in this appendix also apply to 'C24x devices.

This appendix contains a table that compares the TMS320C1x, TMS320C2x, TMS320C2xx, and TMS320C5x instructions alphabetically. Each table entry shows the syntax for the instruction, indicates which devices support the instruction, and describes the operation of the instruction. Section A.1 shows a sample table entry and describes the symbols and abbreviations used in the table.

The TMS320C2x, TMS320C2xx, and TMS320C5x devices have enhanced instructions; these instructions are single mnemonics that perform the functions of several similar instructions. Section A.2 summarizes these enhanced instructions.

Topic **Page**

A.1 Using the Instruction Set Comparison Table A-2
A.2 Enhanced Instructions A-5
A.3 Instruction Set Comparison Table A-6

A.1 Using the Instruction Set Comparison Table

To help you read the comparison table, this section provides an example of a table entry and a list of acronyms.

A.1.1 An Example of a Table Entry

In cases where more than one syntax is used, the first syntax is usually for direct addressing and the second is usually for indirect addressing. Where three or more syntaxes are used, the syntaxes are normally specific to a device.

This is how the AND instruction appears in the table:

Syntax	1x	2x	2xx	5x	Description
AND *dma*	√	√	√	√	**AND With Accumulator**
AND {*ind*} [, *next ARP*]	√	√	√	√	TMS320C1x and TMS320C2x devices: AND the contents of the addressed data-memory location with the 16 LSBs of the accumulator. The 16 MSBs of the accumulator are ANDed with 0s.
AND #*lk* [, *shift*]			√	√	TMS320C2xx and TMS320C5x devices: AND the contents of the addressed data-memory location or a 16-bit immediate value with the contents of the accumulator. The 16 MSBs of the accumulator are ANDed with 0s. If a shift is specified, left shift the constant before the AND. Low-order bits below and high-order bits above the shifted value are treated as 0s.

The first column, *Syntax*, states the mnemonic and the syntaxes for the AND instruction.

The checks in the second through the fifth columns, *1x*, *2x*, *2xx,* and *5x*, indicate the devices that can be used with each of the syntaxes.

1x refers to the TMS320C1x devices.
2x refers to the TMS320C2x devices, including TMS320C25.
2xx refers to the TMS320C2xx devices.
5x refers to the TMS320C5x devices.

In this example, you can use the first two syntaxes with TMS320C1x, TMS320C2x, TMS320C2xx, and TMS320C5x devices, but you can use the last syntax only with TMS320C2xx and TMS320C5x devices.

The sixth column, *Description*, briefly describes how the instruction functions. Often, an instruction functions slightly differently for the different devices: read the entire description before using the instruction.

A.1.2 Symbols and Acronyms Used in the Table

Table A–1 lists the instruction set symbols and acronyms used throughout this appendix.

Table A–1. Symbols and Acronyms Used in the Instruction Set Comparison Table

Symbol	Description	Symbol	Description
lk	16-bit immediate value	INTM	Interrupt mask bit
k	8-bit immediate value	INTR	Interrupt mode bit
{ind}	Indirect address	OV	Overflow bit
ACC	Accumulator	P	Program bus
ACCB	Accumulator buffer	PA	Port address
AR	Auxiliary register	PC	Program counter
ARCR	Auxiliary register compare	PM	Product shifter mode
ARP	Auxiliary register pointer	pma	Program-memory address
BMAR	Block move address register	RPTC	Repeat counter
BRCR	Block repeat count register	shift, $shift_n$	Shift value
C	Carry bit	src	Source address
DBMR	Dynamic bit manipulation register	ST	Status register
dma	Data-memory address	SXM	Sign-extension mode bit
DP	Data-memory page pointer	TC	Test/control bit
dst	Destination address	T	Temporary register
FO	Format status list	TREGn	TMS320C5x temporary register (0–2)
FSX	External framing pulse	TXM	Transmit mode status register
IMR	Interrupt mask register	XF	XF pin status bit

Based on the device, this is how the indirect addressing operand {ind} is interpreted:

{ind}	**'C1x**:	{ * | *+ | *– }
	'C2x:	{ * | *+ | *– | *0+| *0– | *BR0+ | *BR0– }
	'C2xx:	{ * | *+ | *– | *0+| *0– | *BR0+ | *BR0– }
	'C5x:	{ * | *+ | *– | *0+| *0– | *BR0+ | *BR0– }

where the possible options are separated by vertical bars (|). For example:

`ADD {`*`ind`*`}`

is interpreted as:

'C1x devices	ADD { * | *+ | *– }
'C2x devices	ADD { * | *+ | *– | *0+ | *0– | *BR0+ | *BR0– }
'C2xx devices	ADD { * | *+ | *– | *0+ | *0– | *BR0+ | *BR0– }
'C5x devices	ADD { * | *+ | *– | *0+ | *0– | *BR0+ | *BR0– }

Based on the device, these are the sets of values for shift, $shift_1$, and $shift_2$:

shift	**'C1x**:	0–15 (shift of 0–15 bits)
	'C2x:	0–15 (shift of 0–15 bits)
	'C2xx:	0–16 (shift of 0–16 bits)
	'C5x:	0–16 (shift of 0–16 bits)
$shift_1$	**'C1x**:	n/a
	'C2x:	0–15 (shift of 0–15 bits)
	'C2xx:	0–16 (shift of 0–16 bits)
	'C5x:	0–16 (shift of 0–16 bits)
$shift_2$	**'C1x**:	n/a
	'C2x:	n/a
	'C2xx:	0–15 (shift of 0–15 bits)
	'C5x:	0–15 (shift of 0–15 bits)

In some cases, the sets are smaller; in these cases, the valid sets are given in the *Description* column of the table.

A.2 Enhanced Instructions

An enhanced instruction is a single mnemonic that performs the functions of several similar instructions. For example, the enhanced instruction ADD performs the ADD, ADDH, ADDK, and ADLK functions and replaces any of these other instructions at assembly time. For example, when a program using ADDH is assembled for the ’C2xx or ’C5x, ADDH is replaced by an ADD instruction that performs the same function. These enhanced instructions are valid for TMS320C2x, TMS320C2xx, and TMS320C5x devices (not TMS320C1x).

Table A–2 summarizes the enhanced instructions and the functions that the enhanced instructions perform (based on TMS320C1x/2x mnemonics).

Table A–2. Summary of Enhanced Instructions

Enhanced Instruction	Includes These Operations
ADD	ADD, ADDH, ADDK, ADLK
AND	AND, ANDK
BCND	BBNZ, BBZ, BC, BCND, BGEZ, BGZ, BIOZ, BLEZ, BLZ, BNC, BNV, BNZ, BV, BZ
BLDD	BLDD, BLKD
BLDP	BLDP, BLKP
CLRC	CLRC, CNFD, EINT, RC, RHM, ROVM, RSXM, RTC, RXF
LACC	LAC, LACC, LALK, ZALH
LACL	LACK, LACL, ZAC, ZALS
LAR	LAR, LARK, LRLK
LDP	LDP, LDPK
LST	LST, LST1
MAR	LARP, MAR
MPY	MPY, MPYK
OR	OR, ORK
RPT	RPT, RPTK
SETC	CNFP, DINT, SC, SETC, SHM, SOVM, SSXM, STC, SXF
SUB	SUB, SUBH, SUBK

A.3 Instruction Set Comparison Table

Table A–3 contains a comparison of the TMS320C1x, TMS320C2x, TMS320C2xx, and TMS320C5x instructions alphabetically.

Table A–3. Instruction Set Comparison

Syntax	1x	2x	2xx	5x	Description
ABS	√	√	√	√	**Absolute Value of Accumulator** If the contents of the accumulator are less than 0, replace the contents with the 2s complement of the contents. If the contents are ≥ 0, the accumulator is not affected.
ADCB				√	**Add ACCB to Accumulator With Carry** Add the contents of the ACCB and the value of the carry bit to the accumulator. If the result of the addition generates a carry from the accumulator's MSB, the carry bit is set to 1.
ADD *dma* [*, shift*]	√	√	√	√	**Add to Accumulator With Shift**
ADD {*ind*} [*, shift* [*, next ARP*]]	√	√	√	√	TMS320C1x and TMS320C2x devices: Add the contents of the addressed data-memory location to the accumulator; if a shift is specified, left shift the contents of the location before the add. During shifting, low-order bits are zero filled, and high-order bits are sign extended. TMS320C2xx and TMS320C5x devices: Add the contents of the addressed data-memory location or an immediate value to the accumulator; if a shift is specified, left shift the data before the add. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
ADD #*k*			√	√	
ADD #*lk* [*, shift2*]			√	√	
ADDB				√	**Add ACCB to Accumulator** Add the contents of the ACCB to the accumulator.
ADDC *dma*		√	√	√	**Add to Accumulator With Carry**
ADDC {*ind*} [*, next ARP*]		√	√	√	Add the contents of the addressed data-memory location and the carry bit to the accumulator.
ADDH *dma*	√	√	√	√	**Add High to Accumulator**
ADDH {*ind*} [*, next ARP*]	√	√	√	√	Add the contents of the addressed data-memory location to the 16 MSBs of the accumulator. The LSBs are not affected. If the result of the addition generates a carry, the carry bit is set to 1. TMS320C2x, TMS320C2xx, and TMS320C5x devices: If the result of the addition generates a carry from the accumulator's MSB, the carry bit is set to 1.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
ADDK *#k*		√	√	√	**Add to Accumulator Short Immediate** TMS320C1x devices: Add an 8-bit immediate value to the accumulator. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Add an 8-bit immediate value, right justified, to the accumulator with the result replacing the accumulator contents. The immediate value is treated as an 8-bit positive number; sign extension is suppressed.
ADDS *dma*	√	√	√	√	**Add to Accumulator With Sign Extension Suppressed** Add the contents of the addressed data-memory location to the accumulator. The value is treated as a 16-bit unsigned number; sign extension is suppressed.
ADDS {*ind*} [*, next ARP*]	√	√	√	√	
ADDT *dma*		√	√	√	**Add to Accumulator With Shift Specified by T Register** Left shift the contents of the addressed data-memory location by the value in the four LSBs of the T register; add the result to the accumulator. If a shift is specified, left shift the data before the add. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1. TMS320C2xx and TMS320C5x devices: If the result of the addition generates a carry from the accumulator's MSB, the carry bit is set to 1.
ADDT {*ind*} [*, next ARP*]		√	√	√	
ADLK *#lk* [*, shift*]		√	√	√	**Add to Accumulator Long Immediate With Shift** Add a 16-bit immediate value to the accumulator; if a shift is specified, left shift the value before the add. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
ADRK *#k*		√	√	√	**Add to Auxiliary Register Short Immediate** Add an 8-bit immediate value to the current auxiliary register.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
AND *dma* **AND** {*ind*} [*, next ARP*] **AND** *#lk* [*, shift*]	√ √	√ √	√ √ √	√ √ √	**AND With Accumulator** TMS320C1x and TMS320C2x devices: AND the contents of the addressed data-memory location with the 16 LSBs of the accumulator. The 16 MSBs of the accumulator are ANDed with 0s. TMS320C2xx and TMS320C5x devices: AND the contents of the addressed data-memory location or a 16-bit immediate value with the contents of the accumulator. The 16 MSBs of the accumulator are ANDed with 0s. If a shift is specified, left shift the constant before the AND. Low-order bits below and high-order bits above the shifted value are treated as 0s.
ANDB				√	**AND ACCB to Accumulator** AND the contents of the ACCB to the accumulator.
ANDK *#lk* [*, shift*]		√	√	√	**AND Immediate With Accumulator With Shift** AND a 16-bit immediate value with the contents of the accumulator; if a shift is specified, left shift the constant before the AND.
APAC	√	√	√	√	**Add P Register to Accumulator** Add the contents of the P register to the accumulator. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Before the add, left shift the contents of the P register as defined by the PM status bits.
APL [*#lk*] *,dma* **APL** [*#lk,*] {*ind*} [*, next ARP*]				√ √	**AND Data-Memory Value With DBMR or Long Constant** AND the data-memory value with the contents of the DBMR or a long constant. If a long constant is specified, it is ANDed with the contents of the data-memory location. The result is written back into the data-memory location previously holding the first operand. If the result is 0, the TC bit is set to 1; otherwise, the TC bit is cleared.
B *pma* **B** *pma* [, {*ind*} [*, next ARP*]]	√	√	√		**Branch Unconditionally** Branch to the specified program-memory address. TMS320C2x and TMS320C2xx devices: Modify the current AR and ARP as specified.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
B[*D*] *pma* [, {*ind*} [, *next ARP*]]				√	**Branch Unconditionally With Optional Delay** Modify the current auxiliary register and ARP as specified and pass control to the designated program-memory address. If you specify a delayed branch (BD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before branching.
BACC		√	√		**Branch to Address Specified by Accumulator** Branch to the location specified by the 16 LSBs of the accumulator.
BACC[*D*]				√	**Branch to Address Specified by Accumulator With Optional Delay** Branch to the location specified by the 16 LSBs of the accumulator. If you specify a delayed branch (BACCD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before branching.
BANZ *pma* **BANZ** *pma* [, {*ind*} [, *next ARP*]]	√	√	√		**Branch on Auxiliary Register Not Zero** If the contents of the nine LSBs of the current auxiliary register (TMS320C1x) or the contents of the entire current auxiliary register (TMS320C2x) are ≠ 0, branch to the specified program-memory address. TMS320C2x and TMS320C2xx devices: Modify the current AR and ARP (if specified) or decrement the current AR (default). TMS320C1x devices: Decrement the current AR.
BANZ[*D*] *pma* [, {*ind*} [, *next ARP*]]				√	**Branch on Auxiliary Register Not Zero With Optional Delay** If the contents of the current auxiliary register are ≠ 0, branch to the specified program-memory address. Modify the current AR and ARP as specified, or decrement the current AR. If you specify a delayed branch (BANZD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before branching.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
BBNZ *pma* [, {*ind*} [, *next ARP*]]		√	√	√	**Branch on Bit ≠ Zero** If the TC bit = 1, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: If the –p porting switch is used, modify the current AR and ARP as specified.
BBZ *pma* [, {*ind*} [, *next ARP*]] **BBZ** *pma*		√	√	√ √	**Branch on Bit = Zero** If the TC bit = 0, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BC *pma* [, {*ind*} [, *next ARP*]] **BC** *pma*		√	 √	√ √	**Branch on Carry** If the C bit = 1, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BCND *pma*, $cond_1$ [, $cond_2$] [, ...]			√		**Branch Conditionally** Branch to the program-memory address if the specified conditions are met. Not all combinations of conditions are meaningful.
BCND[*D*] *pma*, $cond_1$ [, $cond_2$] [, ...]				√	**Branch Conditionally With Optional Delay** Branch to the program-memory address if the specified conditions are met. Not all combinations of conditions are meaningful. If you specify a delayed branch (BCNDD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before branching.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
BGEZ *pma*	√		√	√	**Branch if Accumulator ≥ Zero**
BGEZ *pma* [, {*ind*} [, *next ARP*]]		√		√	If the contents of the accumulator ≥ 0, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BGZ *pma*	√		√	√	**Branch if Accumulator > Zero**
BGZ *pma* [, {*ind*} [, *next ARP*]]		√		√	If the contents of the accumulator are > 0, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BIOZ *pma*	√		√	√	**Branch on I/O Status = Zero**
BIOZ *pma* [, {*ind*} [, *next ARP*]]		√		√	If the $\overline{\text{BIO}}$ pin is low, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BIT *dma, bit code*		√	√	√	**Test Bit**
BIT {*ind*}, *bit code* [, *next ARP*]		√	√	√	Copy the specified bit of the data-memory value to the TC bit in ST1.
BITT *dma*		√	√	√	**Test Bit Specified by T Register**
BITT {*ind*} [, *next ARP*]		√	√	√	TMS320C2x and TMS320C2xx devices: Copy the specified bit of the data-memory value to the TC bit in ST1. The four LSBs of the T register specify which bit is copied. TMS320C5x devices: Copy the specified bit of the data-memory value to the TC bit in ST1. The four LSBs of the TREG2 specify which bit is copied.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
BLDD *#lk, dma*			√	√	**Block Move From Data Memory to Data Memory**
BLDD *#lk,* {*ind*} [*, next ARP*]			√	√	Copy a block of data memory into data memory. The block of data memory is pointed to by *src*, and the destination block of data memory is pointed to by *dst*.
BLDD *dma, #lk*			√	√	
BLDD {*ind*}*, #lk* [*, next ARP*]			√	√	TMS320C2xx devices: The word of the source and/or the destination space can be pointed to with a long immediate value or a data-memory address. You can use the RPT instruction with BLDD to move consecutive words, pointed to indirectly in data memory, to contiguous program-memory spaces. The number of words to be moved is 1 greater than the number contained in the RPTC at the beginning of the instruction.
BLDD BMAR, *dma*				√	
BLDD BMAR, {*ind*} [*, next ARP*]				√	
BLDD *dma* BMAR				√	
BLDD {*ind*}*,* BMAR [*, next ARP*]				√	
					TMS320C5x devices: The word of the source and/or the destination space can be pointed to with a long immediate value, the contents of the BMAR, or a data-memory address. You can use the RPT instruction with BLDD to move consecutive words, pointed to indirectly in data memory, to a contiguous program-memory space. The number of words to be moved is 1 greater than the number contained in the RPTC at the beginning of the instruction.
BLDP *dma*				√	**Block Move From Data Memory to Program Memory**
BLDP {*ind*} [*, next ARP*]				√	Copy a block of data memory into program memory pointed to by the BMAR. You can use the RPT instruction with BLDP to move consecutive words, indirectly pointed to in data memory, to a contiguous program-memory space pointed to by the BMAR.
BLEZ *pma*	√		√	√	**Branch if Accumulator ≤ Zero**
BLEZ *pma* [*,* {*ind*} [*, next ARP*]]		√	√	√	If the contents of the accumulator are ≤ 0, branch to the specified program-memory address.
					TMS320C2x devices: Modify the current AR and ARP as specified.
					TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
BLKD *dma1, dma2* **BLKD** *dma1,* {*ind*} [*, next ARP*]		√	√	√	**Block Move From Data Memory to Data Memory** Move a block of words from one location in data memory to another location in data memory. Modify the current AR and ARP as specified. RPT or RPTK must be used with BLKD, in the indirect addressing mode, if more than one word is to be moved. The number of words to be moved is 1 greater than the number contained in RPTC at the beginning of the instruction.
BLKP *pma, dma* **BLKP** *pma,* {*ind*} [*, next ARP*]		√	√	√	**Block Move From Program Memory to Data Memory** Move a block of words from a location in program memory to a location in data memory. Modify the current AR and ARP as specified. RPT or RPTK must be used with BLKD, in the indirect addressing mode, if more than one word is to be moved. The number of words to be moved is 1 greater than the number contained in RPTC at the beginning of the instruction.
BLPD #*pma, dma* **BLPD** #*pma,* {*ind*} [*, next ARP*] **BLPD** BMAR, *dma* **BLPD** BMAR, {*ind*} [*, next ARP*]			√ √	√ √ √ √	**Block Move From Program Memory to Data Memory** Copy a block of program memory into data memory. The block of program memory is pointed to by *src*, and the destination block of data memory is pointed to by *dst*. TMS320C2xx devices: The word of the source space can be pointed to with a long immediate value. You can use the RPT instruction with BLPD to move consecutive words that are pointed at indirectly in data memory to a contiguous program-memory space. TMS320C5x devices: The word of the source space can be pointed to with a long immediate value or the contents of the BMAR. You can use the RPT instruction with BLPD to move consecutive words that are pointed at indirectly in data memory to contiguous program-memory spaces.
BLZ *pma* **BLZ** *pma* [*,* {*ind*} [*, next ARP*]]	√	√	√ √	√	**Branch if Accumulator < Zero** If the contents of the accumulator are < 0, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
BNC *pma* [, {*ind*} [, *next ARP*]]		√	√	√	**Branch on No Carry** If the C bit = 0, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BNV *pma* [, {*ind*} [, *next ARP*]]		√	√	√	**Branch if No Overflow** If the OV flag is clear, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BNZ *pma*	√				**Branch if Accumulator ≠ Zero**
BNZ *pma* [, {*ind*} [, *next ARP*]]		√	√	√	If the contents of the accumulator ≠ 0, branch to the specified program-memory address. TMS320C2x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified when the –p porting switch is used.
BSAR [*shift*]				√	**Barrel Shift** In a single cycle, execute a 1- to 16-bit right arithmetic barrel shift of the accumulator. The sign extension is determined by the sign-extension mode bit in ST1.
BV *pma*	√				**Branch on Overflow**
BV *pma* [, {*ind*} [, *next ARP*]]		√	√	√	If the OV flag is set, branch to the specified program-memory address and clear the OV flag. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: To modify the AR and ARP, use the –p porting switch.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
BZ *pma* **BZ** *pma* [, {*ind*} [, *next ARP*]]	√	√	√	√	**Branch if Accumulator = Zero** If the contents of the accumulator = 0, branch to the specified program-memory address. TMS320C2x, TMS320C2xx and TMS320C5x devices: Modify the current AR and ARP as specified. TMS320C2xx and TMS320C5x devices: To modify the AR and ARP, use the –p porting switch.
CALA	√	√	√		**Call Subroutine Indirect** The contents of the accumulator specify the address of a subroutine. Increment the PC, push the PC onto the stack, then load the 12 (TMS320C1x) or 16 (TMS320C2x/C2xx) LSBs of the accumulator into the PC.
CALA[*D*]				√	**Call Subroutine Indirect With Optional Delay** The contents of the accumulator specify the address of a subroutine. Increment the PC and push it onto the stack; then load the 16 LSBs of the accumulator into the PC. If you specify a delayed branch (CALAD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before the call.
CALL *pma* **CALL** *pma* [,{*ind*} [, *next ARP*]]	√	√	√		**Call Subroutine** The contents of the addressed program-memory location specify the address of a subroutine. Increment the PC by 2, push the PC onto the stack, then load the specified program-memory address into the PC. TMS320C2x and TMS320C2xx devices: Modify the current AR and ARP as specified.
CALL[*D*] *pma* [, {*ind*} [, *next ARP*]]				√	**Call Unconditionally With Optional Delay** The contents of the addressed program-memory location specify the address of a subroutine. Increment the PC and push the PC onto the stack; then load the specified program-memory address (symbolic or numeric) into the PC. Modify the current AR and ARP as specified. If you specify a delayed branch (CALLD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before the call.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
CC *pma, cond$_1$ [, cond$_2$] [, ...]*			√		**Call Conditionally** If the specified conditions are met, control is passed to the pma. Not all combinations of conditions are meaningful.
CC[*D*] *pma, cond$_1$ [, cond$_2$] [, ...]*				√	**Call Conditionally With Optional Delay** If the specified conditions are met, control is passed to the pma. Not all combinations of conditions are meaningful. If you specify a delayed branch (CCD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before the call.
CLRC *control bit*			√	√	**Clear Control Bit** Set the specified control bit to a logic 0. Maskable interrupts are enabled immediately after the CLRC instruction executes.
CMPL		√	√	√	**Complement Accumulator** Complement the contents of the accumulator (1s complement).
CMPR *CM*		√	√	√	**Compare Auxiliary Register With AR0** Compare the contents of the current auxiliary register to AR0, based on the following cases: If CM = 00, test whether AR(ARP) = AR0. If CM = 01, test whether AR(ARP) < AR0. If CM = 10, test whether AR(ARP) > AR0. If CM = 11, test whether AR(ARP) ≠ AR0. If the result is true, load a 1 into the TC status bit; otherwise, load a 0 into the TC bit. The comparison does not affect the tested registers. TMS320C5x devices: Compare the contents of the auxiliary register with the ARCR.
CNFD		√	√	√	**Configure Block as Data Memory** Configure on-chip RAM block B0 as data memory. Block B0 is mapped into data-memory locations 512h–767h. TMS320C5x devices: Block B0 is mapped into data-memory locations 512h–1023h.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
CNFP		√	√	√	**Configure Block as Program Memory** Configure on-chip RAM block B0 as program memory. Block B0 is mapped into program-memory locations 65280h–65535h. TMS320C5x devices: Block B0 is mapped into data-memory locations 65024h–65535h.
CONF *2-bit constant*		√			**Configure Block as Program Memory** Configure on-chip RAM block B0/B1/B2/B3 as program memory. For information on the memory mapping of B0/B1/B2/B3, see the *TMS320C2x User's Guide*.
CPL [*#lk,*] *dma* **CPL** [*#lk,*] {*ind*} [*, next ARP*]				√ √	**Compare DBMR or Immediate With Data Value** Compare two quantities: If the two quantities are equal, set the TC bit to 1; otherwise, clear the TC bit.
CRGT				√	**Test for ACC > ACCB** Compare the contents of the ACC with the contents of the ACCB, then load the larger signed value into both registers and modify the carry bit according to the comparison result. If the contents of ACC are greater than or equal to the contents of ACCB, set the carry bit to 1.
CRLT				√	**Test for ACC < ACCB** Compare the contents of the ACC with the contents of the ACCB, then load the smaller signed value into both registers and modify the carry bit according to the comparison result. If the contents of ACC are less than the contents of ACCB, clear the carry bit.
DINT	√	√	√	√	**Disable Interrupts** Disable all interrupts; set the INTM to 1. Maskable interrupts are disabled immediately after the DINT instruction executes. DINT does not disable the unmaskable interrupt $\overline{RS}$; DINT does not affect the IMR.
DMOV *dma* **DMOV** {*ind*} [*, next ARP*]	√ √	√ √	√ √	√ √	**Data Move in Data Memory** Copy the contents of the addressed data-memory location into the next higher address. DMOV moves data only within on-chip RAM blocks. TMS320C2x, TMS320C2xx, and TMS320C5x devices: The on-chip RAM blocks are B0 (when configured as data memory), B1, and B2.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
EINT	√	√	√	√	**Enable Interrupts** Enable all interrupts; clear the INTM to 0. Maskable interrupts are enabled immediately after the EINT instruction executes.
EXAR				√	**Exchange ACCB With ACC** Exchange the contents of the ACC with the contents of the ACCB.
FORT *1-bit constant*		√			**Format Serial Port Registers** Load the FO with a 0 or a 1. If FO = 0, the registers are configured to receive/transmit 16-bit words. If FO = 1, the registers are configured to receive/transmit 8-bit bytes.
IDLE		√	√	√	**Idle Until Interrupt** Forces an executing program to halt execution and wait until it receives a reset or an interrupt. The device remains in an idle state until it is interrupted.
IDLE2				√	**Idle Until Interrupt — Low-Power Mode** Removes the functional clock input from the internal device; this allows for an extremely low-power mode. The IDLE2 instruction forces an executing program to halt execution and wait until it receives a reset or unmasked interrupt.
IN *dma, PA*	√	√	√	√	**Input Data From Port**
IN {*ind*}*, PA* [*, next ARP*]	√	√	√	√	Read a 16-bit value from one of the external I/O ports into the addressed data-memory location. TMS320C1x devices: This is a 2-cycle instruction. During the first cycle, the port address is sent to address lines A2/PA2–A0/PA0; $\overline{DEN}$ goes low, strobing in the data that the addressed peripheral places on data bus D15–D0. TMS320C2x devices: The $\overline{IS}$ line goes low to indicate an I/O access, and the $\overline{STRB}$, R/$\overline{W}$, and READY timings are the same as for an external data-memory read. TMS320C2xx and TMS320C5x devices: The $\overline{IS}$ line goes low to indicate an I/O access, and the $\overline{STRB}$, $\overline{RD}$, and READY timings are the same as for an external data-memory read.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
INTR *K*			√	√	**Soft Interrupt** Transfer program control to the program-memory address specified by *K* (an integer from 0 to 31). This instruction allows you to use your software to execute any interrupt service routine. The interrupt vector locations are spaced apart by two addresses (0h, 2h, 4h, ... , 3Eh), allowing a 2-word branch instruction to be placed at each location.
LAC *dma* [*, shift*]	√	√	√	√	**Load Accumulator With Shift**
LAC {*ind*} [*, shift* [*, next ARP*]]	√	√	√	√	Load the contents of the addressed data-memory location into the accumulator. If a shift is specified, left shift the value before loading it into the accumulator. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
LACB				√	**Load Accumulator With ACCB** Load the contents of the accumulator buffer into the accumulator.
LACC *dma* [*, shift*$_1$]		√	√	√	**Load Accumulator With Shift**
LACC {*ind*} [*, shift*$_1$ [*, next ARP*]]		√	√	√	Load the contents of the addressed data-memory location or the 16-bit constant into the accumulator. If a shift is specified, left shift the value before loading it into the accumulator. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
LACC #*lk* [*, shift*$_2$]		√	√	√	
LACK *8-bit constant*	√	√	√	√	**Load Accumulator Immediate Short** Load an 8-bit constant into the accumulator. The 24 MSBs of the accumulator are zeroed.
LACL *dma*			√	√	**Load Low Accumulator and Clear High Accumulator**
LACL {*ind*} [*, next ARP*]			√	√	Load the contents of the addressed data-memory location or zero-extended 8-bit constant into the 16 LSBs of the accumulator. The MSBs of the accumulator are zeroed. The data is treated as a 16-bit unsigned number. TMS320C2xx: A constant of 0 clears the contents of the accumulator to 0 with no sign extension.
LACL #*k*			√	√	

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
LACT *dma* **LACT** {*ind*} [*, next ARP*]		√ √	√ √	√ √	**Load Accumulator With Shift Specified by T Register** Left shift the contents of the addressed data-memory location by the value specified in the four LSBs of the T register; load the result into the accumulator. If a shift is specified, left shift the value before loading it into the accumulator. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
LALK *#lk* [*, shift*]		√	√	√	**Load Accumulator Long Immediate With Shift** Load a 16-bit immediate value into the accumulator. If a shift is specified, left shift the constant before loading it into the accumulator. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
LAMM *dma* **LAMM** {*ind*} [*, next ARP*]				√ √	**Load Accumulator With Memory-Mapped Register** Load the contents of the addressed memory-mapped register into the low word of the accumulator. The nine MSBs of the data-memory address are cleared, regardless of the current value of DP or the nine MSBs of AR (ARP).
LAR *AR, dma* **LAR** *AR,* {*ind*} [*, next ARP*] **LAR** *AR, #k* **LAR** *AR, #lk*	√ √	√ √	√ √ √ √	√ √ √ √	**Load Auxiliary Register** TMS320C1x and TMS320C2x devices: Load the contents of the addressed data-memory location into the designated auxiliary register. TMS320C25, TMS320C2xx, and TMS320C5x devices: Load the contents of the addressed data-memory location or an 8-bit or 16-bit immediate value into the designated auxiliary register.
LARK *AR, 8-bit constant*	√	√	√	√	**Load Auxiliary Register Immediate Short** Load an 8-bit positive constant into the designated auxiliary register.
LARP *1-bit constant* **LARP** *3-bit constant*	√	√	√	√	**Load Auxiliary Register Pointer** TMS320C1x devices: Load a 1-bit constant into the auxiliary register pointer (specifying AR0 or AR1). TMS320C2x, TMS320C2xx, and TMS320C5x devices: Load a 3-bit constant into the auxiliary register pointer (specifying AR0–AR7).

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
LDP *dma*	√	√	√	√	**Load Data-Memory Page Pointer**
LDP {*ind*} [*, next ARP*]	√	√	√	√	TMS320C1x devices: Load the LSB of the contents of the addressed data-memory location into the DP register. All high-order bits are ignored. DP = 0 defines page 0 (words 0–127), and DP = 1 defines page 1 (words 128–143/255).
LDP *#k*			√	√	TMS320C2x, TMS320C2xx, and TMS320C5x devices: Load the nine LSBs of the addressed data-memory location or a 9-bit immediate value into the DP register. The DP and 7-bit data-memory address are concatenated to form 16-bit data-memory addresses.
LDPK *1-bit constant*	√				**Load Data-Memory Page Pointer Immediate**
LDPK *9-bit constant*		√	√	√	TMS320C1x devices: Load a 1-bit immediate value into the DP register. DP = 0 defines page 0 (words 0–127), and DP = 1 defines page 1 (words 128–143/255).
					TMS320C2x, TMS320C2xx, and TMS320C5x devices: Load a 9-bit immediate into the DP register. The DP and 7-bit data-memory address are concatenated to form 16-bit data-memory addresses. DP ≥ 8 specifies external data memory. DP = 4 through 7 specifies on-chip RAM blocks B0 or B1. Block B2 is located in the upper 32 words of page 0.
LMMR *dma, #lk*				√	**Load Memory-Mapped Register**
LMMR {*ind*}, *#lk* [*, next ARP*]				√	Load the contents of the memory-mapped register pointed at by the seven LSBs of the direct or indirect data-memory value into the long immediate addressed data-memory location. The nine MSBs of the data-memory address are cleared, regardless of the current value of DP or the nine MSBs of AR (ARP).
LPH *dma*		√	√	√	**Load High P Register**
LPH {*ind*} [*, next ARP*]		√	√	√	Load the contents of the addressed data-memory location into the 16 MSBs of the P register; the LSBs are not affected.
LRLK *AR, lk*		√	√	√	**Load Auxiliary Register Long Immediate**
					Load a 16-bit immediate value into the designated auxiliary register.
LST *dma*	√	√	√	√	**Load Status Register**
LST {*ind*} [*, next ARP*]	√	√	√	√	Load the contents of the addressed data-memory location into the ST (TMS320C1x) or into ST0 (TMS320C2x/2xx/5x).

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
LST #*n, dma*		√	√	√	**Load Status Register n**
LST #*n,* {*ind*} [*, next ARP*]		√	√	√	Load the contents of the addressed data-memory location into ST*n*.
LST1 *dma*		√	√	√	**Load ST1**
LST1 {*ind*} [*, next ARP*]		√	√	√	Load the contents of the addressed data-memory location into ST1.
LT *dma*	√	√	√	√	**Load T Register**
LT {*ind*} [*, next ARP*]	√	√	√	√	Load the contents of the addressed data-memory location into the T register (TMS320C1x/2x/2xx) or TREG0 (TMS320C5x).
LTA *dma*	√	√	√	√	**Load T Register and Accumulate Previous Product**
LTA {*ind*} [*, next ARP*]	√	√	√	√	Load the contents of the addressed data-memory location into T register (TMS320C1x/2x/2xx) or TREG0 (TMS320C5x) and add the contents of the P register to the accumulator. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Before the add, shift the contents of the P register as specified by the PM status bits.
LTD *dma*	√	√	√	√	**Load T Register, Accumulate Previous Product, and Move Data**
LTD {*ind*} [*, next ARP*]	√	√	√	√	Load the contents of the addressed data-memory location into the T register (TMS320C1x/2x/2xx) or TREG0 (TMS320C5x), add the contents of the P register to the accumulator, and copy the contents of the specified location into the next higher address (both data-memory locations must reside in on-chip data RAM). TMS320C2x, TMS320C2xx, and TMS320C5x devices: Before the add, shift the contents of the P register as specified by the PM status bits.
LTP *dma*		√	√	√	**Load T Register, Store P Register in Accumulator**
LTP {*ind*} [*, next ARP*]		√	√	√	Load the contents of the addressed data-memory location into the T register (TMS320C1x/2x/2xx) or TREG0 (TMS320C5x). Store the contents of the product register into the accumulator.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
LTS *dma* **LTS** {*ind*} [*, next ARP*]		√	√	√	**Load T Register, Subtract Previous Product** Load the contents of the addressed data-memory location into the T register (TMS320C1x/2x/2xx) or TREG0 (TMS320C5x). Shift the contents of the P register as specified by the PM status bits, and subtract the result from the accumulator.
MAC *pma, dma* **MAC** *pma,* {*ind*} [*, next ARP*]		√	√	√	**Multiply and Accumulate** Multiply a data-memory value by a program-memory value and add the previous product (shifted as specified by the PM status bits) to the accumulator.
MACD *dma, pma* **MACD** *pma,* {*ind*} [*, next ARP*]		√	√	√	**Multiply and Accumulate With Data Move** Multiply a data-memory value by a program-memory value and add the previous product (shifted as specified by the PM status bits) to the accumulator. If the data-memory address is in on-chip RAM block B0, B1, or B2, copy the contents of the address to the next higher address.
MADD *dma* **MADD** {*ind*} [*, next ARP*]				√ √	**Multiply and Accumulate With Data Move and Dynamic Addressing** Multiply a data-memory value by a program-memory value and add the previous product (shifted as defined by the PM status bits) into the accumulator. The program-memory address is contained in the BMAR; this allows for dynamic addressing of coefficient tables. MADD functions the same as MADS, with the addition of data move for on-chip RAM blocks.
MADS *dma* **MADS** {*ind*} [*, next ARP*]				√ √	**Multiply and Accumulate With Dynamic Addressing** Multiply a data-memory value by a program-memory value and add the previous product (shifted as defined by the PM status bits) into the accumulator. The program-memory address is contained in the BMAR; this allows for dynamic addressing of coefficient tables.
MAR *dma* **MAR** {*ind*} [*, next ARP*]	√ √	√ √	√ √	√ √	**Modify Auxiliary Register** Modify the current AR or ARP as specified. MAR acts as NOP in indirect addressing mode.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
MPY *dma*	√	√	√	√	**Multiply**
MPY {*ind*} [*, next ARP*]	√	√	√	√	TMS320C1x and TMS320C2x devices: Multiply the contents of the T register by the contents of the addressed data-memory location; place the result in the P register.
MPY #*k*			√	√	
MPY #*lk*			√	√	TMS320C2xx and TMS320C5x devices: Multiply the contents of the T register (TMS320C2xx) or TREG0 (TMS320C5x) by the contents of the addressed data-memory location or a 13-bit or 16-bit immediate value; place the result in the P register.
MPYA *dma*		√	√	√	**Multiply and Accumulate Previous Product**
MPYA {*ind*} [*, next ARP*]		√	√	√	Multiply the contents of the T register (TMS320C2x/2xx) or TREG0 (TMS320C5x) by the contents of the addressed data-memory location; place the result in the P register. Add the previous product (shifted as specified by the PM status bits) to the accumulator.
MPYK *13-bit constant*	√	√	√	√	**Multiply Immediate**
					Multiply the contents of the T register (TMS320C2x/2xx) or TREG0 (TMS320C5x) by a signed 13-bit constant; place the result in the P register.
MPYS *dma*		√	√	√	**Multiply and Subtract Previous Product**
MPYS {*ind*} [*, next ARP*]		√	√	√	Multiply the contents of the T register (TMS320C2x/2xx) or TREG0 (TMS320C5x) by the contents of the addressed data-memory location; place the result in the P register. Subtract the previous product (shifted as specified by the PM status bits) from the accumulator.
MPYU *dma*		√	√	√	**Multiply Unsigned**
MPYU {*ind*} [*, next ARP*]		√	√	√	Multiply the unsigned contents of the T register (TMS320C2x/2xx) or TREG0 (TMS320C5x) by the unsigned contents of the addressed data-memory location; place the result in the P register.
NEG		√	√	√	**Negate Accumulator**
					Negate (2s complement) the contents of the accumulator.
NMI			√	√	**Nonmaskable Interrupt**
					Force the program counter to the nonmaskable interrupt vector location 24h. NMI has the same effect as a hardware nonmaskable interrupt.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
NOP	√	√	√	√	**No Operation** Perform no operation.
NORM		√	√	√	**Normalize Contents of Accumulator**
NORM {*ind*}		√	√	√	Normalize a signed number in the accumulator.
OPL [#*lk*,] *dma*				√	**OR With DBMR or Long Immediate**
OPL [#*lk*,] {*ind*} [*, next ARP*]				√	If a long immediate is specified, OR it with the value at the specified data-memory location; otherwise, the second operand of the OR operation is the contents of the DBMR. The result is written back into the data-memory location previously holding the first operand.
OR *dma*	√	√	√	√	**OR With Accumulator**
OR {*ind*} [*, next ARP*]	√	√	√	√	TMS320C1x and TMS320C2x devices: OR the 16 LSBs of the accumulator with the contents of the addressed data-memory location. The 16 MSBs of the accumulator are ORed with 0s.
OR #*lk* [*, shift*]			√	√	TMS320C2xx and TMS320C5x devices: OR the 16 LSBs of the accumulator or a 16-bit immediate value with the contents of the addressed data-memory location. If a shift is specified, left-shift before ORing. Low-order bits below and high-order bits above the shifted value are treated as 0s.
ORB				√	**OR ACCB With Accumulator** OR the contents of the ACCB with the contents of the accumulator. ORB places the result in the accumulator.
ORK #*lk* [*, shift*]		√	√	√	**OR Immediate With Accumulator with Shift** OR a 16-bit immediate value with the contents of the accumulator. If a shift is specified, left-shift the constant before ORing. Low-order bits below and high-order bits above the shifted value are treated as 0s.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
OUT *dma, PA* **OUT** {*ind*}*, PA* [*, next ARP*]	√	√	√	√	**Output Data to Port** Write a 16-bit value from a data-memory location to the specified I/O port. TMS320C1x devices: The first cycle of this instruction places the port address onto address lines A2/PA2–A0/PA0. During the same cycle, $\overline{WE}$ goes low and the data word is placed on the data bus D15–D0. TMS320C2x, TMS320C2xx, and TMS320C5x devices: The $\overline{IS}$ line goes low to indicate an I/O access; the $\overline{STRB}$, R/$\overline{W}$, and READY timings are the same as for an external data-memory write.
PAC	√	√	√	√	**Load Accumulator With P Register** Load the contents of the P register into the accumulator. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Before the load, shift the P register as specified by the PM status bits.
POP	√	√	√	√	**Pop Top of Stack to Low Accumulator** Copy the contents of the top of the stack into the 12 (TMS320C1x) or 16 (TMS320C2x/2xx/5x) LSBs of the accumulator and then pop the stack one level. The MSBs of the accumulator are zeroed.
POPD *dma* **POPD** {*ind*} [*, next ARP*]		√	√	√	**Pop Top of Stack to Data Memory** Transfer the value on the top of the stack into the addressed data-memory location and then pop the stack one level.
PSHD *dma* **PSHD** {*ind*} [*, next ARP*]		√	√	√	**Push Data-Memory Value Onto Stack** Copy the addressed data-memory location onto the top of the stack. The stack is pushed down one level before the value is copied.
PUSH	√	√	√	√	**Push Low Accumulator Onto Stack** Copy the contents of the 12 (TMS320C1x) or 16 (TMS320C2x/2xx/5x) LSBs of the accumulator onto the top of the hardware stack. The stack is pushed down one level before the value is copied.
RC		√	√	√	**Reset Carry Bit** Reset the C status bit to 0.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
RET	√	√	√		**Return From Subroutine** Copy the contents of the top of the stack into the PC and pop the stack one level.
RET[*D*]				√	**Return From Subroutine With Optional Delay** Copy the contents of the top of the stack into the PC and pop the stack one level. If you specify a delayed branch (RETD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before the return.
RETC $cond_1$ [, $cond_2$] [, ...]			√		**Return Conditionally** If the specified conditions are met, RETC performs a standard return. Not all combinations of conditions are meaningful.
RETC[*D*] $cond_1$ [, $cond_2$] [, ...]				√	**Return Conditionally With Optional Delay** If the specified conditions are met, RETC performs a standard return. Not all combinations of conditions are meaningful. If you specify a delayed branch (RETCD), the next two instruction words (two 1-word instructions or one 2-word instruction) are fetched and executed before the return.
RETE				√	**Enable Interrupts and Return From Interrupt** Copy the contents of the top of the stack into the PC and pop the stack one level. RETE automatically clears the global interrupt enable bit and pops the shadow registers (stored when the interrupt was taken) back into their corresponding strategic registers. The following registers are shadowed: ACC, ACCB, PREG, ST0, ST1, PMST, ARCR, INDX, TREG0, TREG1, TREG2.
RETI				√	**Return From Interrupt** Copy the contents of the top of the stack into the PC and pop the stack one level. RETI also pops the values in the shadow registers (stored when the interrupt was taken) back into their corresponding strategic registers. The following registers are shadowed: ACC, ACCB, PREG, ST0, ST1, PMST, ARCR, INDX, TREG0, TREG1, TREG2.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
RFSM		√			**Reset Serial Port Frame Synchronization Mode** Reset the FSM status bit to 0.
RHM		√		√	**Reset Hold Mode** Reset the HM status bit to 0.
ROL		√	√	√	**Rotate Accumulator Left** Rotate the accumulator left one bit.
ROLB				√	**Rotate ACCB and Accumulator Left** Rotate the ACCB and the accumulator left by one bit; this results in a 65-bit rotation.
ROR		√	√	√	**Rotate Accumulator Right** Rotate the accumulator right one bit.
RORB				√	**Rotate ACCB and Accumulator Right** Rotate the ACCB and the accumulator right one bit; this results in a 65-bit rotation.
ROVM	√	√	√	√	**Reset Overflow Mode** Reset the OVM status bit to 0; this disables overflow mode.
RPT *dma* **RPT** {*ind*} [*, next ARP*] **RPT** #*k* **RPT** #*lk*		√ √	√ √ √ √	√ √ √ √	**Repeat Next Instruction** TMS320C2x devices: Load the eight LSBs of the addressed value into the RPTC; the instruction following RPT is executed the number of times indicated by RPTC + 1. TMS320C2xx and TMS320C5x devices: Load the eight LSBs of the addressed value or an 8-bit or 16-bit immediate value into the RPTC; the instruction following RPT is repeated *n* times, where *n* is RPTC+1.
RPTB *pma*				√	**Repeat Block** RPTB repeats a block of instructions the number of times specified by the memory-mapped BRCR without any penalty for looping. The BRCR must be loaded before RPTB is executed.
RPTK #*k*		√	√	√	**Repeat Instruction as Specified by Immediate Value** Load the 8-bit immediate value into the RPTC; the instruction following RPTK is executed the number of times indicated by RPTC + 1.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
RPTZ *#lk*				√	**Repeat Preceded by Clearing the Accumulator and P Register** Clear the accumulator and product register and repeat the instruction following RPTZ *n* times, where *n* = *lk* +1.
RSXM		√	√	√	**Reset Sign-Extension Mode** Reset the SXM status bit to 0; this suppresses sign extension on shifted data values for the following arithmetic instructions: ADD, ADDT, ADLK, LAC, LACT, LALK, SBLK, SUB, and SUBT.
RTC		√	√	√	**Reset Test/Control Flag** Reset the TC status bit to 0.
RTXM		√			**Reset Serial Port Transmit Mode** Reset the TXM status bit to 0; this configures the serial port transmit section in a mode where it is controlled by an FSX.
RXF		√	√	√	**Reset External Flag** Reset XF pin and the XF status bit to 0.
SACB				√	**Store Accumulator in ACCB** Copy the contents of the accumulator into the ACCB.
SACH *dma* [*, shift*]	√	√	√	√	**Store High Accumulator With Shift**
SACH {*ind*} [*, shift* [*, next ARP*]]	√	√	√	√	Copy the contents of the accumulator into a shifter. Shift the entire contents by zero, one, or four bits (TMS320C1x) or from zero to seven bits (TMS320C2x/2xx/5x), and then copy the 16 MSBs of the shifted value into the addressed data-memory location. The accumulator is not affected.
SACL *dma*	√				**Store Low Accumulator With Shift**
SACL *dma* [*, shift*]	√	√	√	√	TMS320C1x devices: Store the 16 LSBs of the accumulator into the addressed data-memory location. A shift value of 0 must be specified if the ARP is to be changed. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Store the 16 LSBs of the accumulator into the addressed data-memory location. If a shift is specified, shift the contents of the accumulator before storing. Shift values are zero, one, or four bits (TMS320C20) or from zero to seven bits (TMS320C2x/2xx/5x).
SACL {*ind*} [*, shift* [*, next ARP*]]		√	√	√	

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
SAMM *dma* **SAMM** {*ind*} [*, next ARP*]				√ √	**Store Accumulator in Memory-Mapped Register** Store the low word of the accumulator in the addressed memory-mapped register. The upper nine bits of the data address are cleared, regardless of the current value of DP or the nine MSBs of AR (ARP).
SAR *AR, dma* **SAR** *AR,* {*ind*} [*, next ARP*]	√ √	√ √	√ √	√ √	**Store Auxiliary Register** Store the contents of the specified auxiliary register in the addressed data-memory location.
SATH				√	**Barrel-Shift Accumulator as Specified by T Register 1** If bit 4 of TREG1 is a 1, barrel-shift the accumulator right by 16 bits; otherwise, the accumulator is unaffected.
SATL				√	**Barrel-Shift Low Accumulator as Specified by T Register 1** Barrel-shift the accumulator right by the value specified in the four LSBs of TREG1.
SBB				√	**Subtract ACCB From Accumulator** Subtract the contents of the ACCB from the accumulator. The result is stored in the accumulator; the accumulator buffer is not affected.
SBBB				√	**Subtract ACCB From Accumulator With Borrow** Subtract the contents of the ACCB and the logical inversion of the carry bit from the accumulator. The result is stored in the accumulator; the accumulator buffer is not affected. Clear the carry bit if the result generates a borrow.
SBLK #*lk* [*, shift*]		√	√	√	**Subtract From Accumulator Long Immediate With Shift** Subtract the immediate value from the accumulator. If a shift is specified, left shift the value before subtracting. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
SBRK #*k*		√	√	√	**Subtract From Auxiliary Register Short Immediate** Subtract the 8-bit immediate value from the designated auxiliary register.
SC		√	√	√	**Set Carry Bit** Set the C status bit to 1.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
SETC *control bit*			√	√	**Set Control Bit** Set the specified control bit to a logic 1. Maskable interrupts are disabled immediately after the SETC instruction executes.
SFL		√	√	√	**Shift Accumulator Left** Shift the contents of the accumulator left one bit.
SFLB				√	**Shift ACCB and Accumulator Left** Shift the concatenation of the accumulator and the ACCB left one bit. The LSB of the ACCB is cleared to 0, and the MSB of the ACCB is shifted into the carry bit.
SFR		√	√	√	**Shift Accumulator Right** Shift the contents of the accumulator right one bit. If SXM = 1, SFR produces an arithmetic right shift. If SXM = 0, SFR produces a logic right shift.
SFRB				√	**Shift ACCB and Accumulator Right** Shift the concatenation of the accumulator and the ACCB right 1 bit. The LSB of the ACCB is shifted into the carry bit. If SXM = 1, SFRB produces an arithmetic right shift. If SXM = 0, SFRB produces a logic right shift.
SFSM		√			**Set Serial Port Frame Synchronization Mode** Set the FSM status bit to 1.
SHM		√		√	**Set Hold Mode** Set the HM status bit to 1.
SMMR *dma, #lk* **SMMR** {*ind*}, *#lk* [, *next ARP*]				√ √	**Store Memory-Mapped Register** Store the memory-mapped register value, pointed at by the seven LSBs of the data-memory address, into the long immediate addressed data-memory location. The nine MSBs of the data-memory address of the memory-mapped register are cleared, regardless of the current value of DP or the upper nine bits of AR(ARP).
SOVM	√	√	√	√	**Set Overflow Mode** Set the OVM status bit to 1; this enables overflow mode. (The ROVM instruction clears OVM.)

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
SPAC	√	√	√	√	**Subtract P Register From Accumulator** Subtract the contents of the P register from the contents of the accumulator. TMS320C2x, TMS320C2xx, and TMS320C5x devices: Before the subtraction, shift the contents of the P register as specified by the PM status bits.
SPH *dma*		√	√	√	**Store High P Register**
SPH {*ind*} [*, next ARP*]		√	√	√	Store the high-order bits of the P register (shifted as specified by the PM status bits) at the addressed data-memory location.
SPL *dma*		√	√	√	**Store Low P Register**
SPL {*ind*} [*, next ARP*]		√	√	√	Store the low-order bits of the P register (shifted as specified by the PM status bits) at the addressed data-memory location.
SPLK #*lk, dma*			√	√	**Store Parallel Long Immediate**
SPLK #*lk,* {*ind*} [*, next ARP*]				√	Write a full 16-bit pattern into a memory location. The parallel logic unit (PLU) supports this bit manipulation independently of the ALU, so the accumulator is unaffected.
SPM *2-bit constant*		√	√	√	**Set P Register Output Shift Mode** Copy a 2-bit immediate value into the PM field of ST1. This controls shifting of the P register as shown below: PM = 00 Multiplier output is not shifted. PM = 01 Multiplier output is left shifted one place and zero filled. PM = 10 Multiplier output is left shifted four places and zero filled. PM = 11 Multiplier output is right shifted six places and sign extended; the LSBs are lost.
SQRA *dma*		√	√	√	**Square and Accumulate Previous Product**
SQRA {*ind*} [*, next ARP*]		√	√	√	Add the contents of the P register (shifted as specified by the PM status bits) to the accumulator. Then load the contents of the addressed data-memory location into the T register (TMS320C2x/2xx) or TREG0 (TMS320C5x), square the value, and store the result in the P register.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
SQRS *dma*		√	√	√	**Square and Subtract Previous Product**
SQRS {*ind*} [*, next ARP*]		√	√	√	Subtract the contents of the P register (shifted as specified by the PM status bits) to the accumulator. Then load the contents of the addressed data-memory location into the T register (TMS320C2x/2xx) or TREG0 (TMS320C5x), square the value, and store the result in the P register.
SST *dma*	√	√	√	√	**Store Status Register**
SST {*ind*} [*, next ARP*]	√	√	√	√	Store the contents of the ST (TMS320C1x) or ST0 (TMS320C2x/2xx/5x) in the addressed data-memory location.
SST #*n, dma*			√	√	**Store Status Register n**
SST #*n,* {*ind*} [*, next ARP*]			√	√	Store ST*n* in data memory.
SST1 *dma*		√	√	√	**Store Status Register ST1**
SST1 {*ind*} [*, next ARP*]		√	√	√	Store the contents of ST1 in the addressed data-memory location.
SSXM		√	√	√	**Set Sign-Extension Mode**
					Set the SXM status bit to 1; this enables sign extension.
STC		√	√	√	**Set Test/Control Flag**
					Set the TC flag to 1.
STXM		√			**Set Serial Port Transmit Mode**
					Set the TXM status bit to 1.
SUB *dma* [*, shift*]	√	√	√	√	**Subtract From Accumulator With Shift**
SUB {*ind*} [*, shift* [*, next ARP*]]	√	√	√	√	TMS320C1x and TMS320C2x devices: Subtract the contents of the addressed data-memory location from the accumulator. If a shift is specified, left shift the value before subtracting. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
SUB #*k*			√	√	
SUB #*lk* [*, shift*$_2$]			√	√	
					TMS320C2xx and TMS320C5x devices: Subtract the contents of the addressed data-memory location or an 8- or 16-bit constant from the accumulator. If a shift is specified, left shift the data before subtracting. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
SUBB *dma*		√	√	√	**Subtract From Accumulator With Borrow**
SUBB {*ind*} [*, next ARP*]		√	√	√	Subtract the contents of the addressed data-memory location and the value of the carry bit from the accumulator. The carry bit is affected in the normal manner.
SUBC *dma*	√	√	√	√	**Conditional Subtract**
SUBC {*ind*} [*, next ARP*]	√	√	√	√	Perform conditional subtraction. SUBC can be used for division.
SUBH *dma*	√	√	√	√	**Subtract From High Accumulator**
SUBH {*ind*} [*, next ARP*]	√	√		√	Subtract the contents of the addressed data-memory location from the 16 MSBs of the accumulator. The 16 LSBs of the accumulator are not affected.
SUBK *#k*		√	√	√	**Subtract From Accumulator Short Immediate**
					Subtract an 8-bit immediate value from the accumulator. The data is treated as an 8-bit positive number; sign extension is suppressed.
SUBS *dma*	√	√	√	√	**Subtract From Low Accumulator With Sign Extension Suppressed**
SUBS {*ind*} [*, next ARP*]	√	√	√	√	Subtract the contents of the addressed data-memory location from the accumulator. The data is treated as a 16-bit unsigned number; sign extension is suppressed.
SUBT *dma*		√	√	√	**Subtract From Accumulator With Shift Specified by T Register**
SUBT {*ind*} [*, next ARP*]		√	√	√	Left shift the data-memory value as specified by the four LSBs of the T register (TMS320C2x/2xx) or TREG1 (TMS320C5x), and subtract the result from the accumulator. If a shift is specified, left shift the data-memory value before subtracting. During shifting, low-order bits are zero filled, and high-order bits are sign extended if SXM = 1.
SXF		√	√	√	**Set External Flag**
					Set the XF pin and the XF status bit to 1.
TBLR *dma*	√	√	√	√	**Table Read**
TBLR {*ind*} [*, next ARP*]	√	√	√	√	Transfer a word from program memory to a data-memory location. The program-memory address is in the 12 (TMS320C1x) or 16 (TMS320C2x/2xx/5x) LSBs of the accumulator.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
TBLW *dma* **TBLW** {*ind*} [, *next ARP*]	√ √	√ √	√ √	√ √	**Table Write** Transfer a word from data-memory to a program-memory location. The program-memory address is in the 12 (TMS320C1x) or 16 (TMS320C2x/2xx/5x) LSBs of the accumulator.
TRAP		√	√	√	**Software Interrupt** The TRAP instruction is a software interrupt that transfers program control to program-memory address 30h (TMS320C2x) or 22h (TMS320C2xx/5x) and pushes the PC + 1 onto the hardware stack. The instruction at address 30h or 22h may contain a branch instruction to transfer control to the TRAP routine. Putting the PC + 1 on the stack enables an RET instruction to pop the return PC.
XC *n, cond*$_1$ [, *cond*$_2$] [, ...]				√	**Execute Conditionally** Execute conditionally the next *n* instruction words where $1 \leq n \leq 2$. Not all combinations of conditions are meaningful.
XOR *dma* **XOR** {*ind*} [, *next ARP*] **XOR** #*lk* [, *shift*]	√ √	√ √	√ √ √	√ √ √	**Exclusive-OR With Accumulator** TMS320C1x and TMS320C2x devices: Exclusive-OR the contents of the addressed data-memory location with 16 LSBs of the accumulator. The MSBs are not affected. TMS320C2xx and TMS320C5x devices: Exclusive-OR the contents of the addressed data-memory location or a 16-bit immediate value with the accumulator. If a shift is specified, left shift the value before performing the exclusive-OR operation. Low-order bits below and high-order bits above the shifted value are treated as 0s.
XORB				√	**Exclusive-OR of ACCB With Accumulator** Exclusive-OR the contents of the accumulator with the contents of the ACCB. The results are placed in the accumulator.
XORK #*lk* [, *shift*]		√	√	√	**Exclusive-OR Immediate With Accumulator With Shift** Exclusive-OR a 16-bit immediate value with the accumulator. If a shift is specified, left shift the value before performing the exclusive-OR operation. Low-order bits below and high-order bits above the shifted value are treated as 0s.

Table A–3. Instruction Set Comparison (Continued)

Syntax	1x	2x	2xx	5x	Description
XPL [#*lk*,] *dma* **XPL** [#*lk*,] {*ind*} [*, next ARP*]				√ √	**Exclusive-OR of Long Immediate or DBMR With Addressed Data-Memory Value** If a long immediate value is specified, exclusive OR it with the addressed data-memory value; otherwise, exclusive OR the DBMR with the addressed data-memory value. Write the result back to the data-memory location. The accumulator is not affected.
ZAC	√	√	√	√	**Zero Accumulator** Clear the contents of the accumulator to 0.
ZALH *dma* **ZALH** {*ind*} [*, next ARP*]	√ √	√ √	√ √	√ √	**Zero Low Accumulator and Load High Accumulator** Clear the 16 LSBs of the accumulator to 0 and load the contents of the addressed data-memory location into the 16 MSBs of the accumulator.
ZALR *dma* **ZALR** {*ind*} [*, next ARP*]		√ √	√ √	√ √	**Zero Low Accumulator, Load High Accumulator With Rounding** Load the contents of the addressed data-memory location into the 16 MSBs of the accumulator. The value is rounded by 1/2 LSB; that is, the 15 LSBs of the accumulator (0–14) are cleared and bit 15 is set to 1.
ZALS *dma* **ZALS** {*ind*} [*, next ARP*]	√ √	√ √	√ √	√ √	**Zero Accumulator, Load Low Accumulator With Sign Extension Suppressed** Load the contents of the addressed data-memory location into the 16 LSBs of the accumulator. The 16 MSBs are zeroed. The data is treated as a 16-bit unsigned number.
ZAP				√	**Zero the Accumulator and Product Register** The accumulator and product register are zeroed. The ZAP instruction speeds up the preparation for a repeat multiply/accumulate.
ZPR				√	**Zero the Product Register** The product register is cleared.

Appendix B

Submitting ROM Codes to TI

The size of a printed circuit board is a consideration in many DSP applications. To make full use of the board space, Texas Instruments offers a ROM code option that reduces the chip count and provides a single-chip solution. This option allows you to use a code-customized processor for a specific application while taking advantage of:

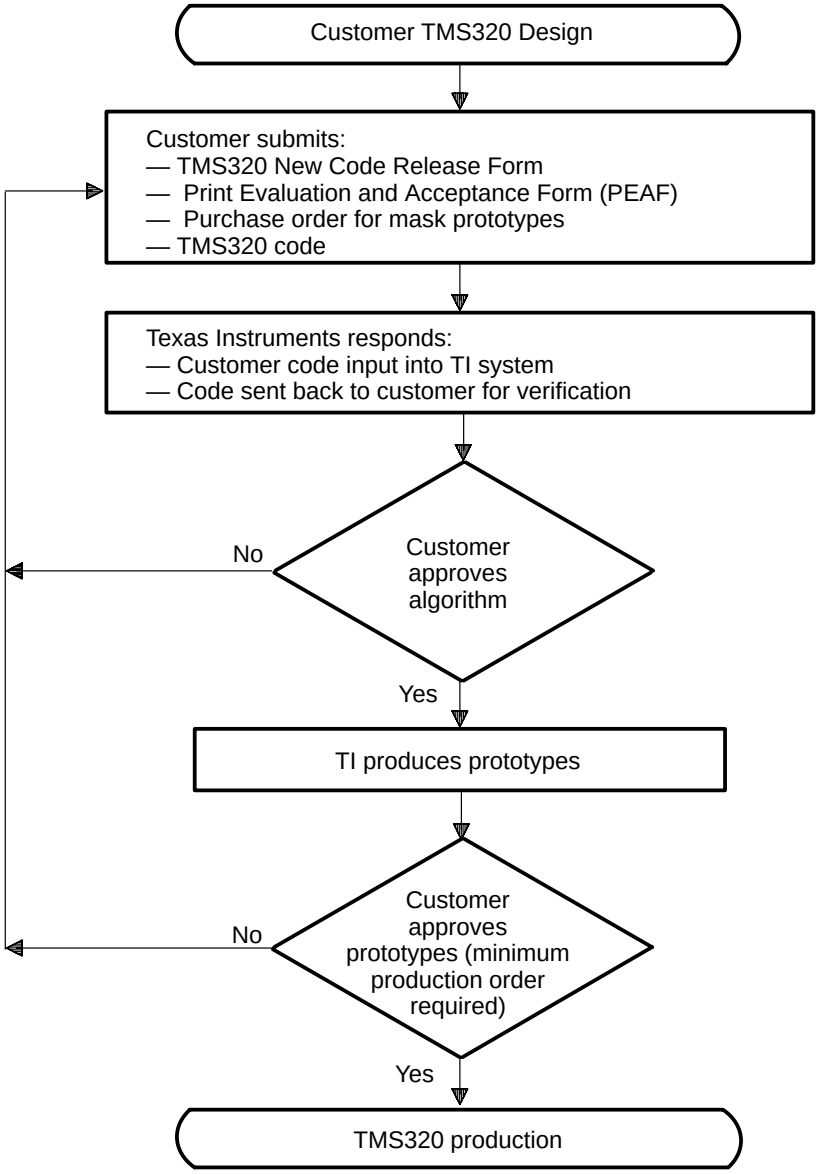
- Greater memory expansion
- Lower system cost
- Less hardware and wiring
- Smaller PCB

If a routine or algorithm is used often, it can be programmed into the on-chip ROM of a TMS320 DSP. TMS320 programs can also be expanded by using external memory; this reduces chip count and allows for a more flexible program memory. Multiple functions are easily implemented by a single device, thus enhancing system capabilities.

TMS320 development tools are used to develop, test, refine, and finalize the algorithms. The microprocessor/microcomputer (MP/$\overline{\text{MC}}$) mode is available on all ROM-coded TMS320 DSP devices when accesses to either on-chip or off-chip memory are required. The microprocessor mode is used to develop, test, and refine a system application. In this mode of operation, the TMS320 acts as a standard microprocessor by using external program memory. When the algorithm has been finalized, the code can be submitted to Texas Instruments for masking into the on-chip program ROM. At that time, the TMS320 becomes a microcomputer that executes customized programs from the on-chip ROM. Should the code need changing or upgrading, the TMS320 can once again be used in the microprocessor mode. This shortens the field-upgrade time and prevents the possibility of inventory obsolescence.

Figure B–1 illustrates the procedural flow for developing and ordering TMS320 masked parts. When ordering, there is a one-time, nonrefundable charge for mask tooling. A minimum production order per year is required for any masked-ROM device. ROM codes will be deleted from the TI system one year after the final delivery.

Figure B–1. TMS320 ROM Code Procedural Flow Chart



The TMS320 ROM code may be submitted in one of the following forms:

- 3-1/2-in floppy: COFF format from macro-assembler/linker (preferred)
- Modem (BBS): COFF format from macro-assembler/linker
- EPROM (others): TMS27C64
- PROM: TBP28S166, TBP28S86

When code is submitted to TI for masking, the code is reformatted to accommodate the TI mask-generation system. System-level verification by the customer is, therefore, necessary to ensure the reformatting remains transparent and does not affect the execution of the algorithm. The formatting changes involve the removal of address-relocation information (the code address begins at the base address of the ROM in the TMS320 device and progresses without gaps to the last address of the ROM) and the addition of data in the reserved locations of the ROM for device ROM test. Because these changes have been made, a checksum comparison is not a valid means of verification.

With each masked-device order, the customer must sign a disclaimer that states:

> The units to be shipped against this order were assembled, for expediency purposes, on a prototype (that is, nonproduction qualified) manufacturing line, the reliability of which is not fully characterized. Therefore, the anticipated inherent reliability of these prototype units cannot be expressly defined.

and a release that states:

> Any masked ROM device may be resymbolized as TI standard product and resold as though it were an unprogrammed version of the device, at the convenience of Texas Instruments.

The use of the ROM-protect feature does not hold for this release statement. Additional risk and charges are involved when the ROM-protect feature is selected. Contact the nearest TI Field Sales Office for more information on procedures, leadtimes, and cost associated with the ROM-protect feature.

Appendix C

Design Considerations for Using the XDS510 Emulator

This appendix assists you in meeting the design requirements of the Texas Instruments XDS510™ emulator for IEEE-1149.1 designs and discusses the XDS510 cable (manufacturing part number 2617698-0001). This cable is identified by a label on the cable pod marked *JTAG 3/5V* and supports both standard 3-V and 5-V target system power inputs.

The term *JTAG,* as used in this book, refers to TI scan-based emulation, which is based on the IEEE 1149.1 standard.

For more information concerning the IEEE 1149.1 standard, contact IEEE Customer Service:

Address: IEEE Customer Service
445 Hoes Lane, PO Box 1331
Piscataway, NJ 08855-1331

Phone: (800) 678–IEEE in the US and Canada
(908) 981–1393 outside the US and Canada

FAX: (908) 981–9667 Telex: 833233

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C.1 Designing Your Target System's Emulator Connector (14-Pin Header)

JTAG target devices support emulation through a dedicated emulation port. This port is accessed directly by the emulator and provides emulation functions that are a superset of those specified by IEEE 1149.1. To communicate with the emulator, *your target system must have a 14-pin header* (two rows of seven pins) with the connections that are shown in Figure C–1. Table C–1 describes the emulation signals.

Although you can use other headers, the recommended unshrouded, straight header has these DuPont connector systems part numbers:

- 65610–114
- 65611–114
- 67996–114
- 67997–114

Figure C–1. 14-Pin Header Signals and Header Dimensions

TMS	1	2	$\overline{\text{TRST}}$
TDI	3	4	GND
PD (V_{CC})	5	6	no pin (key)†
TDO	7	8	GND
TCK_RET	9	10	GND
TCK	11	12	GND
EMU0	13	14	EMU1

Header Dimensions:
Pin-to-pin spacing, 0.100 in. (X,Y)
Pin width, 0.025-in. square post
Pin length, 0.235-in. nominal

† While the corresponding female position on the cable connector is plugged to prevent improper connection, the cable lead for pin 6 is present in the cable and is grounded, as shown in the schematics and wiring diagrams in this appendix.

Table C–1. 14-Pin Header Signal Descriptions

Signal	Description	Emulator† State	Target† State
EMU0	Emulation pin 0	I	I/O
EMU1	Emulation pin 1	I	I/O
GND	Ground		
PD(V_{CC})	Presence detect. Indicates that the emulation cable is connected and that the target is powered up. PD must be tied to V_{CC} in the target system.	I	O
TCK	Test clock. TCK is a 10.368-MHz clock source from the emulation cable pod. This signal can be used to drive the system test clock.	O	I
TCK_RET	Test clock return. Test clock input to the emulator. May be a buffered or unbuffered version of TCK.	I	O
TDI	Test data input	O	I
TDO	Test data output	I	O
TMS	Test mode select	O	I
$\overline{\text{TRST}}$‡	Test reset	O	I

† I = input; O = output

‡ Do not use pullup resistors on $\overline{\text{TRST}}$: it has an internal pulldown device. In a low-noise environment, $\overline{\text{TRST}}$ can be left floating. In a high-noise environment, an additional pulldown resistor may be needed. (The size of this resistor should be based on electrical current considerations.)

C.2 Bus Protocol

The IEEE 1149.1 specification covers the requirements for the test access port (TAP) bus slave devices and provides certain rules, summarized as follows:

- The TMS and TDI inputs are sampled on the rising edge of the TCK signal of the device.
- The TDO output is clocked from the falling edge of the TCK signal of the device.

When these devices are daisy-chained together, the TDO of one device has approximately a half TCK cycle setup time before the next device's TDI signal. This timing scheme minimizes race conditions that would occur if both TDO and TDI were timed from the same TCK edge. The penalty for this timing scheme is a reduced TCK frequency.

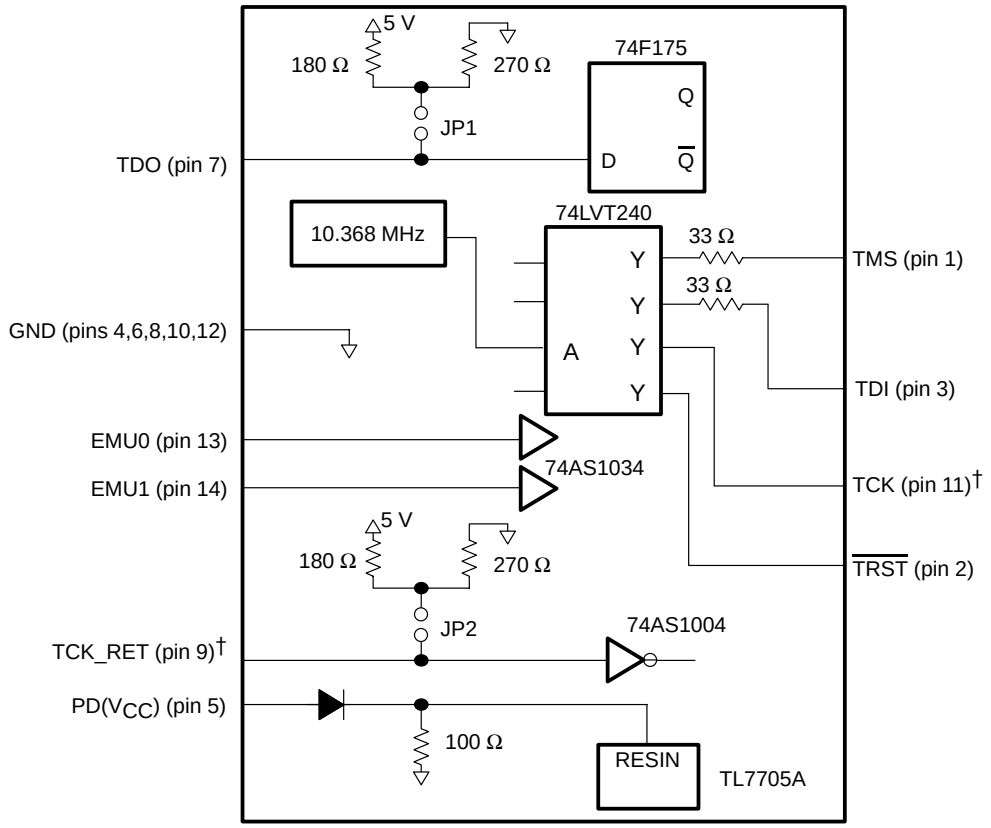
The IEEE 1149.1 specification does not provide rules for bus master (emulator) devices. Instead, it states that the device expects a bus master to provide bus slave compatible timings. The XDS510 provides timings that meet the bus slave rules.

C.3 Emulator Cable Pod

Figure C–2 shows a portion of the emulator cable pod. The functional features of the pod are:

- ❏ TDO and TCK_RET can be parallel-terminated inside the pod if required by the application. By default, these signals are not terminated.
- ❏ TCK is driven with a 74LVT240 device. Because of the high-current drive (32-mA I_{OL}/I_{OH}), this signal can be parallel-terminated. If TCK is tied to TCK_RET, you can use the parallel terminator in the pod.
- ❏ TMS and TDI can be generated from the falling edge of TCK_RET, according to the IEEE 1149.1 bus slave device timing rules.
- ❏ TMS and TDI are series terminated to reduce signal reflections.
- ❏ A 10.368-MHz test clock source is provided. You can also provide your own test clock for greater flexibility.

Figure C–2. Emulator Cable Pod Interface



† The emulator pod uses TCK_RET as its clock source for internal synchronization. TCK is provided as an optional target system test clock source.

C.4 Emulator Cable Pod Signal Timing

Figure C–3 shows the signal timings for the emulator cable pod. Table C–2 defines the timing parameters illustrated in the figure. These timing parameters are calculated from values specified in the standard data sheets for the emulator and cable pod and are for reference only. Texas Instruments does not test or guarantee these timings.

The emulator pod uses TCK_RET as its clock source for internal synchronization. TCK is provided as an optional target system test clock source.

Figure C–3. Emulator Cable Pod Timings

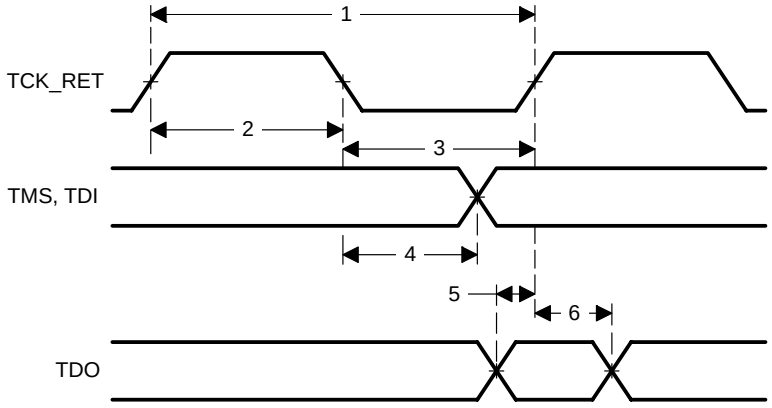


Table C–2. Emulator Cable Pod Timing Parameters

No.	Parameter	Description	Min	Max	Unit
1	$t_{c(TCK)}$	Cycle time, TCK_RET	35	200	ns
2	$t_{w(TCKH)}$	Pulse duration, TCK_RET high	15		ns
3	$t_{w(TCKL)}$	Pulse duration, TCK_RET low	15		ns
4	$t_{d(TMS)}$	Delay time, TMS or TDI valid for TCK_RET low	6	20	ns
5	$t_{su(TDO)}$	Setup time, TDO to TCK_RET high	3		ns
6	$t_{h(TDO)}$	Hold time, TDO from TCK_RET high	12		ns

C.5 Emulation Timing Calculations

The examples in this section help you calculate emulation timings in your system. For actual target timing parameters, see the appropriate data sheet for the device you are emulating.

The examples use the following assumptions:

$t_{su(TTMS)}$	Setup time, target TMS or TDI to TCK high	10 ns
$t_{d(TTDO)}$	Delay time, target TDO from TCK low	15 ns
$t_{d(bufmax)}$	Delay time, target buffer maximum	10 ns
$t_{d(bufmin)}$	Delay time, target buffer minimum	1 ns
$t_{bufskew}$	Skew time, target buffer between two devices in the same package: $[t_{d(bufmax)} - t_{d(bufmin)}] \times 0.15$	1.35 ns
$t_{TCKfactor}$	Duty cycle, assume a 40/60% duty cycle clock	0.4 (40%)

Also, the examples use the following values from Table C–2 on page C-6:

$t_{d(TMSmax)}$	Delay time, emulator TMS or TDI from TCK_RET low, maximum	20 ns
$t_{su(TDOmin)}$	Setup time, TDO to emulator TCK_RET high, minimum	3 ns

There are two key timing paths to consider in the emulation design:

- The TCK_RET-to-TMS or TDI path, called $t_{pd(TCK\_RET-TMS/TDI)}$ (propagation delay time)
- The TCK_RET-to-TDO path, called $t_{pd(TCK\_RET-TDO)}$

In the examples, the worst-case path delay is calculated to determine the maximum system test clock frequency.

Example C–1. Key Timing for a Single-Processor System Without Buffers

- The following example calculates key timing for a single-processor system without buffers.

$$t_{pd\,(TCK\_RET\text{-}TMS/TDI)} = \frac{\left[t_{d\,(TMSmax)} + t_{su\,(TTMS)}\right]}{t_{TCKfactor}}$$

$$= \frac{(20\text{ ns} + 10\text{ ns})}{0.4}$$

$$= 75\text{ ns, or } 13.3\text{ MHz}$$

$$t_{pd\,(TCK\_RET–TDO)} = \frac{\left[t_{d\,(TTDO)} + t_{su\,(TDOmin)}\right]}{t_{TCKfactor}}$$

$$= \frac{(15\text{ ns} + 3\text{ ns})}{0.4}$$

$$= 45\text{ ns, or } 22.2\text{ MHz}$$

In the preceding example, the TCK_RET-to-TMS/TDI path is the limiting factor because it requires more time to complete.

- The following example calculates key timing for a single- or multiple-processor system with buffered input and output:

$$t_{pd\,(TCK\_RET\text{-}TMS/TDI)} = \frac{\left[t_{d\,(TMSmax)} + t_{su\,(TTMS)} + t_{bufskew}\right]}{t_{TCKfactor}}$$

$$= \frac{(20\text{ ns} + 10\text{ ns} + 1.35\text{ ns})}{0.4}$$

$$= 78.4\text{ ns, or } 12.7\text{ MHz}$$

$$t_{pd\,(TCK\_RET–TDO)} = \frac{\left[t_{d\,(TTDO)} + t_{su\,(TDOmin)} + t_{d\,(bufmax)}\right]}{t_{TCKfactor}}$$

$$= \frac{(15\text{ ns} + 3\text{ ns} + 10\text{ ns})}{0.4}$$

$$= 70\text{ ns, or } 14.3\text{ MHz}$$

In the preceding example, the TCK_RET-to-TMS/TDI path is the limiting factor becaise it requires more time to complete.

In a multiprocessor application, it is necessary to ensure that the EMU0 and EMU1 lines can go from a logic-low level to a logic-high level in less than 10 µs, this parameter is called rise time, t_r. This can be calculated as follows:

$$
\begin{aligned}
t_r &= 5(R_{pullup} \times N_{devices} \times C_{load\_per\_device}) \\
&= 5(4.7\ k\Omega \times 16 \times 15\ pF) \\
&= 5(4.7 \times 10^3\ \Omega \times 16 \times 15 = no\ ^{-12}\ F) \\
&= 5(1128 \times 10^{-9}\ ) \\
&= 5.64\ \mu s
\end{aligned}
$$

C.6 Connections Between the Emulator and the Target System

It is extremely important to provide high-quality signals between the emulator and the JTAG target system. You must supply the correct signal buffering, test clock inputs, and multiple processor interconnections to ensure proper emulator and target system operation.

Signals applied to the EMU0 and EMU1 pins on the JTAG target device can be either input or output. In general, these two pins are used as both input and output in multiprocessor systems to handle global run/stop operations. EMU0 and EMU1 signals are applied only as inputs to the XDS510 emulator header.

C.6.1 Buffering Signals

If the distance between the emulation header and the JTAG target device is greater than 6 inches, the emulation signals must be buffered. If the distance is less than 6 inches, no buffering is necessary. Figure C–4 shows the simpler, no-buffering situation.

The distance between the header and the JTAG target device must be no more than 6 inches. The EMU0 and EMU1 signals must have pullup resistors connected to V_{CC} to provide a signal rise time of less than 10 μs. A 4.7-kΩ resistor is suggested for most applications.

Figure C–4. Emulator Connections Without Signal Buffering

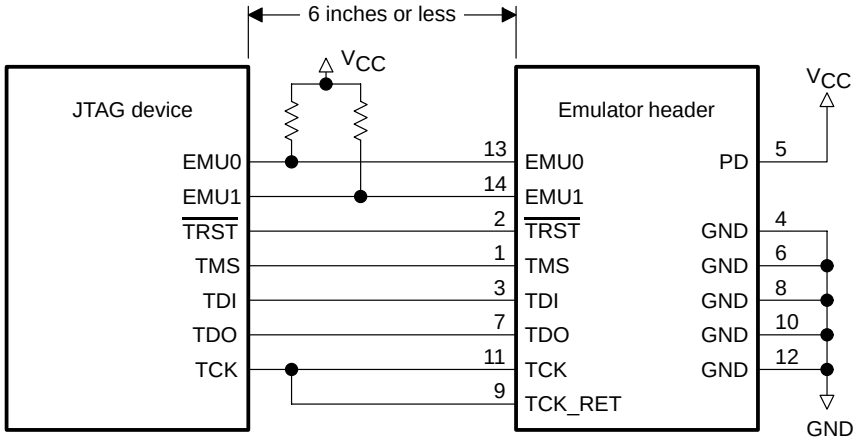
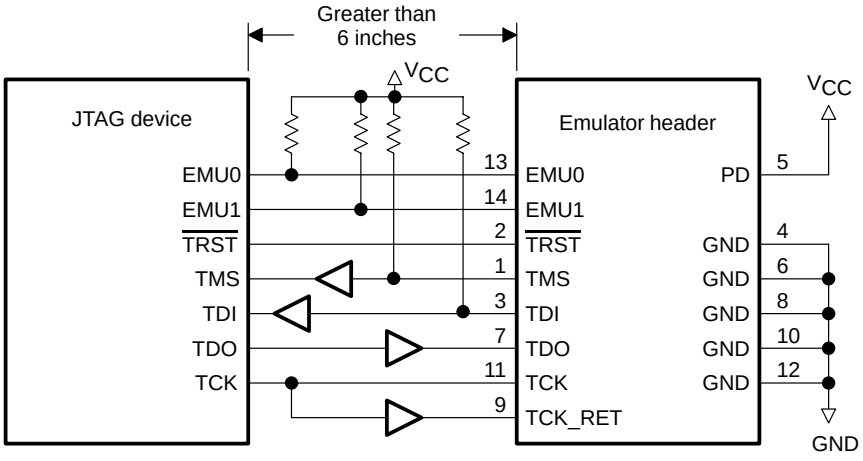


Figure C–5 shows the connections necessary for buffered transmission signals. The distance between the emulation header and the processor is greater than 6 inches. Emulation signals TMS, TDI, TDO, and TCK_RET are buffered through the same device package.

Figure C–5. Emulator Connections With Signal Buffering



The EMU0 and EMU1 signals must have pullup resistors connected to V_{CC} to provide a signal rise time of less than 10 μs. A 4.7-kΩ resistor is suggested for most applications.

The input buffers for TMS and TDI should have pullup resistors connected to V_{CC} to hold these signals at a known value when the emulator is not connected. A resistor value of 4.7 kΩ or greater is suggested.

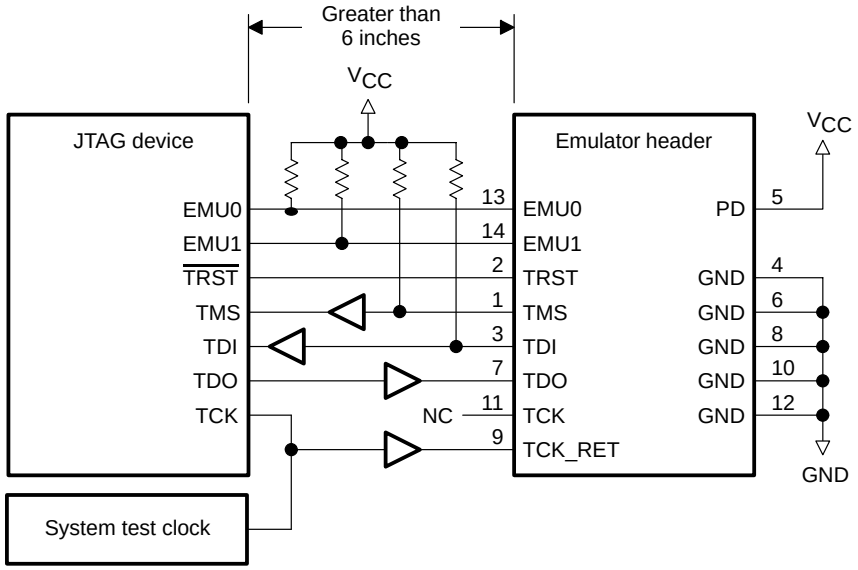
To have high-quality signals (especially the processor TCK and the emulator TCK_RET signals), you may have to employ special care when routing the printed wiring board trace. You also may have to use termination resistors to match the trace impedance. The emulator pod provides optional internal parallel terminators on the TCK_RET and TDO. TMS and TDI provide fixed series termination.

Because $\overline{\text{TRST}}$ is an asynchronous signal, it should be buffered as needed to ensure sufficient current to all target devices.

C.6.2 Using a Target-System Clock

Figure C–6 shows an application with the system test clock generated in the target system. In this application, the emulator's TCK signal is left unconnected.

Figure C–6. Target-System-Generated Test Clock



Note: When the TMS and TDI lines are buffered, pullup resistors must be used to hold the buffer inputs at a known level when the emulator cable is not connected.

There are two benefits in generating the test clock in the target system:

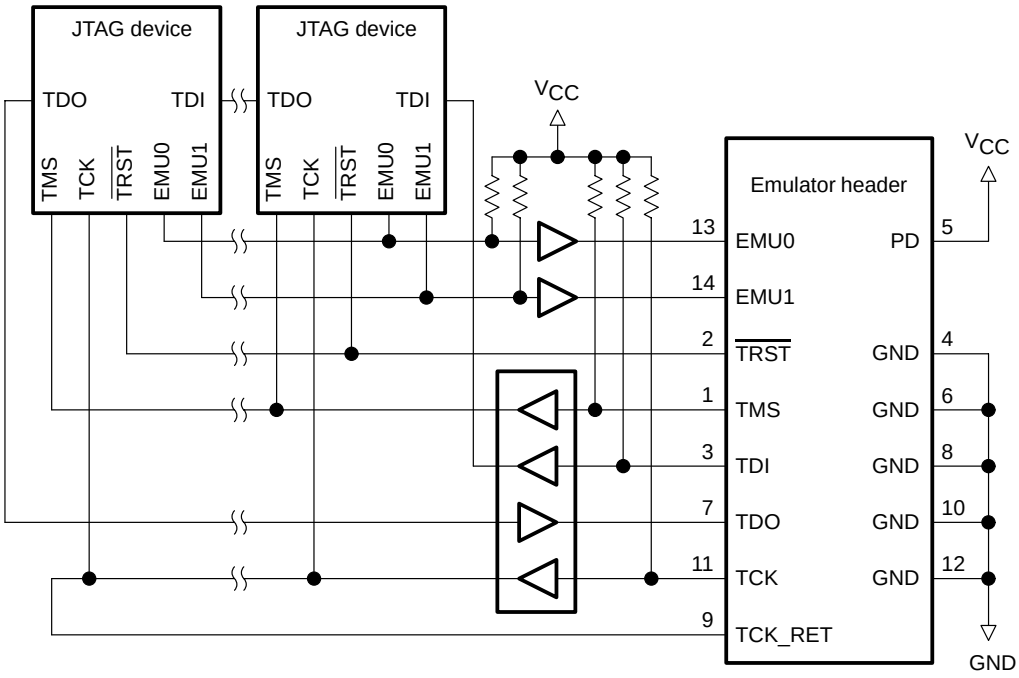
- The emulator provides only a single 10.368-MHz test clock. If you allow the target system to generate your test clock, you can set the frequency to match your system requirements.
- In some cases, you may have other devices in your system that require a test clock when the emulator is not connected. The system test clock also serves this purpose.

C.6.3 Configuring Multiple Processors

Figure C–7 shows a typical daisy-chained multiprocessor configuration that meets the minimum requirements of the IEEE 1149.1 specification. The emulation signals are buffered to isolate the processors from the emulator and provide adequate signal drive for the target system. One of the benefits of this interface is that you can slow down the test clock to eliminate timing problems. Follow these guidelines for multiprocessor support:

- The processor TMS, TDI, TDO, and TCK signals must be buffered through the same physical device package for better control of timing skew.
- The input buffers for TMS, TDI, and TCK must have pullup resistors connected to V_{CC} to hold these signals at a known value when the emulator is not connected. A resistor value of 4.7 kΩ or greater is suggested.
- Buffering EMU0 and EMU1 is optional but highly recommended to provide isolation. These are not critical signals and do not have to be buffered through the same physical package as TMS, TCK, TDI, and TDO.

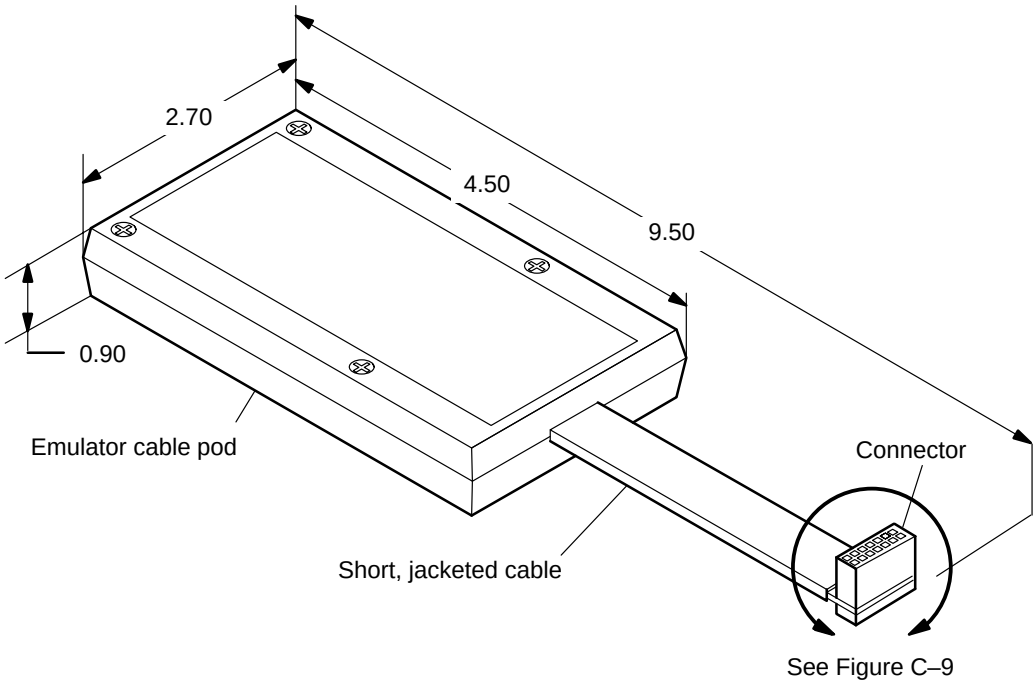
Figure C–7. Multiprocessor Connections



C.7 Physical Dimensions for the 14-Pin Emulator Connector

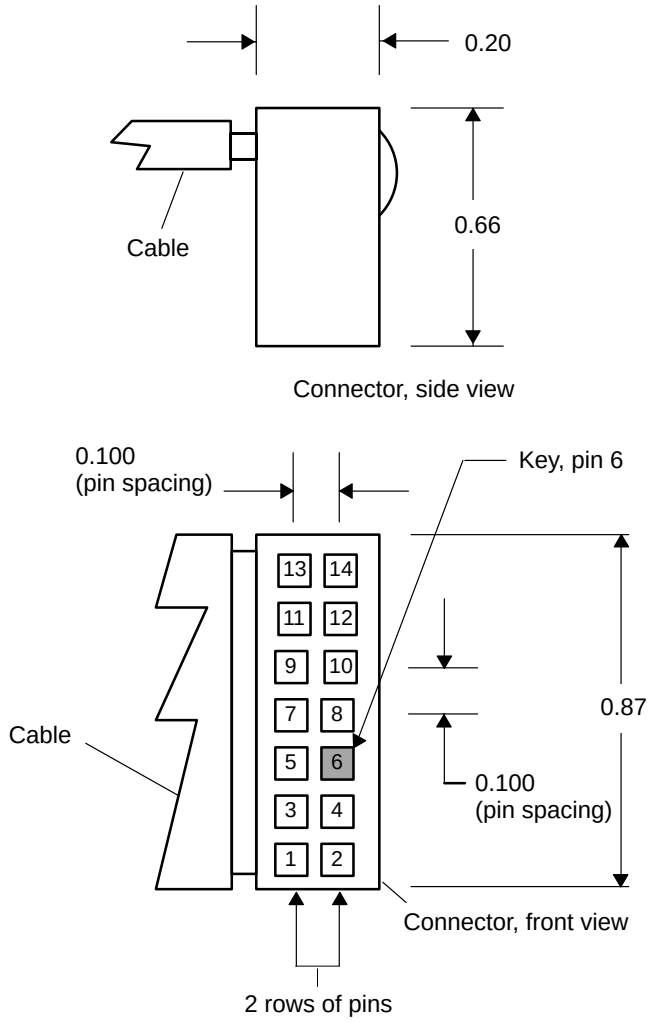
The JTAG emulator target cable consists of a 3-foot section of jacketed cable that connects to the emulator, an active cable pod, and a short section of jacketed cable that connects to the target system. The overall cable length is approximately 3 feet 10 inches. Figure C–8 and Figure C–9 show the physical dimensions for the target cable pod and short cable. The cable pod box is nonconductive plastic with four recessed metal screws.

Figure C–8. Pod/Connector Dimensions



Note: All dimensions are in inches and are nominal dimensions, unless otherwise specified. Pin-to-pin spacing on the connector is 0.100 inches in both the X and Y planes.

Figure C–9. 14-Pin Connector Dimensions



Note: All dimensions are in inches and are nominal dimensions, unless otherwise specified. Pin-to-pin spacing on the connector is 0.100 inches in both the X and Y planes.

C.8 Emulation Design Considerations

This section describes the use and application of the scan path linker (SPL), which can simultaneously add all four secondary JTAG scan paths to the main scan path. It also describes the use of the emulation pins and the configuration of multiple processors.

C.8.1 Using Scan Path Linkers

You can use the TI ACT8997 scan path linker (SPL) to divide the JTAG emulation scan path into smaller, logically connected groups of 4 to 16 devices. As described in the *Advanced Logic and Bus Interface Logic Data Book*, the SPL is compatible with the JTAG emulation scanning. The SPL is capable of adding any combination of its four secondary scan paths into the main scan path.

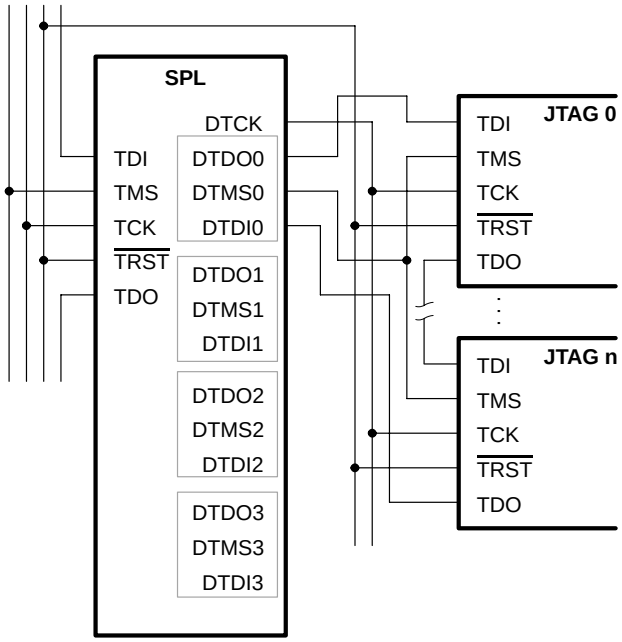
A system of multiple, secondary JTAG scan paths has better fault tolerance and isolation than a single scan path. Since an SPL has the capability of adding all secondary scan paths to the main scan path simultaneously, it can support global emulation operations, such as starting or stopping a selected group of processors.

TI emulators do not support the nesting of SPLs (for example, an SPL connected to the secondary scan path of another SPL). However, you can have multiple SPLs on the main scan path.

Scan path selectors are not supported by this emulation system. The TI ACT8999 scan path selector is similar to the SPL, but it can add only one of its secondary scan paths at a time to the main JTAG scan path. Thus, global emulation operations are not assured with the scan path selector.

You can insert an SPL on a backplane so that you can add up to four device boards to the system without the jumper wiring required with nonbackplane devices. You connect an SPL to the main JTAG scan path in the same way you connect any other device. Figure C–10 shows how to connect a secondary scan path to an SPL.

Figure C–10. Connecting a Secondary JTAG Scan Path to a Scan Path Linker



The $\overline{\text{TRST}}$ signal from the main scan path drives all devices, even those on the secondary scan paths of the SPL. The TCK signal on each target device on the secondary scan path of an SPL is driven by the SPL's DTCK signal. The TMS signal on each device on the secondary scan path is driven by the respective DTMS signals on the SPL.

DTDO0 on the SPL is connected to the TDI signal of the first device on the secondary scan path. DTDI0 on the SPL is connected to the TDO signal of the last device in the secondary scan path. Within each secondary scan path, the TDI signal of a device is connected to the TDO signal of the device before it. If the SPL is on a backplane, its secondary JTAG scan paths are on add-on boards; if signal degradation is a problem, you may need to buffer both the $\overline{\text{TRST}}$ and DTCK signals. Although degradation is less likely for DTMS*n* signals, you may also need to buffer them for the same reasons.

C.8.2 Emulation Timing Calculations for a Scan Path Linker (SPL)

The examples in this section help you to calculate the key emulation timings in the SPL secondary scan path of your system. For actual target timing parameters, see the appropriate device data sheet for your target device.

The examples use the following assumptions:

$t_{su(TTMS)}$	Setup time, target TMS/TDI to TCK high	10 ns
$t_{d(TTDO)}$	Delay time, target TDO from TCK low	15 ns
$t_{d(bufmax)}$	Delay time, target buffer, maximum	10 ns
$t_{d(bufmin)}$	Delay time, target buffer, minimum	1 ns
$t_{(bufskew)}$	Skew time, target buffer, between two devices in the same package: $[t_{d(bufmax)} - t_{d(bufmin)}] \times 0.15$	1.35 ns
$t_{(TCKfactor)}$	Duty cycle, TCK assume a 40/60% clock	0.4 (40%)

Also, the examples use the following values from the SPL data sheet:

$t_{d(DTMSmax)}$	Delay time, SPL DTMS/DTDO from TCK low, maximum	31 ns
$t_{su(DTDLmin)}$	Setup time, DTDI to SPL TCK high, minimum	7 ns
$t_{d(DTCKHmin)}$	Delay time, SPL DTCK from TCK high, minimum	2 ns
$t_{d(DTCKLmax)}$	Delay time, SPL DTCK from TCK low, maximum	16 ns

There are two key timing paths to consider in the emulation design:

- ❑ The TCK-to-DTMS/DTDO path, called $t_{pd(TCK-DTMS)}$
- ❑ The TCK-to-DTDI path, called $t_{pd(TCK-DTDI)}$

In the following two examples, the worst-case path delay is calculated to determine the maximum system test clock frequency.

Example C–2. Key Timing for a Single-Processor System Without Buffering (SPL)

- The following example calculates key timing for a single-processor system without buffering (SPL):

$$t_{pd\,(TCK\text{-}DTMS)} = \frac{\left[t_{d\,(DTMSmax)} + t_{d\,(DTCKHmin)} + t_{su\,(TTMS)}\right]}{t_{TCKfactor}}$$

$$= \frac{(31\text{ ns} + 2\text{ ns} + 10\text{ ns})}{0.4}$$

$$= 107.5\text{ ns, or } 9.3\text{ MHz}$$

$$t_{pd\,(TCK\text{-}DTDI)} = \frac{\left[t_{d\,(TTDO)} + t_{d\,(DTCKLmax)} + t_{su\,(DTDLmin)}\right]}{t_{TCKfactor}}$$

$$= \frac{(15\text{ ns} + 16\text{ ns} + 7\text{ ns})}{0.4}$$

$$= 9.5\text{ ns, or } 10.5\text{ MHz}$$

In the preceding example, the TCK-to-DTMS/DTDL path is the limiting factor.

- The following example calculates key timing for a single- or multiprocessor-system with buffered input and output (SPL):

$$t_{pd\,(TCK\text{-}TDMS)} = \frac{\left[t_{d\,(DTMSmax)} + t_{(DTCKHmin)} + t_{su\,(TTMS)} + t_{(bufskew)}\right]}{t_{TCKfactor}}$$

$$= \frac{(31\text{ ns} + 2\text{ ns} + 10\text{ ns} + 1.35\text{ ns})}{0.4}$$

$$= 110.9\text{ ns, or } 9.0\text{ MHz}$$

$$t_{pd\,(TCK–DTDI)} = \frac{\left[t_{d\,(TTDO)} + t_{d\,(DTCKLmax)} + t_{su\,(DTDLmin)} + t_{d\,(bufskew)}\right]}{t_{TCKfactor}}$$

$$= \frac{(15\text{ ns} + 15\text{ ns} + 7\text{ ns} + 10\text{ ns})}{0.4}$$

$$= 120\text{ ns, or } 8.3\text{ MHz}$$

In the preceding example, the TCK-to-DTDI path is the limiting factor.

C.8.3 Using Emulation Pins

The EMU0/1 pins of TI devices are bidirectional, 3-state output pins. When in an inactive state, these pins are at high impedance. When the pins are active, they provide one of two types of output:

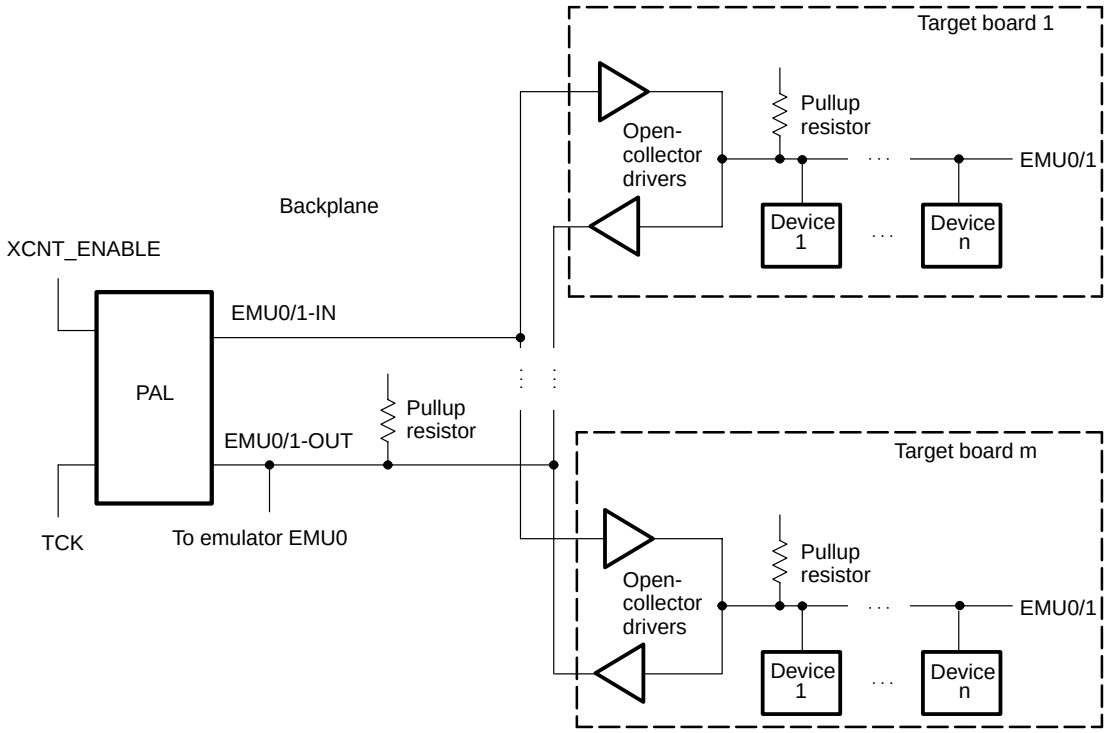
- Signal event. The EMU0/1 pins can be configured via software to signal internal events. In this mode, driving one of these pins low can cause devices to signal such events. To enable this operation, the EMU0/1 pins function as open-collector sources. External devices such as logic analyzers can also be connected to the EMU0/1 signals in this manner. If such an external source is used, it must also be connected via an open-collector source.

- External count. The EMU0/1 pins can be configured via software as totem-pole outputs for driving an external counter. If the output of more than one device is configured for totem-pole operation, then these devices can be damaged. The emulation software detects and prevents this condition. However, the emulation software has no control over external sources on the EMU0/1 signal. Therefore, all external sources must be inactive when any device is in the external count mode.

TI devices can be configured by software to halt processing if their EMU0/1 pins are driven low. This feature combined with the signal event output, allows one TI device to halt all other TI devices on a given event for system-level debugging.

If you route the EMU0/1 signals between multiple boards, they require special handling because they are more complex than normal emulation signals. Figure C–11 shows an example configuration that allows any processor in the system to stop any other processor in the system. Do not tie the EMU0/1 pins of more than 16 processors together in a single group without using buffers. Buffers provide the crisp signals that are required during a RUNB (run benchmark) debugger command or when the external analysis counter feature is used.

Figure C–11. EMU0/1 Configuration to Meet Timing Requirements of Less Than 25 ns



Notes: 1) The low time on EMU0/1-IN must be at least one TCK cycle and less than 10 μs. Software sets the EMU0/1-OUT pin to a high state.

2) To enable the open-collector driver and pullup resistor on EMU1 to provide rise/fall times of less than 25 ns, the modification shown in this figure is suggested. Rise times of more than 25 ns can cause the emulator to detect false edges during the RUNB command or when the external counter selected from the debugger analysis menu is used.

These seven important points apply to the circuitry shown in Figure C–11 and the timing shown in Figure C–12:

- ❏ Open-collector drivers isolate each board. The EMU0/1 pins are tied together on each board.
- ❏ At the board edge, the EMU0/1 signals are split to provide both input and output connections. This is required to prevent the open-collector drivers from acting as latches that can be set only once.
- ❏ The EMU0/1 signals are bused down the backplane. Pullup resistors must be installed as required.

- The bused EMU0/1 signals go into a programmable logic array device PAL®, whose function is to generate a low pulse on the EMU0/1-IN signal when a low level is detected on the EMU0/1-OUT signal. This pulse must be longer than one TCK period to affect the devices but less than 10 μs to avoid possible conflicts or retriggering once the emulation software clears the device's pins.
- During a RUNB debugger command or other external analysis count, the EMU0/1 pins on the target device become totem-pole outputs. The EMU1 pin is a ripple carry-out of the internal counter. EMU0 becomes a *processor-halted* signal. During a RUNB or other external analysis count, the EMU0/1-IN signal to all boards must remain in the high (disabled) state. You must provide some type of external input (XCNT_ENABLE) to the PAL® to disable the PAL® from driving EMU0/1-IN to a low state.
- If you use sources other than TI processors (such as logic analyzers) to drive EMU0/1, their signal lines must be isolated by open-collector drivers and be inactive during RUNB and other external analysis counts.
- You must connect the EMU0/1-OUT signals to the emulation header or directly to a test bus controller.

Figure C–12. Suggested Timings for the EMU0 and EMU1 Signals

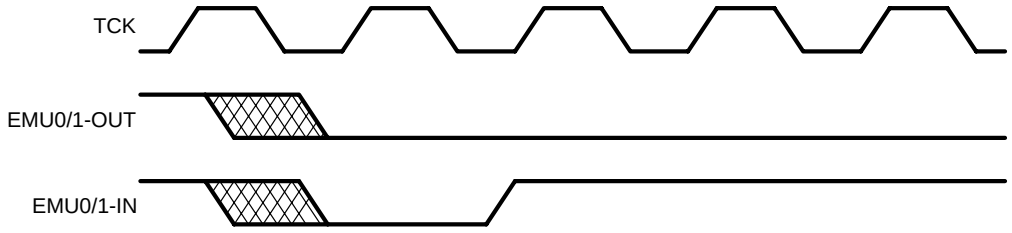
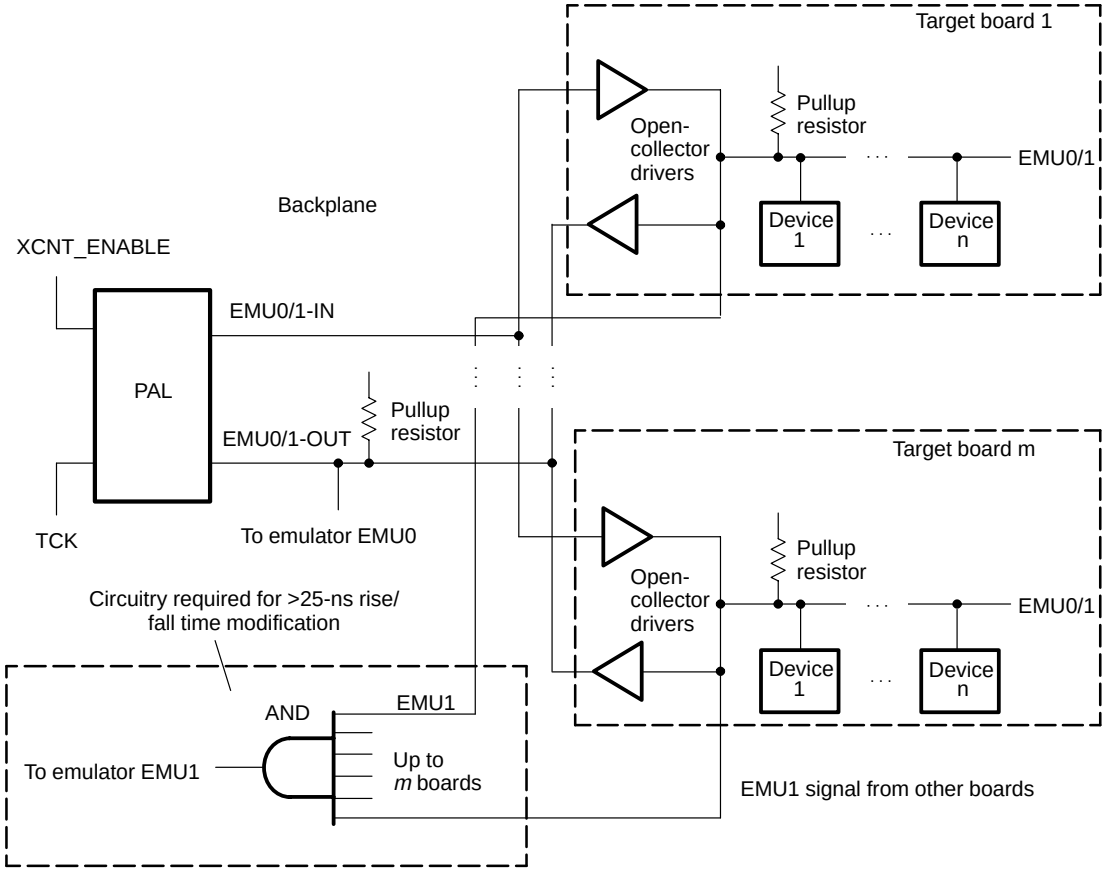


Figure C–13. EMU0/1 Configuration With Additional AND Gate to Meet Timing Requirements of Greater Than 25 ns

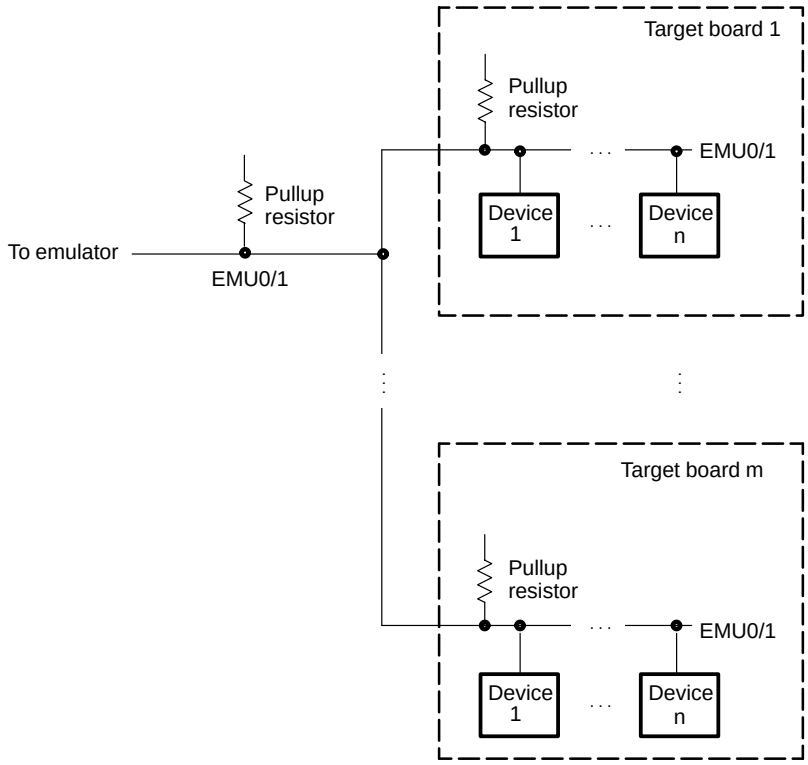


Notes: 1) The low time on EMU0/1-IN must be at least one TCK cycle and less than 10 μs. Software sets the EMU0/1-OUT pin to a high state.

2) To enable the open-collector driver and pullup resistor on EMU1 to provide rise/fall time of greater than 25 ns, the modification shown in this figure is suggested. Rise times of more than 25 ns can cause the emulator to detect false edges during the RUNB command or when the external counter selected from the debugger analysis menu is used.

You do not need to have devices on one target board stop devices on another target board using the EMU0/1 signals (see the circuit in Figure C–14). In this configuration, the global-stop capability is lost. It is important not to overload EMU0/1 with more than 16 devices.

Figure C–14. EMU0/1 Configuration Without Global Stop

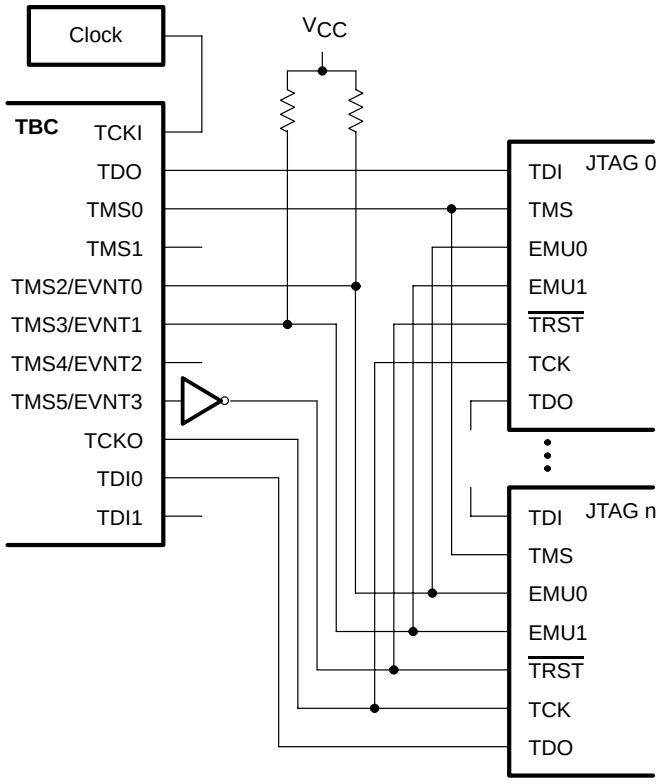


Note: The open-collector driver and pullup resistor on EMU1 must be able to provide rise/fall times of less than 25 ns. Rise times of more than 25 ns can cause the emulator to detect false edges during the RUNB command or when the external counter selected from the debugger analysis menu is used. If this condition cannot be met, then the EMU0/1 signals from the individual boards must be ANDed together (as shown in Figure C–14) to produce an EMU0/1 signal for the emulator.

C.8.4 Performing Diagnostic Applications

For systems that require built-in diagnostics, it is possible to connect the emulation scan path directly to a TI ACT8990 test bus controller (TBC) instead of the emulation header. The TBC is described in the Texas Instruments *Advanced Logic and Bus Interface Logic Data Book*. Figure C–15 shows the scan path connections of *n* devices to the TBC.

Figure C–15. TBC Emulation Connections for n JTAG Scan Paths



In the system design shown in Figure C–15, the TBC emulation signals TCKI, TDO, TMS0, TMS2/EVNT0, TMS3/EVNT1, TMS5/EVNT3, TCKO, and TDI0 are used, and TMS1, TMS4/EVNT2, and TDI1 are not connected. The target devices' EMU0 and EMU1 signals are connected to V_{CC} through pullup resistors and tied to the TBC's TMS2/EVNT0 and TMS3/EVNT1 pins, respectively. The TBC's TCKI pin is connected to a clock generator. The TCK signal for the main JTAG scan path is driven by the TBC's TCKO pin.

On the TBC, the TMS0 pin drives the TMS pins on each device on the main JTAG scan path. TDO on the TBC connects to TDI on the first device on the main JTAG scan path. TDI0 on the TBC is connected to the TDO signal of the last device on the main JTAG scan path. Within the main JTAG scan path, the TDI signal of a device is connected to the TDO signal of the device before it. $\overline{\text{TRST}}$ for the devices can be generated either by inverting the TBC's TMS5/EVNT3 signal for software control or by logic on the board itself.

Appendix D

Glossary

A

A0–A15: Collectively, the external address bus; the 16 pins are used in parallel to address external data memory, program memory, or I/O space.

ACC: See *accumulator.*

ACCH: *Accumulator high word.* The upper 16 bits of the accumulator. See also *accumulator*.

ACCL: *Accumulator low word.* The lower 16 bits of the accumulator. See also *accumulator*.

accumulator: A 32-bit register that stores the results of operations in the central arithmetic logic unit (CALU) and provides an input for subsequent CALU operations. The accumulator also performs shift and rotate operations.

address: The location of program code or data stored in memory.

addressing mode: A method by which an instruction interprets its operands to acquire the data it needs. See also *direct addressing*; *immediate addressing*; *indirect addressing*.

analog-to-digital (A/D) converter: A circuit that translates an analog signal to a digital signal.

AR: See *auxiliary register*.

AR0–AR7: *Auxiliary registers 0 through 7.* See *auxiliary register*.

ARAU: See *auxiliary register arithmetic unit (ARAU).*

ARB: See *auxiliary register pointer buffer (ARB)*.

ARP: See *auxiliary register pointer (ARP).*

auxiliary register: One of eight 16-bit registers (AR7–AR0) used as pointers to addresses in data space. The registers are operated on by the auxiliary register arithmetic unit (ARAU) and are selected by the auxiliary register pointer (ARP).

auxiliary register arithmetic unit (ARAU): A 16-bit arithmetic unit used to increment, decrement, or compare the contents of the auxiliary registers. Its primary function is manipulating auxiliary register values for indirect addressing.

auxiliary register pointer (ARP): A 3-bit field in status register ST0 that points to the current auxiliary register.

auxiliary register pointer buffer (ARB): A 3-bit field in status register ST1 that holds the previous value of the auxiliary register pointer (ARP).

B

B0: An on-chip block of dual-access RAM that can be configured as either data memory or program memory, depending on the value of the CNF bit in status register ST1.

B1: An on-chip block of dual-access RAM available for data memory.

B2: An on-chip block of dual-access RAM available for data memory.

$\overline{\text{BIO}}$ pin: A general-purpose input pin that can be tested by conditional instructions that cause a branch when an external device drives $\overline{\text{BIO}}$ low.

bit-reversed indexed addressing: A method of indirect addressing that allows efficient I/O operations by resequencing the data points in a radix-2 fast Fourier transform (FFT) program. The direction of carry propagation in the ARAU is reversed.

boot loader: A built-in segment of code that transfers code from an 8-bit external source to a 16-bit external program destination at reset.

$\overline{\text{BOOT}}$ pin: The pin that enables the on-chip boot loader. When $\overline{\text{BOOT}}$ is held low, the processor executes the boot loader program after a hardware reset. When $\overline{\text{BOOT}}$ is held high, the processor skips execution of the boot loader and accesses off-chip program-memory at reset.

$\overline{\text{BR}}$: *Bus request pin.* This pin is tied to the $\overline{\text{BR}}$ signal, which is asserted when a global data memory access is initiated.

branch: A switching of program control to a nonsequential program-memory address.

C

C bit: See *carry bit.*

CALU: See *central arithmetic logic unit (CALU).*

carry bit: Bit 9 of status register ST1; used by the CALU for extended arithmetic operations and accumulator shifts and rotates. The carry bit can be tested by conditional instructions.

central arithmetic logic unit (CALU): The 32-bit wide main arithmetic logic unit for the 'C24x CPU that performs arithmetic and logic operations. It accepts 32-bit values for operations, and its 32-bit output is held in the accumulator.

CLK register: *CLKOUT1-pin control register.* Bit 0 of determines whether the CLKOUT1 signal is available at the CLKOUT1 pin.

CLKIN: *Input clock signal*. A clock source signal supplied to the on-chip clock generator at the CLKIN/X2 pin or generated internally by the on-chip oscillator. The clock generator divides or multiplies CLKIN to produce the CPU clock signal, CLKOUT1.

CLKOUT1: *Master clock output signal*. The output signal of the on-chip clock generator. The CLKOUT1 high pulse signifies the CPU's logic phase (when internal values are changed), and the CLKOUT1 low pulse signifies the CPU's latch phase (when the values are held constant).

CLKOUT1 cycle: See *CPU cycle*.

CLKOUT1-pin control register: See *CLK register*.

clock mode (clock generator): One of the modes which sets the internal CPU clock frequency to a fraction or multiple of the frequency of the input clock signal CLKIN.

CNF bit: *DARAM configuration bit*. Bit 12 in status register ST1. CNF is used to determine whether the on-chip RAM block B0 is mapped to program space or data space.

codec: A device that codes in one direction of transmission and decodes in another direction of transmission.

COFF: *Common object file format.* A system of files configured according to a standard developed by AT&T. These files are relocatable in memory space.

context saving/restoring: Saving the system status when the device enters a subroutine (such as an interrupt service routine) and restoring the system status when exiting the subroutine. On the ’C24x, only the program counter value is saved and restored automatically; other context saving and restoring must be performed by the subroutine.

CPU: *Central processing unit*. The ’C24x CPU is the portion of the processor involved in arithmetic, shifting, and Boolean logic operations, as well as the generation of data- and program-memory addresses. The CPU includes the central arithmetic logic unit (CALU), the multiplier, and the auxiliary register arithmetic unit (ARAU).

CPU cycle: The time required for the CPU to go through one logic phase (during which internal values are changed) and one latch phase (during which the values are held constant).

current AR: See *current auxiliary register*.

current auxiliary register: The auxiliary register pointed to by the auxiliary register pointer (ARP). The auxiliary registers are AR0 (ARP = 0) through AR7 (ARP = 7). See also *auxiliary register*; *next auxiliary register*.

current data page: The data page indicated by the content of the data page pointer (DP). See also *data page*; *DP*.

D

D0–D15: Collectively, the external data bus; the 16 pins are used in parallel to transfer data between the ’C24x and external data memory, program memory, or I/O space.

DARAM: *Dual-access RAM*. RAM that can be accessed twice in a single CPU clock cycle. For example, your code can read from and write to DARAM in the same clock cycle.

DARAM configuration bit (CNF): See *CNF bit*.

data-address generation logic: Logic circuitry that generates the addresses for data memory reads and writes. This circuitry, which includes the auxiliary registers and the ARAU, can generate one address per machine cycle. See also *program-address generation logic.*

data page: One block of 128 words in data memory. Data memory contains 512 data pages. Data page 0 is the first page of data memory (addresses 0000h–007Fh); data page 511 is the last page (addresses FF80h–FFFFh). See also *data page pointer (DP)*; *direct addressing*.

data page 0: Addresses 0000h–007Fh in data memory; contains the memory-mapped registers, a reserved test/emulation area for special information transfers, and the scratch-pad RAM block (B2).

data page pointer (DP): A 9-bit field in status register ST0 that specifies which of the 512 data pages is currently selected for direct address generation. When an instruction uses direct addressing to access a data-memory value, the DP provides the nine MSBs of the data-memory address, and the instruction provides the seven LSBs.

data-read address bus (DRAB): A 16-bit internal bus that carries the address for each read from data memory.

data read bus (DRDB): A 16-bit internal bus that carries data from data memory to the CALU and the ARAU.

data-write address bus (DWAB): A 16-bit internal bus that carries the address for each write to data memory.

data write bus (DWEB): A 16-bit internal bus that carries data to both program memory and data memory.

decode phase: The phase of the pipeline in which the instruction is decoded. See also *pipeline*; *instruction-fetch phase*; *operand-fetch phase; instruction-execute phase*.

direct addressing: One of the methods used by an instruction to address data-memory. In direct addressing, the data-page pointer (DP) holds the nine MSBs of the address (the current data page), and the instruction word provides the seven LSBs of the address (the offset). See also *indirect addressing*.

DIV2/DIV1: Two pins used together to determine the clock mode of the ’C24x clock generator ($\div 2$, $\times 1$, $\times 2$, or $\times 4$).

DP: See *data page pointer (DP).*

DRAB: See *data-read address bus (DRAB).*

DRDB: See *data read bus (DRDB).*

$\overline{\text{DS}}$: *Data memory select pin*. The ’C24x asserts $\overline{\text{DS}}$ to indicate an access to external data memory (local or global).

DSWS: *Data-space wait-state bit(s).* A value in the wait-state generator control register (WSGR) that determines the number of wait states applied to reads from and writes to off-chip data space.

dual-access RAM: See *DARAM*.

dummy cycle: A CPU cycle in which the CPU intentionally reloads the program counter with the same address.

DWAB: See *data-write address bus (DWAB).*

DWEB: See *data write bus (DWEB).*

E

execute phase: The fourth phase of the pipeline; the phase in which the instruction is executed. See also *pipeline*; *instruction-fetch phase*; *instruction-decode phase*; *operand-fetch phase*.

external interrupt: A hardware interrupt triggered by an external event sending an input through an interrupt pin.

F

FIFO buffer: *First-in, first-out buffer.* A portion of memory in which data is stored and then retrieved in the same order in which it was stored. The synchronous serial port has two four-word-deep FIFO buffers: one for its transmit operation and one for its receive operation.

flash memory: Electronically erasable and programmable, nonvolatile (read-only) memory.

G

general-purpose input/output pins: Pins that can be used to accept input signals and/or send output signals but are not linked to specific uses. These pins are the input pin $\overline{\text{BIO}}$, the output pin XF, and the input/output pins IO0, IO1, IO2, and IO3.

global data space: One of the four ’C24x address spaces. The global data space can be used to share data with other processors within a system and can serve as additional data space. See also *local data space*.

GREG: *Global memory allocation register*. A memory-mapped register used for specifying the size of the global data memory. Addresses not allocated by the GREG for global data memory are available for local data memory.

H

hardware interrupt: An interrupt triggered through physical connections with on-chip peripherals or external devices.

$\overline{\text{HOLD}}$: An input signal that allows external devices to request control of the external buses. If an external device drives the $\overline{\text{HOLD}}$/$\overline{\text{INT1}}$ pin low and the CPU sends an acknowledgement at the $\overline{\text{HOLDA}}$ pin, the external device has control of the buses until it drives $\overline{\text{HOLD}}$ high or a nonmaskable hardware interrupt is generated. If $\overline{\text{HOLD}}$ is not used, it should be pulled high.

$\overline{\text{HOLDA}}$: *$\overline{\text{HOLD}}$ acknowledge signal*. An output signal sent to the $\overline{\text{HOLDA}}$ pin by the CPU in acknowledgement of a properly initiated HOLD operation. When $\overline{\text{HOLDA}}$ is low, the processor is in a holding state and the address, data, and memory-control lines are available to external circuitry.

HOLD operation: An operation on the 'C24x that allows for direct memory access of external memory and I/O devices. A HOLD operation is initiated by a $\overline{\text{HOLD}}$/$\overline{\text{INT1}}$ interrupt. When the corresponding interrupt service routine executes an IDLE instruction, the external buses enter the high-impedance state and the $\overline{\text{HOLDA}}$ signal is asserted. The buses return to their normal state, and the HOLD operation is concluded, when the processor exits the IDLE state.

I

$\overline{\text{IACK}}$: See *interrupt acknowledge signal ($\overline{\text{IACK}}$)*.

ICR: See *interrupt control register (ICR)*.

IFR: See *interrupt flag register (IFR)*.

immediate addressing: One of the methods for obtaining data values used by an instruction; the data value is a constant embedded directly into the instruction word; data memory is not accessed.

immediate operand/immediate value: A constant given as an operand in an instruction that is using immediate addressing.

IMR: See *interrupt mask register (IMR)*.

indirect addressing: One of the methods for obtaining data values used by an instruction. When an instruction uses indirect addressing, data memory is addressed by the current auxiliary register. See also *direct addressing*.

input clock signal: See *CLKIN*.

input shifter: A 16- to 32-bit left barrel shifter that shifts incoming 16-bit data from 0 to 16 positions left relative to the 32-bit output.

instruction-decode phase: The second phase of the pipeline; the phase in which the instruction is decoded. See also *pipeline*; *instruction-fetch phase*; *operand-fetch phase; instruction-execute phase*.

instruction-execute phase: The fourth phase of the pipeline; the phase in which the instruction is executed. See also *pipeline*; *instruction-fetch phase*; *instruction-decode phase*; *operand-fetch phase*.

instruction-fetch phase: The first phase of the pipeline; the phase in which the instruction is fetched from program-memory. See also *pipeline*; *instruction-decode phase*; *operand-fetch phase; instruction-execute phase*.

instruction register (IR): A 16-bit register that contains the instruction being executed.

instruction word: A 16-bit value representing all or half of an instruction. An instruction that is fully represented by 16 bits uses one instruction word. An instruction that must be represented by 32 bits uses two instruction words (the second word is a constant).

$\overline{\text{INT1}}$–$\overline{\text{INT3}}$: Three external pins used to generate general-purpose hardware interrupts.

internal interrupt: A hardware interrupt caused by an on-chip peripheral.

interrupt: A signal sent to the CPU that (when not masked or disabled) forces the CPU into a subroutine called an interrupt service routine (ISR). This signal can be triggered by an external device, an on-chip peripheral, or an instruction (INTR, NMI, or TRAP).

interrupt acknowledge signal ($\overline{\text{IACK}}$): An output signal that indicates an interrupt has been received and that the program counter is fetching the interrupt vector that will force the processor into the appropriate interrupt service routine.

interrupt control register (ICR): A 16-bit register used to differentiate $\overline{\text{HOLD}}$ and $\overline{\text{INT1}}$ and to individually mask and flag $\overline{\text{INT2}}$ and $\overline{\text{INT3}}$.

interrupt flag register (IFR): A 16-bit memory-mapped register that indicates pending interrupts. Read the IFR to identify pending interrupts and write to the IFR to clear selected interrupts. Writing a 1 to any IFR flag bit clears that bit to 0.

interrupt latency: The delay between the time an interrupt request is made and the time it is serviced.

interrupt mask register (IMR): A 16-bit memory-mapped register used to mask external and internal interrupts. Writing a 1 to any IMR bit position enables the corresponding interrupt (when INTM = 0).

interrupt mode bit (INTM): Bit 9 in status register ST0; either enables all maskable interrupts that are not masked by the IMR or disables all maskable interrupts.

interrupt service routine (ISR): A module of code that is executed in response to a hardware or software interrupt.

interrupt trap: See *interrupt service routine (ISR)*.

interrupt vector: A branch instruction that leads the CPU to an interrupt service routine (ISR).

interrupt vector location: An address in program memory where an interrupt vector resides. When an interrupt is acknowledged, the CPU branches to the interrupt vector location and fetches the interrupt vector.

INTM bit: See *interrupt mode bit (INTM)*.

I/O-mapped register: One of the on-chip registers mapped to addresses in I/O (input/output) space. These registers, which include the registers for the on-chip peripherals, must be accessed with the IN and OUT instructions. See also *memory-mapped register*.

IR: See *instruction register (IR).*

$\overline{\text{IS}}$: *I/O space select pin*. The 'C24x asserts $\overline{\text{IS}}$ to indicate an access to external I/O space.

ISR: See *interrupt service routine (ISR)*.

ISWS: *I/O-space wait-state bit(s).* A value in the wait-state generator control register (WSGR) that determines the number of wait states applied to reads from and writes to off-chip I/O space.

L

latch phase: The phase of a CPU cycle during which internal values are held constant. See also *logic phase*; *CLKOUT1*.

local data space: The portion of data-memory addresses that are not allocated as global by the global memory allocation register (GREG). If none of the data-memory addresses are allocated for global use, all of data space is local. See also *global data space*.

logic phase: The phase of a CPU cycle during which internal values are changed. See also *latch phase*; *CLKOUT1*.

long-immediate value: A 16-bit constant given as an operand of an instruction that is using immediate addressing.

LSB: *Least significant bit.* The lowest order bit in a word. When used in plural form (LSBs), refers to a specified number of low-order bits, beginning with the lowest order bit and counting to the left. For example, the four LSBs of a 16-bit value are bits 0 through 3. See also *MSB*.

M

machine cycle: See *CPU cycle*.

maskable interrupt: A hardware interrupt that can be enabled or disabled through software. See also *nonmaskable interrupt*.

master clock output signal: See *CLKOUT1*.

master phase: See *logic phase*.

memory-mapped register: One of the on-chip registers mapped to addresses in data memory. See also *I/O-mapped register*.

microcomputer mode: A mode in which the on-chip ROM or flash memory is enabled. This mode is selected with the MP/$\overline{\text{MC}}$ pin. See also *MP/$\overline{\text{MC}}$ pin*; *microprocessor mode*.

microprocessor mode: A mode in which the on-chip ROM or flash memory is disabled. This mode is selected with the MP/$\overline{\text{MC}}$ pin. See also *MP/$\overline{\text{MC}}$ pin*; *microcomputer mode*.

microstack (MSTACK): A register used for temporary storage of the program counter (PC) value when an instruction needs to use the PC to address a second operand.

MIPS: Million instructions per second.

MODE bit: Bit 4 of the interrupt control register (ICR); determines whether the $\overline{\text{HOLD}}$/$\overline{\text{INT1}}$ pin is only negative-edge sensitive or both negative- and positive-edge sensitive.

MP/$\overline{\text{MC}}$ pin: A pin that indicates whether the processor is operating in microprocessor mode or microcomputer mode. MP/$\overline{\text{MC}}$ high selects microprocessor mode; MP/$\overline{\text{MC}}$ low selects microcomputer mode.

MSB: *Most significant bit.* The highest order bit in a word. When used in plural form (MSBs), refers to a specified number of high-order bits, beginning with the highest order bit and counting to the right. For example, the eight MSBs of a 16-bit value are bits 15 through 8. See also *LSB*.

MSTACK: See *microstack*.

multiplier: A part of the CPU that performs 16-bit × 16-bit multiplication and generates a 32-bit product. The multiplier operates using either signed or unsigned 2s-complement arithmetic.

N

next AR: See *next auxiliary register*.

next auxiliary register: The register that is pointed to by the auxiliary register pointer (ARP) when an instruction that modifies ARP is finished executing. See also *auxiliary register*; *current auxiliary register*.

$\overline{\text{NMI}}$: A hardware interrupt that uses the same logic as the maskable interrupts but cannot be masked. It is often used as a soft reset. See also *maskable interrupt*; *nonmaskable interrupt*.

nonmaskable interrupt: An interrupt that can be neither masked by the interrupt mask register (IMR) nor disabled by the INTM bit of status register ST0.

NPAR: *Next program address register.* Part of the program-address generation logic. This register provides the address of the next instruction to the program counter (PC), the program address register (PAR), the micro stack (MSTACK), or the stack.

O

operand: A value to be used or manipulated by an instruction; specified in the instruction.

operand-fetch phase: The third phase of the pipeline; the phase in which an operand or operands are fetched from memory. See also *pipeline*; *instruction-fetch phase*; *instruction-decode phase; instruction-execute phase*.

output shifter: 32- to 16-bit barrel left shifter. Shifts the 32-bit accumulator output from 0 to 7 bits left for quantization management, and outputs either the 16-bit high or low half of the shifted 32-bit data to the data write bus (DWEB).

OV bit: *Overflow flag bit.* Bit 12 of status register ST0; indicates whether the result of an arithmetic operation has exceeded the capacity of the accumulator.

overflow (in a register): A condition in which the result of an arithmetic operation exceeds the capacity of the register used to hold that result.

overflow mode: The mode in which an overflow in the accumulator causes the accumulator to be loaded with a preset value. If the overflow is in the positive direction, the accumulator is loaded with its most positive number. If the overflow is in the negative direction, the accumulator is filled with its most negative number.

OVM bit: *Overflow mode bit.* Bit 11 of status register ST0; enables or disables overflow mode. See also *overflow mode*.

P

PAB: See *program address bus (PAB).*

PAR: *Program address register.* A register that holds the address currently being driven on the program address bus for as many cycles as it takes to complete all memory operations scheduled for the current machine cycle.

PC: See *program counter (PC).*

PCB: *Printed circuit board*.

pending interrupt: A maskable interrupt that has been successfully requested but is awaiting acknowledgement by the CPU.

pipeline: A method of executing instructions in an assembly line fashion. The ’C24x pipeline has four independent phases. During a given CPU cycle, four different instructions can be active, each at a different stage of completion. See also *instruction-fetch phase*; *instruction-decode phase*; *operand-fetch phase; instruction-execute phase*.

PLL: Phase lock loop circuit.

PM bits: See *product shift mode bits (PM).*

power-down mode: The mode in which the processor enters a dormant state and dissipates considerably less power than during normal operation. This mode is initiated by the execution of an IDLE instruction. During a power-down mode, all internal contents are maintained so that operation continues unaltered when the power-down mode is terminated. The contents of all on-chip RAM also remains unchanged.

PRDB: See *program read bus (PRDB).*

PREG: See *product register (PREG).*

product register (PREG): A 32-bit register that holds the results of a multiply operation.

product shifter: A 32-bit shifter that performs a 0-, 1-, or 4-bit left shift, or a 6-bit right shift of the multiplier product based on the value of the product shift mode bits (PM).

product shift mode: One of four modes (no-shift, shift-left-by-one, shift-left-by-four, or shift-right-by-six) used by the product shifter.

product shift mode bits (PM): Bits 0 and 1 of status register ST1; they identify which of four shift modes (no-shift, left-shift-by-one, left-shift-by-four, or right-shift-by-six) will be used by the product shifter.

program address bus (PAB): A 16-bit internal bus that provides the addresses for program-memory reads and writes.

program-address generation logic: Logic circuitry that generates the addresses for program memory reads and writes, and an operand address in instructions that require two registers to address operands. This circuitry can generate one address per machine cycle. See also *data-address generation logic.*

program control logic: Logic circuitry that decodes instructions, manages the pipeline, stores status of operations, and decodes conditional operations.

program counter (PC): A register that indicates the location of the next instruction to be executed.

program read bus (PRDB): A 16-bit internal bus that carries instruction code and immediate operands, as well as table information, from program memory to the CPU.

$\overline{\text{PS}}$: *Program select pin*. The ’C24x asserts $\overline{\text{PS}}$ to indicate an access to external program memory.

PSLWS: *Lower program-space wait-state bits.* A value in the wait-state generator control register (WSGR) that determines the number of wait states applied to reads from and writes to off-chip lower program space (addresses 0000h–7FFFh). See also *PSUWS*.

PSUWS: *Upper program-space wait-state bits.* A value in the wait-state generator control register (WSGR) that determines the number of wait states applied to reads from and writes to off-chip upper program space (addresses 8000h–FFFFh). See also *PSLWS*.

R

RAMEN: *RAM enable pin*. This pin enables or disables on-chip single-access RAM.

$\overline{\text{RD}}$: *Read select pin*. The ’C24x asserts $\overline{\text{RD}}$ to request a read from external program, data, or I/O space. $\overline{\text{RD}}$ can be connected directly to the output enable pin of an external device.

READY: *External device ready pin*. Used to create wait states externally. When this pin is driven low, the ’C24x waits one CPU cycle and then tests READY again. After READY is driven low, the ’C24x does not continue processing until READY is driven high.

repeat counter (RPTC): A 16-bit register that counts the number of times a single instruction is repeated. RPTC is loaded by an RPT instruction.

reset: A way to bring the processor to a known state by setting the registers and control bits to predetermined values and signaling execution to start at address 0000h.

reset pin ($\overline{\text{RS}}$): A pin that causes a reset.

reset vector: The interrupt vector for reset.

return address: The address of the instruction to be executed when the CPU returns from a subroutine or interrupt service routine.

RPTC: See *repeat counter (RPTC).*

$\overline{\text{RS}}$: *Reset pin*. When driven low, causes a reset on any ’C24x device.

R/$\overline{\text{W}}$: *Read/write pin.* Indicates the direction of transfer between the ’C24x and external program, data, or I/O space.

S

SARAM: *Single-access RAM*. RAM that can accessed (read from or written to) once in a single CPU cycle.

scratch-pad RAM: Another name for DARAM block B2 in data space (32 words).

short-immediate value: An 8-, 9-, or 13-bit constant given as an operand of an instruction that is using immediate addressing.

sign bit: The MSB of a value when it is seen by the CPU to indicate the sign (negative or positive) of the value.

sign extend: Fill the unused high order bits of a register with copies of the sign bit in that register.

sign-extension mode (SXM) bit: Bit 10 of status register ST1; enables or disables sign extension in the input shifter. It also differentiates between logic and arithmetic shifts of the accumulator.

single-access RAM: See *SARAM*.

slave phase: See *latch phase*.

software interrupt: An interrupt caused by the execution of an INTR, NMI, or TRAP instruction.

software stack: A program control feature that allows you to extend the hardware stack into data memory with the PSHD and POPD instructions. The stack can be directly stored and recovered from data memory, one word at time. This feature is useful for deep subroutine nesting or protection against stack overflow.

ST0 and ST1: See *status registers ST0 and ST1*.

stack: A block of memory reserved for storing return addresses for subroutines and interrupt service routines. The ’C24x stack is 16 bits wide and eight levels deep.

status registers ST0 and ST1: Two 16-bit registers that contain bits for determining processor modes, addressing pointer values, and indicating various processor conditions and arithmetic logic results. These registers can be stored into and loaded from data memory, allowing the status of the machine to be saved and restored for subroutines.

$\overline{\text{STRB}}$: *External access active strobe*. The ’C24x asserts $\overline{\text{STRB}}$ during accesses to external program, data, or I/O space.

SXM bit: See *sign-extension mode bit (SXM).*

T

TC bit: *Test/control flag bit*. Bit 11 of status register ST1; stores the results of test operations done in the central arithmetic logic unit (CALU) or the auxiliary register arithmetic unit (ARAU). The TC bit can be tested by conditional instructions.

temporary register (TREG): A 16-bit register that holds one of the operands for a multiply operation; the dynamic shift count for the LACT, ADDT, and SUBT instructions; or the dynamic bit position for the BITT instruction.

TOS: *Top of stack.* Top level of the 8-level last-in, first-out hardware stack.

TREG: See *temporary register (TREG).*

TTL: *Transistor-to-transistor logic*.

V

vector: See *interrupt vector*.

vector location: See *interrupt vector location*.

W

wait state: A CLKOUT1 cycle during which the CPU waits when reading from or writing to slower external memory.

wait-state generator: An on-chip peripheral that generates a limited number of wait states for a given off-chip memory space (program, data, or I/O). Wait states are set in the wait-state generator control register (WSGR).

$\overline{\text{WE}}$: *Write enable pin*. The ’C24x asserts $\overline{\text{WE}}$ to request a write to external program, data, or I/O space.

WSGR: *Wait-state generator control register.* This register, which is mapped to I/O memory, controls the wait-state generator.

X

XF bit: *XF-pin status bit.* Bit 4 of status register ST1 that is used to read or change the logic level on the XF pin.

XF pin: *External flag pin*. A general-purpose output pin whose status can be read or changed by way of the XF bit in status register ST1.

Z

zero fill: A way to fill the unused low or high order bits in a register by inserting 0s.

Appendix E

Summary of Updates in This Document

This appendix provides a summary of the updates in this version of the document. Updates within paragraphs appear in a **bold typeface**.

Page:	Change or Add:
4–5	Changed Figure 4–2, *Data Memory Map for 'C24x.* The word count for DARAM (B0) and DARAM (B1) was changed to : 256/512, device dependent.
6–5	Changed Figure 6–2, *System Control Register* (*SYSCR*). The CLKSRC1 and CLKSRC0 read/write notation was changed to : R/W–1*, not affected by reset. Set to 1 by power-on reset.
6–37	Changed Figure 6–10, *External Interrupt Control Registers.* The alpha suffix CR was added to each register designation. At address 7072h, Type A was changed to Type A–NMI.
8–49	Changed the text note in the center of the page to : BLDD does not work with core memory-mapped registers **such as IMR, IFR, and GREG**.

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