

An Implementation of FFT, DCT, and Other Transforms on the TMS320C30

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An Implementation of FFT, DCT, and Other Transforms on the TMS320C30

Abstract

This book describes the several types of transforms and related algorithms used on the TMS320C30 family of digital signal processors. These include:

- ❑ The Fast Fourier Transforms (FFTs)
 - the complex radix-2 FFT
 - the complex radix-4 FFT
 - the real valued radix-2
- ❑ The Discrete Hartley Transform (DHT)
- ❑ The Discrete Cosine Transform (DCT)

The book contains:

- ❑ A description of transforms and their implementation on the TMS320C30 family of digital signal processors.
- ❑ A description and comparison of the different kinds of transforms: the FFTs, the Hartley transform and the Cosine transform
- ❑ A description of the features of the TMS320C30 that allow the efficient implementation of these algorithms
- ❑ Outlines of specific descriptions of implementations, transforms and TMS320C30 C Compiler facts
- ❑ Implementation issues
- ❑ Several graphics and tables detailing
 - Forms and flowgraphs of FFTs



- Memory requirements for FFT and Hartley transforms
- Differences in FFT and DCT timing

The end of the book contains 17 appendices with actual TMS320C30 source code for performing transforms.



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This report describes the implementation of several Fast Fourier Transforms (FFTs) and related algorithms on the TMS320C30. The TMS320C30 is the first device in the third generation of 32-bit floating-point Digital Signal Processors (DSPs) in the Texas Instruments TMS320 family. The algorithms considered here are the complex radix-2 FFT, the complex radix-4 FFT, the real-valued radix-2 FFT (both forward and inverse transforms), the Discrete Hartley Transform (DHT), and the Discrete Cosine Transform (DCT). These transforms have many applications, such as in image processing, sonar, and radar.

The introduction briefly describes transforms and their implementation on the TMS320 family of processors. Next, the different kinds of FFTs (including the real FFT), the closely-related Hartley transform, and the Cosine transform are described and compared. This is followed by a description of the TMS320C30 features that permit efficient implementations of these algorithms. Then, specific implementations, transforms, and TMS320C30 C Compiler facts are outlined. Finally, the report discusses some implementation issues, and the appendices list actual TMS320C30 code for performing transforms.

The powerful architecture and instruction set of the TMS320C30 permit flexible and compact coding of the algorithms in assembly language while preserving close correspondence to a high-level language implementation. The efficiency of the architecture and the speed of the device make faster realization of real and complex transforms possible. With the availability of a C compiler, these routines can be put in C-callable form and used as faster versions of FFT C functions.

Introduction

The Fast Fourier Transform (FFT) is an important tool used in Digital Signal Processing (DSP) applications. Its development by Cooley and Tuckey gave impetus to the establishment of DSP as an independent discipline. The well-structured form of the FFT has also made it one of the benchmarks in assessing the performance of number-crunching devices and systems.

In recent years, because of the popularity of this signal-processing tool, there have been efforts to improve its performance by advances both at the algorithmic level and in hardware implementation. Researchers have been developing efficient algorithms to increase the execution speed of FFTs while keeping requirements for memory size low. On the other hand, developers of VLSI systems are including features in their designs that improve system performance for applications requiring FFTs. In particular, single-chip programmable DSP devices, currently available or under development, can realize FFTs with speeds that allow the implementation of very complex systems in realtime.

The Texas Instruments TMS320 family consists of five generations of programmable digital signal processors. The TMS32010 introduced the first generation, which today encompasses more than twelve devices with various speeds, interfacing capabilities, and price/performance combinations. FFT implementations on the TMS32010 can be found in the appendix of the book by Burrus and Parks [1].

The second-generation TMS320 devices (the TMS32020, the TMS320C25, and their spinoffs) enhanced the architecture and speed capabilities of the first generation. Examples of FFT programs implemented on the TMS32020 can be found in an application report in the book *Digital Signal Processing Applications with the TMS320 Family* [2]. Such programs are easily extended to the TMS320C25 because of the code compatibility between devices.

The architectural and speed improvements on the processors from one generation to the next have made the FFT computation faster and the programming easier. These advantages have reached a new high level in the third generation. The TMS320C30 is the first device in the third generation, and this report examines implementation of the FFT algorithms on it. The fourth generation (TMS320C4x) is a new set of floating-point devices, while the fifth generation (TMS320C5x) is a continuation of the fixed-point devices. Since software compatibility is maintained within the fixed-point and the floating-point devices, the existing FFT implementations will also be applicable to these new generations.

The Fourier Transform of an analog signal $x(t)$, given as

$$X(\omega) = \int_{-\infty}^{\infty} x(t) e^{-j\omega t} dt \quad (1)$$

determines the frequency content of the signal $x(t)$. In other words, for every frequency, the Fourier transform $X(\omega)$ determines the contribution of a sinusoid of that frequency in the composition of the signal $x(t)$. For computations on a digital computer, the signal $x(t)$ is sampled at discrete-time instants. If the input signal is digitized, a sequence of numbers $x(n)$ is available instead of the continuous-time signal $x(t)$. Then, the Fourier transform takes the form

$$X(e^{j\omega}) = \sum_{n=-\infty}^{\infty} x(n) e^{-j\omega n} \quad (2)$$

The resulting transform $X(e^{j\omega})$ is a periodic function of ω , and it needs to be computed for only one period. The actual computation of the Fourier transform of a stream of data presents difficulties because $X(e^{j\omega})$ is a continuous function in ω . Since the transform must be computed at discrete points, the properties of the Fourier transform led to the definition of the *Discrete Fourier Transform* (DFT), given by

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-j \frac{2\pi kn}{N}} \quad (3)$$

When $x(n)$ consists of N points $x(0), x(1), \dots, x(N-1)$, the frequency-domain representation is given by the set of N points $X(k), k=0, 1, \dots, N-1$. Equation (3) is often written in the form

$$X(k) = \sum_{n=0}^{N-1} x(n) W_N^{nk} \quad (4)$$

where $W_N^{nk} = e^{-j 2\pi nk / N}$. The factor W_N is sometimes referred to as the *twiddle factor*. A detailed description of the DFT can be found in references [1,3,4]. The computational requirements of the DFT increase rapidly with increasing block size N , having an impact on the real-time system performance. This problem was alleviated with the development of special fast algorithms, collectively known as Fast Fourier Transform (FFT). With an FFT, the computational burden increases much less rapidly with N , and for any given N , the FFT computational load, measured in terms of required multiplications and additions, is smaller than a brute-force computation of the DFT.

The definition of the FFT is identical to the DFT: only the method of computation differs. To achieve the efficiency of an FFT, it is important that N be a highly composite number. Typically, the length N of the FFT is a power of 2: $N = 2^M$, and the whole algorithm breaks down into a repeated application of an elementary transform known as a *butterfly*. If N is not a power of 2, the sequence $x(n)$ is appended with enough zeroes to make the total length a power of 2. Again, references [1,3,4] contain a detailed development of the FFT. Reference [2] also discusses the same topic.

Different Forms of the FFT

Over the years, researchers have developed different forms of FFT for more efficient computation. Special cases, such as those in which the input is a sequence of real numbers, have been investigated, and even more sophisticated algorithms have been developed. The general form of the FFT *butterfly* is given in Figure 1.

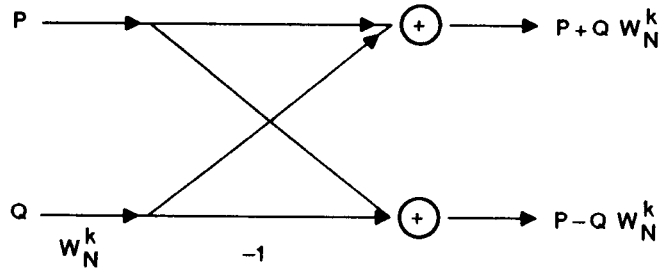


Figure 1. Radix-2 Butterfly for Decimation in Time

If the inputs to the butterfly are the two complex numbers P and Q , the outputs will be the complex numbers P' and Q' , such that

$$P' = P + Q W_N^k \quad (5)$$

and

$$Q' = P - Q W_N^k \quad (6)$$

The quantities P , Q , and P' , Q' represent different points in the array being transformed, and they may or may not occupy adjacent locations in that array. In an in-place computation, the result P' will overwrite P , and Q' will overwrite Q . W_N^k represents again the twiddle factor, and its exponent is determined by the location of the corresponding butterfly in the FFT algorithm.

Figure 2 shows an alternate form of the same FFT butterfly.

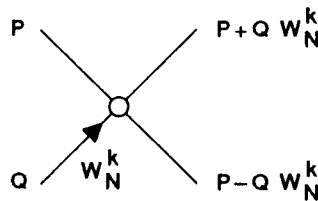


Figure 2. Alternate Form of Radix-2 Butterfly for Decimation in Time.

Although the notation is now less descriptive, it creates a clearer picture when several butterflies are put together to form an FFT. Using the first notation, Figure 3 is the flowgraph of an 8-point FFT example.

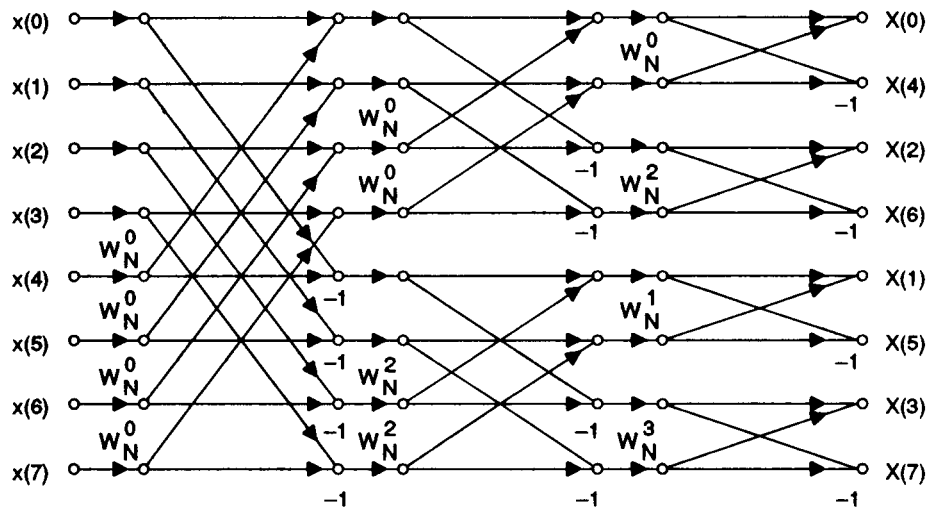


Figure 3. Example of 8-Point FFT with Decimation in Time.

Note that the input sequence $x(n)$ is in the correct order, while the output $X(k)$ is scrambled. Actually, this scrambling occurs in a very systematic way, called bit-reversed order: If you express the indices of a scrambled sequence in binary and you reverse this number, the result is the order that this particular point occupies. For instance, $X(3)$ occupies the sixth position in the output (when counting from the zero position). In binary form, $3_{10} = 011_2$, and if bit-reversed, you get $110_2 = 6_{10}$, which is the position that $X(3)$ occupies. It turns out that the third position is occupied by $X(6)$, and to restore the correct order at the output, you need only to swap these two numbers.

The same procedure can be repeated with all the scrambled numbers not occupying the position that their index suggests. If the input sequence $x(n)$ is rearranged to appear in bit-reversed form, the output $X(k)$ appears in the correct order, as shown in Figure 4.

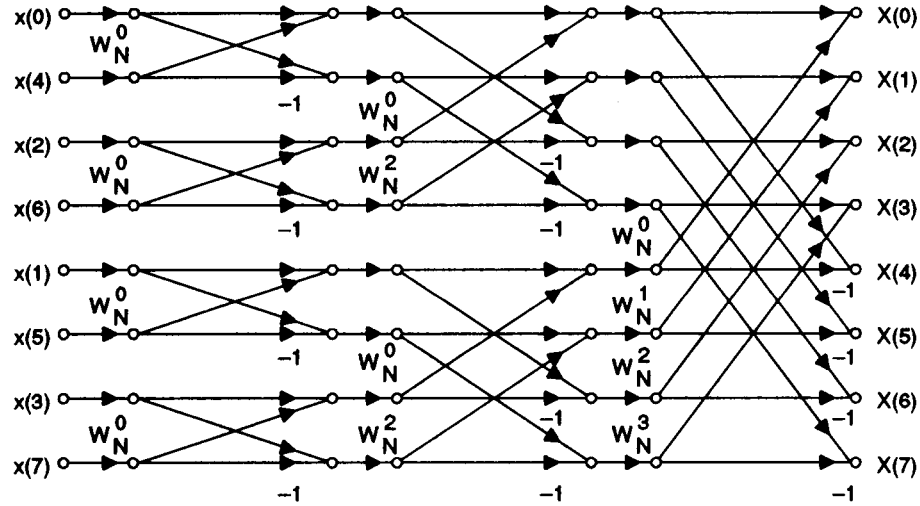


Figure 4. Alternate Form of 8-Point FFT with Decimation in Time. The Input Is in Bit-Reversed Order and the Output Is in the Correct Order.

Since the only difference between Figures 3 and 4 is a rearrangement of the butterflies, the computational load and the final results are identical. In terms of implementation, this rearrangement means that the nesting of the two innermost loops in the FFT routine is interchanged.

The butterflies and the FFT configurations presented thus far implement the FFT with a *decimation in time*. This terminology essentially describes a way of grouping the terms of the DFT definition; see Equation (3). An alternative way of grouping the DFT terms together is called *decimation in frequency*. Figures 5 and 6 show the same example of an 8-point FFT: Figure 5 with the input in correct order and the output in bit-reversed order, and Figure 6 vice-versa, and using the decimation in frequency (DIF).

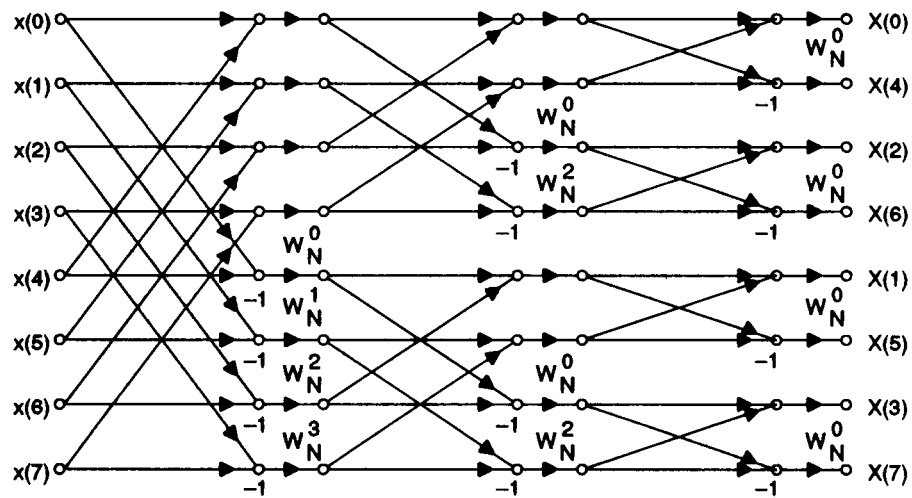


Figure 5. Example of an 8-Point FFT with Decimation in Frequency.

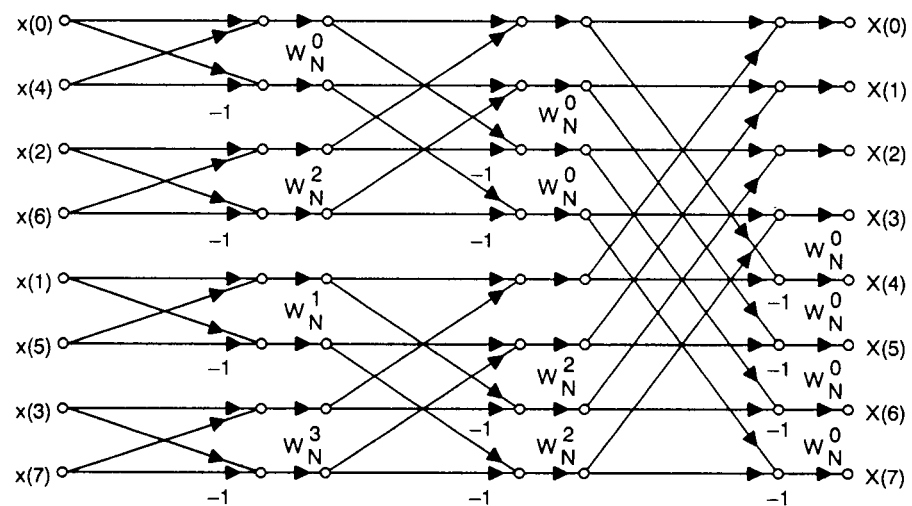


Figure 6. Alternate Form of 8-Point FFT with Decimation in Frequency. The Input Is in Bit-Reversed Order and the Output Is in the Correct Order

Pictorially, the difference between decimation in time and decimation in frequency is that the twiddle factor appears at the input of the butterfly in the first, and at the output in the second. Otherwise, the two methods are identical in terms of results. However, depending on what is the most convenient order of getting the twiddle factors and where the longest-span butterfly appears, you may prefer one method over the other.

The butterfly shown in Figure 1 (or Figure 2) is the smallest element in a radix-2 FFT. The radix of the FFT represents the number of inputs that are combined in a butterfly. The Fast Fourier Transform is usually explained around the radix-2 algorithm for conceptual simplicity. If, however, higher-order radices are used, more computational savings can be achieved. These savings increase with the radix, but there is very little improvement above radix 4. That's why the radix-2 and radix-4 FFTs are the most commonly used algorithms.

In radix-4 FFT, each butterfly has 4 inputs and 4 outputs, essentially combining two stages of a radix-2 algorithm in one. Figure 7 shows this combination graphically.

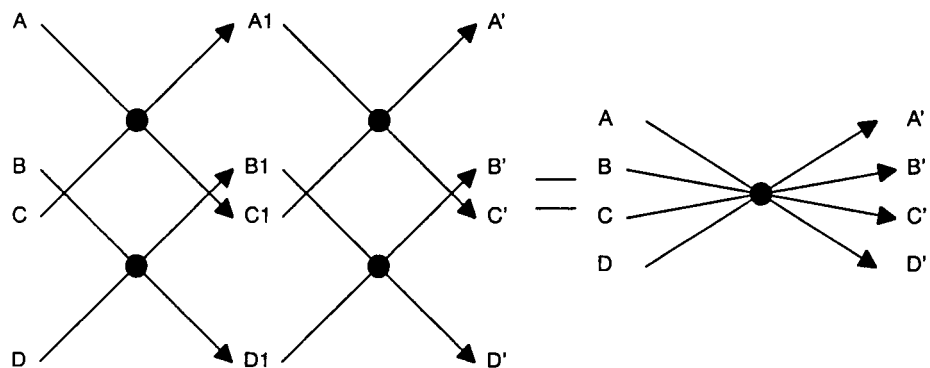


Figure 7. Butterfly for Radix-4, Decimation-in-Time FFT.

Although four radix-2 butterflies are combined into one radix-4 butterfly, the computational load of the latter is less than four times the load of a radix-2 butterfly. Examples of radix-4, 16-point FFTs are shown in Figures 8 and 9 for decimation in time and decimation in frequency, respectively.

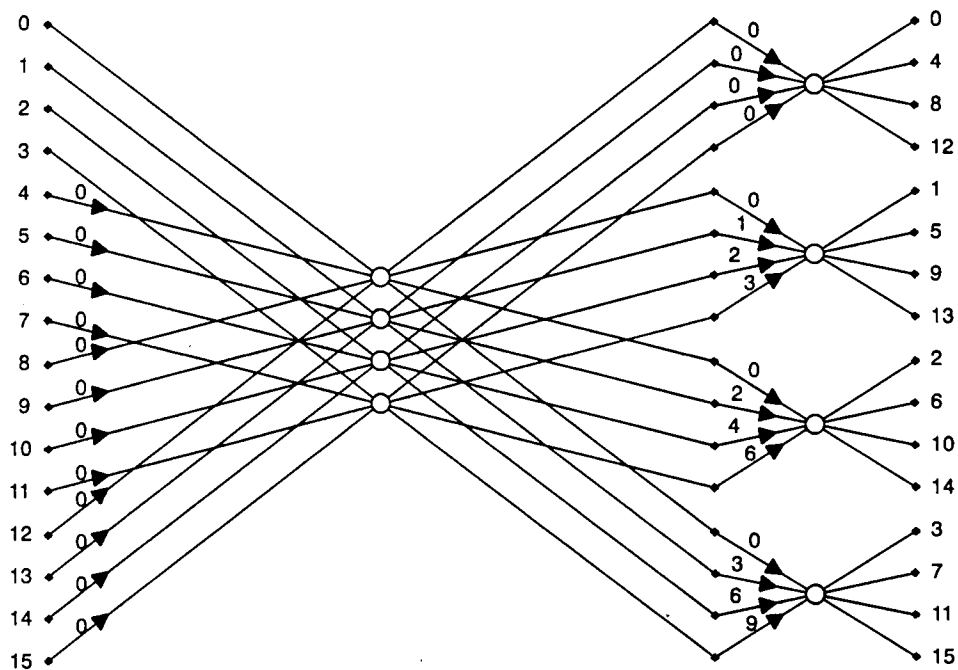


Figure 8. Example of a 16-Point, Radix-4, Decimation-in-Time FFT.

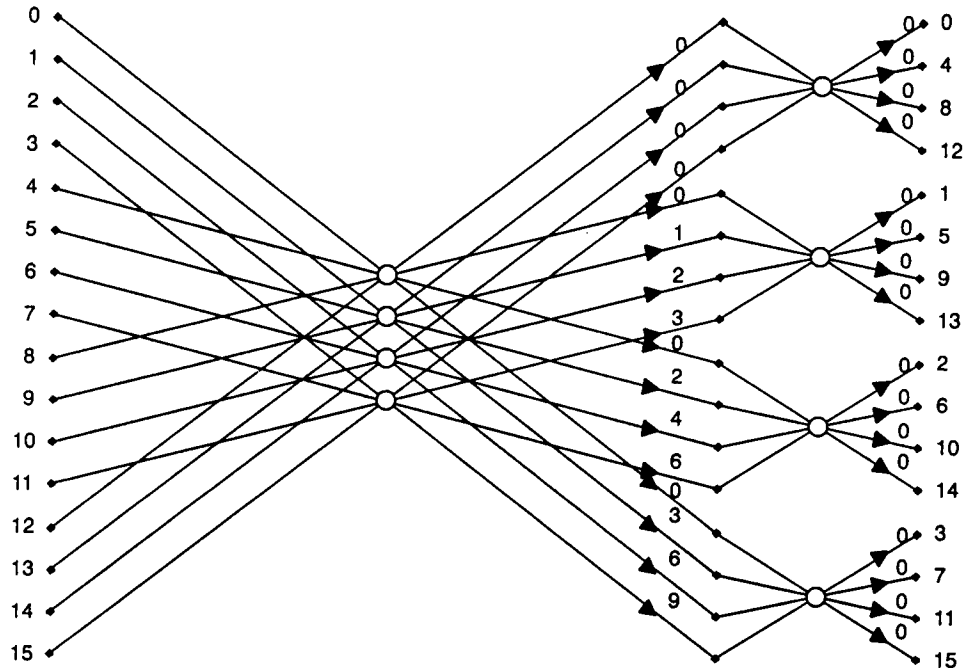


Figure 9. Example of a 16-Point, Radix-4, Decimation-in-Frequency FFT.

These configurations take the incoming sequence in order and produce the frequency-domain result in digit-reversed form. It is a simple matter to rearrange the FFT and have the input in digit-reversed form and the output in order.

Digit reversal is similar to bit reversal, except that the number whose digits are reversed is written in base 4 (equal to the radix) rather than base 2. For example, the output value $X(14)$ in a 16-point, radix-4 FFT occupies position eleven (again starting from zero) because $14_{10} = 32_4$ and, reversing the digits of the number, $23_4 = 11_{10}$. To restore the output to the correct order, the contents of locations with digit-reversed indices should be swapped. However, since the TMS320C30 has a special bit-reversed addressing mode, it is desirable to have the output of the radix-4 computation in bit-reversed rather than digit-reversed form. This is accomplished quite simply if, in each radix-4 butterfly, the two middle output legs are interchanged. That is, whenever the output of the butterfly is the four numbers A' , B' , C' , and D' , instead of storing them in that order, store them in the order A' , C' , B' , and D' , as shown in Figure 10.

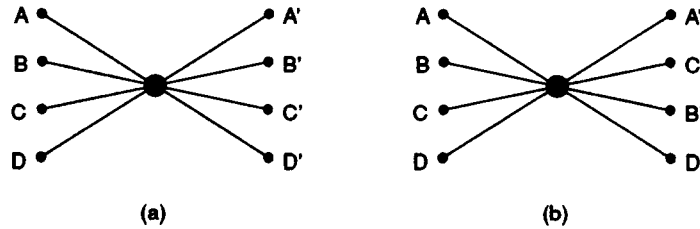


Figure 10. Radix-4 Butterflies. (a) Regularly-Ordered Output, (b) Bit-Reversed Output.

References [5, 6] explain why this simple rearrangement puts the result in bit-reversed order.

Features of the TMS320C30

The TMS320C30 is the first device introduced in the third generation of the TMS320 Digital Signal Processors [7,8]. It has many architectural features that permit very efficient implementation of algorithms. Some of those features pertinent to the FFT implementation are discussed in this section.

The two most salient characteristics of the TMS320C30 device are its high speed (60-ns cycle time) and floating-point arithmetic. The higher speed makes the implementation of real-time application easier than in earlier processors, even when the other architectural advantages are not considered. Each instruction executes in a single cycle under mild pipeline restrictions. The device automatically takes care of any potential conflicts. The pipeline should be observed closely (e.g., using the trace capability of the simulator) only if code optimization for speed is required.

The floating-point capability permits the handling of numbers of high dynamic range without concern for overflows. In FFT programs, in particular, the computed values tend to increase from one stage to the next, as discussed in reference [2]. Then, the fixed-point arithmetic will cause overflows if the incoming numbers are large enough and no provisions are made for scaling. All these considerations are eliminated with the floating-point capability of the TMS320C30. The TMS320C30 performs floating-point arithmetic with the same speed as any fixed point operation; no performance is sacrificed for this feature.

There are eight extended-precision registers, R0—R7, that can be used as accumulators or general-purpose registers, and eight auxiliary registers, AR0—AR7, for addressing and integer arithmetic. For many applications, these registers are sufficient for temporary storage of values, and there is no need to use memory locations. This is the case with the radix-2 FFT algorithm, where no locations are required other than those for the transformation of incoming data to be transformed. Also, arithmetic using these registers greatly increases the programming efficiency. The two index registers, IR0 and IR1, are used for indexing the contents of the auxiliary registers AR0—AR7, thus making the access of the butterfly legs and the twiddle factors easy.

A powerful structure in the TMS320C30 is the block-repeat capability that has the form

```

                RPTB    LABEL
                put instructions here
LABEL          last instruction

```

Whatever occurs after the RPTB instruction and up to the LABEL is repeated one time more than the number included in the repeat counter register, RC. The RC register must be initialized before entering the block-repeat construct. The net effect is that the repeated code behaves as if it were straight-line coded (no penalty for looping), with program size equal to the one in looped code. In this way, the FFT butterfly, being the core of the program, can be implemented in a block-repeat form, thereby saving execution time while preserving the clarity of the program and conserving program space.

A bit-reversed addressing mode is available to eliminate the need for swapping memory locations at the beginning or the end of the FFT (depending on the FFT type). When you use this addressing mode, you access a sequence of data points in bit-reversed order rather than sequentially, and you can recover the points in the correct order during retrieval of the data instead of spending extra cycles to accomplish it in software.

Implementation of Radix-2 and Radix-4 Complex FFTs

Because of the powerful architecture and the instruction set of the TMS320C30, the assembly language program follows closely the flow of a high-level language program; this makes it easy to read and debug. It also keeps the size of the program small and reduces the requirements for program memory. Appendix A presents an example of code for a Radix-2 complex FFT, while Appendix B is a radix-4 complex FFT. The program memory requirements for these programs (as well as others to be discussed later) are given in Table 1.

Table 1. Program Memory Requirements for the Core of the FFT and Hartley Transforms

Routine Type	Program Size
Radix-2, complex FFT	50 words
Radix-4, complex FFT	170 words
Radix-2, real FFT	68 words
Radix-2, real inverse FFT	76 words
Hartley transform	71 words

The numbers in the table correspond only to the core program and do not include the sine/cosine tables for the twiddle factors, any input/output, or any bit-reversing operations. Note also that they are independent of the FFT data size.

The data memory requirements are, of course, dependent on the FFT size. The maximum length of a complex, radix-2 FFT that can be implemented entirely on the internal memory of the TMS320C30 is 1024 points. In the present implementation, the 1024-point radix-4 FFT requires a few more locations (about 7) than are available on-chip.

The code (provided in the appendices) has been written to be independent of the FFT length. The length N , together with the sine/cosine tables for the twiddle factors, should be provided separately to maintain the generic nature of the core FFT program. An example of a file with the sine/cosine tables for a 64-point FFT is given in the Appendix F. Note that the FFT size and the number of stages are declared `.global` in both files (i.e., the main routine and the file with the table) so that the core program gets the actual values during linking.

To reduce the storage requirements of a sine/cosine table, a full sine and a cosine cycle are overlapped. The table stores $5/4$ of a full sine wave, with the cosine table starting with a phase delay of $1/4$ cycle from the sine table. This table size is larger than actually needed, and it is selected merely for testing convenience of the algorithms. The minimum table size for a radix-2 complex FFT includes $1/2$ of a full sine wave, and $1/2$ of a full cosine wave. If these two half waves are combined using the above quarter-cycle phase delay, the minimum table size for this kind of FFT is $3/4$ of a full sine wave. For instance, for a 1024-point FFT, the table can be the first 768 points of a sine wave, where a full cycle would be 1024 points. In the case of a radix-4 complex FFT, the minimum table size should include $3/4$ of a sine and $3/4$ of a cosine wave. Overlapping these requirements, we get the minimum table size of a radix-4 algorithm to be one full sine wave.

An example of a linking file is also included in Appendix F to show how the different segments are assigned. For a complete description of the assembler and linker, consult the corresponding manual [6].

The timing of the FFT routines was done using the cycle-counting capability of the TMS320C30 simulator. For the conversion of the number of cycles into seconds, a cycle time of 50 ns was used. The timing refers only to the core FFT computation, ignoring read-in and write-out requirements, since such requirements are application-dependent. Also, no bit reversal is counted (although it may be included in the program), since it is performed as part of the read-in or read-out. Table 2 gives the timing for the different FFT routines and for the Hartley transform.

Table 2. FFT Timing in Milliseconds[†]

Transform Size	Radix-2 Complex FFT	Radix-4 Complex FFT	Radix-2 Real FFT	Radix-2 Real Inverse FFT	Hartley Transform
64	0.101	0.103	0.047	0.053	0.068
128	0.211	—	0.099	0.110	0.151
256	0.453	0.520	0.215	0.241	0.336
512	0.991	—	0.476	0.535	0.943
1024	2.175	2.533	1.055	1.193	2.025
1024	1.972				

[†]Improvements have been made and are shown in this table. You may obtain the latest code from the BBS, (713) 274-2323.

The last entry in this table represents the timing of the radix-2, DIT routine generated at the University of Erlangen [18] and given in Appendix A. These numbers are typically used for benchmarking.

Implementation of Real FFT

The development of FFT algorithms is centered mostly around the assumption that the input sequence consists of complex numbers (as does the output). This assumption guarantees the generality of the algorithm. However, in a large number of actual applications, the input is a sequence of real numbers. If this condition is taken into consideration, additional computational savings can be achieved because the FFT of a real sequence demonstrates the following symmetries: Assuming that the FFT output $X(k)$ is complex,

$$X(k) = R(k) + j I(k) \quad (7)$$

and that the sequence has length N , $R(k)$ and $I(k)$ should satisfy the following relations:

$$R(k) = R(N-k), \quad k = 1, \dots, N/2-1 \quad (8)$$

$$I(k) = -I(N-k), \quad k = 1, \dots, N/2-1 \quad (9)$$

$$I(0) = I(N/2) = 0. \quad (10)$$

In other words, the real part of the transform is symmetric around zero frequency, while the imaginary part is antisymmetric. Similar conditions hold if the transform is expressed in terms of magnitude and phase.

The savings are due to the fact that not all points need to be computed. Since the not-computed points do not need to be saved either, there are also storage savings. An efficient algorithm for real-valued FFTs is described in [10]. This algorithm was implemented in the present study in such a way that, given the sequence of N real numbers $x(0), x(1), \dots, x(N-1)$, the resulting FFT, consisting of complex numbers, is stored as $R(0), R(1), \dots, R(N/2), I(N/2-1), I(N/2-2), \dots, I(1)$. $R(k)$ and $I(k)$ represent the real and imaginary parts of the complex number $X(k)$. Figure 11 shows the memory arrangement for the FFT. Note that the input to the real FFT should be bit-reversed, but the bit reversal can be done as the data is brought in. With this arrangement, an N -point FFT uses exactly N memory locations. If the full array $X(k)$ is needed, the following relations should be used:

$$X(0) = R(0) \quad (11)$$

$$X(k) = R(k) + j I(k), \quad K = 1, \dots, N/2-1 \quad (12)$$

$$X(N/2) = R(N/2) \quad (13)$$

$$X(k) = R(N-k) - j I(N-k), \quad k = N/2+1, \dots, N-1 \quad (14)$$

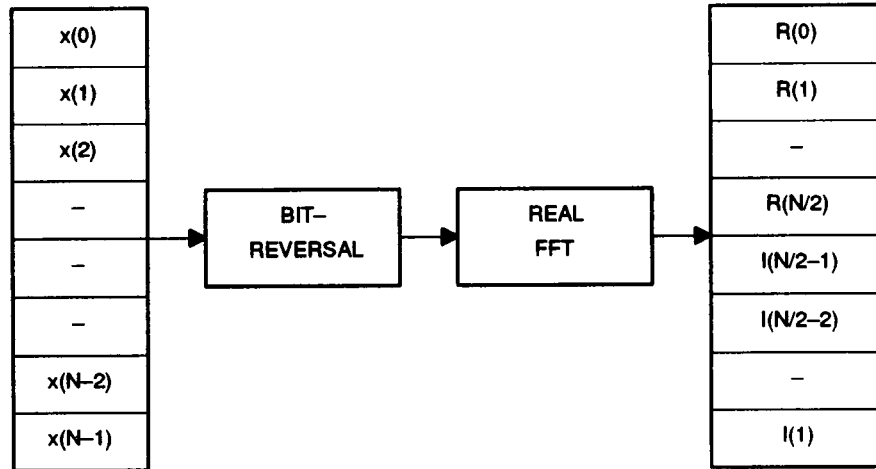


Figure 11. Memory Arrangement of a Real FFT.

It is expected that, in most signal processing applications, there will be no need to reconstruct the full $X(k)$ array and that the output shown in Figure 11 will be sufficient for any further processing.

Appendix C contains TMS320C30 routines implementing a radix-2 real FFT and its inverse. The implementation of the forward transformation is based on the FORTRAN programs contained in [10]. The inverse transformation assumes that the input data are given in the order presented at the output of the forward transformation and produces a time signal in the proper order (i.e., bit-reversing takes place at the end of the program). Viewed another way, the inverse real FFT operates as shown in Figure 11 but with the arrows reversed (and inverse FFT taking the place of the FFT).

The timing for the real-valued FFT (both forward and inverse) is included in Table 2, and the corresponding program sizes are shown in Table 1. As you can see, the real-valued FFT is considerably faster than the corresponding complex FFT because not all the computations need be performed. Furthermore, there are data storage savings because only half the values must be stored. As a result, the maximum length of real-valued FFT that can be implemented on the TMS320C30 without using any external memory is 2048 points. Of course, if all the values are needed, they can be recovered using the symmetry conditions mentioned earlier. To achieve the efficiencies of real FFT and not use any extra memory locations during the computation, the decimation-in-time method is applied [10]. Decimation in time requires the bit-reversal operation in the forward transform to be performed at the beginning of the program rather than at the end. The reverse is true for bit-reversing in the inverse transform.

The Discrete Hartley Transform

Another transform that has attracted attention recently is the Discrete Hartley Transform (DHT)[11, 12]. The DHT is applicable to real-valued signals and is closely related to the real-valued FFT. Comparison of references [10] and [12] describing the implementation of the two algorithms on FORTRAN programs shows that their implementation on the TMS320C30 should be similar. And indeed, this is the case.

The DHT pair is defined for a real-valued sequence $x(n)$, $n = 0, \dots, N-1$, by the following equations:

$$H(k) = \sum_{n=0}^{N-1} x(n) \text{cas}(2\pi k n / N), \quad k=0, \dots, N-1 \quad (15)$$

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} H(k) \text{cas}(2\pi k n / N), \quad k=0, \dots, N-1 \quad (16)$$

where $\text{cas}(x) = \cos(x) + \sin(x)$. The DHT demonstrates a symmetry that is convenient for implementations: The same program can be used for both the forward and the inverse transforms, and the result is correct within a scale factor. Also, the real FFT and the DHT can be derived from each other [12].

A radix-2 Hartley transform was implemented on the TMS320C30, and the corresponding code is included in Appendix D. This code follows the structure of the real FFT in Appendix C. Tables 1 and 2 show the program memory requirements and the timing for the execution of Hartley transforms of different sizes. The sine/cosine table sizes are the same as in the case of a real FFT.

The Discrete Cosine Transform

The Discrete Cosine Transform (DCT), since its introduction in 1974 [13], has gained popularity in speech and image processing applications because of its near-optimal behavior. This discussion is based on the paper by Lee [14]. The DCT code was developed and implemented by Paul Wilhelm of the University of Washington.

If $x(n)$, $n=0, \dots, N-1$ is a time-domain signal and $X(k)$ is the corresponding DCT, $x(n)$ and $X(k)$ are related by the following equations:

$$x(k) = \frac{2}{N} \sum_{n=0}^{N-1} e(n) x(n) \cos \frac{(2k+1)\pi n}{2N} \quad (17)$$

$$x(n) = \sum_{k=0}^{N-1} e(k) X(k) \cos \frac{(2k+1)\pi n}{2N} \quad (18)$$

$$e(0) = 1/\sqrt{2} \quad (19)$$

$$e(k) = 1, \quad \text{for } k \neq 0 \quad (20)$$

Appendix E shows an implementation of the DCT based on the paper by Lee [14]. The appendix contains the algorithms for both the forward and the inverse transformations and an example of a table for a 16-point DCT. Note that, because of the structure of the algorithm, the cosine table needed contains actually the inverses of the cosines (within a scale factor), and it is not stored in the natural order. Instead, it is generated by the following C pseudocode:

```
for [k=2, i=0; k=N/2; k*=2]
  for [j=k/2; j<N/2; j+=k]{
    cos__table[i++] = 1/[2*cos[j*pi/(2*N)]];
    cos__table[i++] = 1/[2*cos[(N-j)*pi/(2*N)]];
  }
cos__table[N-2] = cos[pi/4];
cos__table[N-1] = 2/N;
```


The last entry to the table is not part of the cosine itself; it is a constant that is used by the algorithm, and it is placed at the end of the cosine table for convenience.

Table 3 shows the timing of the forward and inverse transforms for different transform lengths. The difference in the timing between the forward and the inverse transforms is due to the fact that more time was expended to optimize the performance of the inverse transform. Since four of the smallest butterflies were done simultaneously in the center program loop, the minimum permissible array size to be transformed is 8.

Table 3. DCT Timing in Milliseconds

Transform Size	Forward Transform	Inverse Transform
16	0.019	0.017
64	0.875	0.073
128	0.192	0.161
256	0.418	0.347
512	0.912	0.754
1024	1.982	1.652

Other Related Transforms

In addition to the FFT types mentioned earlier (complex, real, decimation-in-time, decimation-in-frequency, etc.), newer forms of the FFT have been developed to reduce the computational load. One of the latest in the literature is the *Split-Radix* FFT. The Split-Radix FFT [16] has the lowest number of multiplies and adds of any known algorithm. It achieves this efficiency by combining certain radix-2 and radix-4 butterflies, but, as a result, the classical concept of FFT stages is lost. The new structure uses a rather complicated indexing scheme, which is the price paid for the reduced multiplies/adds. Since, on the TMS320C30, multiplies/adds are not more expensive computationally than any other operation, the indexing scheme wipes out the gains of the reduced arithmetic. Actually, an implementation of the split-radix FFT showed it to be slower than the radix-2 FFT, one of the main reasons being that the block-repeat structure could no longer be used effectively.

Very often, there is a question on what the different benchmark numbers mean. A useful comparison of execution times for different algorithms on different machines has been made [17]. Table 4 presents a small segment of the resulting information that is relevant to the present discussion: the timing in seconds for the radix-8, mix-radix, and split-radix algorithms that were implemented on various machines. Different operating systems and compilers have been used, as shown. The execution times of Table 4 should be compared with the 0.0010055 s that it takes to implement a 1024-point, radix-2, real FFT on a TMS320C30. As can be seen, the TMS320C30 compares favorably to all the other machines investigated.

Table 4. Execution Times in Seconds for a 1024-Point Real FFT. The Numbers Should Be Compared with 0.001055 s of a 1024-Point Real FFT on the TMS320C30

Machine	Radix-8	Mix-radix	Split-radix
VAX 750 UNIX BSD4.2 f77	0.3634	0.3902	0.3021
VAX 750 UNIX BSD4.2 f77 -O	0.2376	0.2948	0.2089
VAX 750 UNIX BSD4.3 f77	0.2545	0.2600	0.2371
VAX 750 UNIX BSD4.3 f77 -O	0.1825	0.2127	0.1672
VAX 785 ULTRIX f77	0.1046	0.1107	0.1101
VAX 785 ULTRIX f77 -O	0.0796	0.0943	0.0811
VAX 785 VMS FOR/NOOPTM	0.0767	0.0871	0.0975
VAX 785 VMS FOR/OPTM	0.0539	0.0641	0.0633
VAX 8600 VMS FOR/OPTM	0.0217	0.0243	0.0235
MICROVAX VMS FOR/NOOPTM	0.1671	0.1846	0.1864
MICROVAX VMS FOR/OPTM	0.1299	0.1527	0.1419
DEC-10 TOPS-10 FOR/NOOPTM	0.0940	0.1184	0.0991
DEC-10 TOPS-10 FOR/OPTM	0.0885	0.1110	0.0845
CDC 855 FTN5,OPT=0	0.0277	0.0319	0.0338
CDC 855 FTN5,OPT=1	0.0277	0.0316	0.0337
CDC 855 FTN5,OPT=2	0.0182	0.0171	0.0151
CDC 855 FTN5,OPT=3	0.0180	0.0173	0.0150
SUN 3/50 UNIX BSD4.2 f77 -O -f68881	0.2518	0.3365	0.2103
SUN 3/50 UNIX BSD4.2 f77 -f68881	0.2806	0.3897	0.2802
SUN 3/50 UNIX BSD4.2 f77 -O	0.7586	1.047	0.6955
SUN 3/50 UNIX BSD4.2 f77	0.7476	1.029	0.7033
SUN 3/160 UNIX BSD4.2 f77	0.6037	0.6895	0.5660
SUN 3/160 UNIX BSD4.2 f77 -pfa	0.0983	0.1060	0.0946
SUN 3/260 UNIX BSD4.3 f77	0.3689	0.4126	0.3390
SUN 3/260 UNIX BSD4.3 f77 -O	0.3530	0.4142	0.3297
Pyramid 90X UNIX BSD4.2 f77 -O	0.2053	0.2244	0.1416
Pyramid 90X UNIX BSD4.2 f77	0.2206	0.2457	0.1326
HP-1000 21MX-E FTN7X	0.9400	1.248	0.9478
Apple MAC Microsoft FOR	2.6670	3.1600	2.8260
AST PC Microsoft FOR	1.5040	2.0800	1.4630

The TMS320C30 C Compiler

The C compiler for the TMS320C30 permits easy porting of high-level language programs to the DSP device. If the CPU loading of a particular application is not very high, the C compiler can create programs that run on the TMS320C30 in real time. If, however, the result is non-realtime, it may be necessary to use assembly language for more efficient coding.

In most cases, only a portion of the code needs to be written in assembly language. Typically, there are a few code segments where the device spends most of the time and which, when optimized in assembly language, yield the necessary performance improvement. By following the conventions outlined in the run-time environment of the C compiler [15], you can write these time-critical routines in assembly language and call them in a C program. This is also true for the FFT routines. In appendices A, B, and C, the radix-2, radix-4, and real FFT routines mentioned earlier are also put in a C-callable form by adding the necessary interface at the beginning and the end of the code. The tables with the sines and cosines are again assumed to be supplied during link time.

Issues in FFT Implementation

There are many ways of actually implementing the FFT code (and the other transformations), taking into consideration the different possibilities of program locations, the data locations, the ways of input and output, etc. Since it is impractical to cover every possible case, this report has concentrated on a configuration in which the use of external memory is minimized. With the source code and additional explanations provided, you should be able to customize the FFT implementation for a particular application.

Use of External Memory

In these implementations, only on-chip memory was used, and that's why the maximum transform size considered was 1024 points long (2048 for a real transform). Often, though, applications call for use of external memory for program or data or both. When external memory is used, the structure of the code does not change at all; it is only the timing that may be affected.

Fast external memory can be selected so that no wait states are necessary. But even when there are no wait states, accessing external memory may impose some limitations. For instance, you can make only one external memory access in a full cycle, but you can make two accesses of internal memory in each cycle. Also, because of multiplexing of the busses, pipeline conflicts may arise if both program and data are placed on the same external port. Resolution of such conflicts causes extra cycles for the execution. The section on pipelining in the *TMS320C30 User's Guide* explains in detail what kind of potential conflicts may occur.

To minimize or avoid such conflicts, there are some simple steps that the designer can take. The TMS320C30 has three separate memory areas (one on-chip, one accessed by the primary bus, and one accessed by the expansion bus) that can be combined. For instance, the program can be placed on the expansion port and the data on the primary port. Or the data can first be brought into internal memory and then operated upon. Alternatively, the program may be relocated to internal memory. A related approach is to use the cache. All the transforms are implemented as loops that are executed many times. If you activate the on-chip cache after the first access of the code, the instructions execute from the cache instead of the external memory.

If there are additional conflicts, they can typically be resolved by some rearrangement of the code. For instance, consecutively writing to external memory takes two cycles per write. If, however, a write is followed by some internal operation, then the second cycle of the write is transparent, and the actual cost is one cycle.

Bit Reversal

The TMS320C30 has a special form of the indirect addressing mode for the bit-reversing operation that is required at the beginning or the end of an FFT. Through this addressing mode, the scrambled data are accessed in their proper order. This addressing mode works as follows:

Let AR_n ($n=0..7$) be the auxiliary register pointing to the array with scrambled data. The index register $IR0$ contains a number equal to one-half the size of the FFT. Then, after every access of the data, AR_n is incremented by $IR0$ using the construct

$$*AR_n + + [IR0]B$$

This causes the contents of AR_n to be incremented by the contents of $IR0$, but if there is a carry in this incrementing, the carry propagates to the right instead of to the left. The result is the generation of the addresses in a bit-reversed order. The bit-reversed addressing mode works correctly if the array with the data is aligned in memory so that the first memory address is a multiple of the FFT size. This can be achieved if the first memory address has zeros for the last M bits, where $M = \log_2 N$, with N being the FFT size. For example, in the case of a 1024-point FFT, the last 10 bits of the memory address of the first datum should be zeros.

In the implementation of the complex FFT, the output is complex even when the input is real. So, there is a need to consider both the real and the imaginary parts of the data array. The above description of the bit-reversed addressing mode assumed that the real and the imaginary parts are stored as separate arrays in the memory. In this case, each of the arrays (real or imaginary parts) can be accessed as described. However, in most cases (including this report), the real and imaginary points alternate in the same array.

In this arrangement, the following simple modification achieves the same goal: set IRO equal to N instead of $N/2$, and access the N points of the transform. At every access, the auxiliary register is pointing to the real part of the FFT. The imaginary part is located in the next higher location, and it can be easily accessed.

With the bit-reversed addressing mode, the unscrambling of the data can take place when the FFT result is accessed for further processing or for I/O. It is possible, though, that certain applications demand the reordering of the data in the same array. Such a rearrangement can be done very simply for a complex FFT with the following code.

; DO THE BIT-REVERSING EXPLICITLY

```

        LDI    @FFTSIZ,RC      ; RC = FFT SIZE
        SUBI   1,RC           ; RC SHOULD BE ONE LESS THAN DESIRED #
        LDI    @FFTSIZ,IRO    ; IRO = FFT SIZE
        LDI    @INPUT,ARO
        LDI    @INPUT,AR1
*
        RPTB   BITRV
        CMPI   AR1,ARO        ; EXCHANGE LOCATIONS ONLY
        BGE    CONT          ; IF ARO<AR1
        LDF    *ARO,R0
        ||     LDF    *AR1,R1  ; EXCHANGE REAL PARTS
        STF    R0,*AR1
        ||     STF    R1,*ARO
        LDF    *+ARO,R0
        ||     LDF    *+AR1,R1 ; EXCHANGE IMAGINARY PARTS
        STF    R0,*+AR1
        ||     STF    R1,*+ARO
        CONT   NOP    *ARO++(2)
        BITRV  NOP    *AR1++(IRO)B

```

Note that AR1 is pointing to the bit-reversed version of the address contained in ARO. For real-valued FFT, or for FFTs that store the real and the imaginary parts in separate arrays, the real-FFT routine in Appendix C contains a modified example of the above code.

Use of DMA

If the signal to be transformed arrives as a continuous stream of data, the DMA could be used to collect the new data while the data already collected are processed. In this case, the data source address of the DMA points to the memory location corresponding to a serial port, or to another port associated with an external device. The destination is a memory space designated for storage.

There are two ways to use such buffers. One possibility is to designate one buffer as the temporary storage and the other buffer as the working area. When the storage buffer receives the necessary amount of data, the data is transferred to the working area, and the DMA starts refilling the storage buffer. Alternatively, the two buffers are considered equivalent: when the processor finishes processing and outputting the data from one and the DMA has filled the other, the two buffers switch functions; i.e., the DMA starts filling the first buffer while the CPU is processing the data in the buffer just filled.

Test Vector

For testing purposes, a vector with 64 (quasi-random) data points and the corresponding FFT values is given in Appendix F. In this way, if any of the routines is implemented, the test vectors can be used to verify the correct functionality of the routines. Together with the test vectors, Appendix C gives a sine/cosine table for a 64-point transform, and the linking file for such a transform.

Summary

This report examined implementations of fast transforms on the Texas Instruments TMS320C3x floating-point devices. The transforms considered were several forms of the FFT, the Discrete Hartley Transform, and the Discrete Cosine Transform. Because of the powerful architecture of the device, the implementation was done easily and efficiently. It was shown that a TMS320C30 executes the FFTs several times faster than large computers such as VAX and SUN workstations. With the availability of the C compiler, these routines can be put in C-callable form and be used to compute the corresponding transforms efficiently.

Appendices

Appendices A to F contain the TMS320C30 assembly language programs for the different algorithms considered. The contents of the appendices are as follows:

Appendix A: Radix-2 Complex FFT.

composed of

- A1: Generic Program to Do a Looped-Code Radix-2 FFT Computation on the TMS320C30.
- A2: `fft__2` - Radix-2 Complex FFT to Be Called as a C Function.
- A3: Complex, Radix-2 DIT FFT - R2DIT.ASM.
- A4: Complex, Radix-2 DIT FFT - R2DITB.ASM.
- A5: TWID1KBR.ASM - Table with Twiddle Factors for a FFT up to a Length of 1024 Complex Points.

Appendix B: Radix-4 Complex FFT.

composed of

- B1: Generic Program to Do a Looped-Code Radix-4 FFT on the TMS320C30.
- B2: `fft__4` - Radix-4 Complex FFT to Be Called as a C Function.

Appendix C: Radix-2 Real FFT.

composed of

- C1: Generic Program to Do a Radix-2 Real FFT Computation on the TMS320C30.
- C2: `fft__rl` - Radix-2 Real FFT to Be Called as a C Function.
- C3: Generic Program to Do a Radix-2 Real Inverse FFT Computation on the TMS320C30.

Appendix D: Discrete Hartley Transform.

composed of

- D1: Generic Program to Do a Radix-2 Hartley Transform on the TMS320C30.

Appendix E: Discrete Cosine Transform.

composed of

- E1: A Fast Cosine Transform.
- E2: A Fast Cosine Transform (Inverse Transform).
- E3: FCT Cosine Tables File.
- E4: Data File.

Appendix F: Test Vectors, 64-Point Sine Table, Link Command File.
composed of

- F1: Example of a 64-Point Vector to Test the FFT Routines.
- F2: File to Be Linked with the Source Code for a 64-Point,
Radix-4 FFT.
- F3: Link Command File.

The first three appendices contain the code for the radix-2, complex radix-4, and real radix-2 FFT transformations. These routines are given in both the regular form and in a C-callable form. Furthermore, the contents of a file with the twiddle factors are given, as well as an example of a link command file for a 64-point FFT. Note that the source code of these routines can be downloaded from the TI DSP bulletin board (BBS) by calling (713) 274-2323. For questions regarding the BBS, call the TI DSP hotline at (713) 274-2320.

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Mr. Raimund Meyer and Mr. Karl Schwarz (Lehrstuhl für Nachrichtentechnik, University of Erlangen) provided the fast routines of Appendix A to do 1024-point, radix-2, DIT FFT. Mr. Paul Wilhelm of the University of Washington provided the routines for the Fast Cosine Transform (FCT) together with the related explanations and the test vector in Appendix E. Their contributions are gratefully acknowledged.

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Appendix A. Radix-2 Complex FFT

Appendix A1. Generic Program to Do a Looped-Code Radix-2 FFT Computation on the TMS320C30

```

* GENERIC PROGRAM TO DO A LOOPED-CODE RADIX-2 FFT COMPUTATION ON THE
* TMS320C30.
*
* THE PROGRAM IS TAKEN FROM THE BURGESS & PARKS BOOK, P. 111. THE (COMPLET)
* DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE IN-PLACE, BUT THE
* RESULT IS MOVED TO ANOTHER MEMORY SECTION TO DEMONSTRATE THE BIT-REVERSED
* ADDRESSING. THE INTRINSIC FACTORS ARE SUPPLIED IN A TABLE PUT IN A DATA
* SECTION. THIS DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC
* NATURE OF THE PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FFT N AND
* LOG2(N) ARE DEFINED IN A .GLOBAL DIRECTIVE AND SPECIFIED DURING LINKING.
*
* AUTHOR: PANOS E. PAPANTONIOULIS
*
* TEARS INSTRUCTIONS
*
* .GLOBAL FFT
* .GLOBAL N
* .GLOBAL M
* .GLOBAL SINE
*
* .SECTION ".IN",.1024
* .BSS OUTP,1024
*
* .TEXT
*
* INITIALIZE
*
* .WORD FFT
*
* .SPACE 100
*
* FFTSIZ
* .WORD N
* LOGFFT
* .WORD M
* SINTAB
* .WORD SINE
* INPUT
* .WORD INP
* OUTPUT
* .WORD OUTP
*
* FFT:
*
* LDH FFTSIZ,IR1
* LSH -2,IR1
* LDH 0,ARG
* LDH FFTSIZ,IR0
* LSH 1,IR0
* LDH FFTSIZ,R7
* LDH 1,ARG
*
* LDH 1,ARG
*
* OUTER LOOP
*
* MIP
* LDH ++ARG(1)
* LDH INPUT,ARG
*
* ; ENTRY POINT FOR EXECUTION
* ; FFT SIZE
* ; LOG2(N)
* ; ADDRESS OF SINE TABLE
*
* ; MEMORY WITH INPUT DATA
* ; MEMORY WITH OUTPUT DATA
*
* ; STARTING LOCATION OF THE PROGRAM
* ; RESERVE 100 WORDS FOR VECTORS, ETC.
*
* ; COMMAND TO LOAD DATA PAGE POINTER
*
* ; IR1=H/4, POINTER FOR SIN/COS TABLE
* ; ARG HOLDS THE CURRENT STAGE NUMBER
* ; IRC=2*H/4 (BECAUSE OF REAL/IMAG)
* ; R7=H/2
* ; INITIALIZE REPEAT COUNTER OF FIRST
* ; LOOP
* ; INITIALIZE IE INDEX (ARG-1E)
*
* ; CURRENT FFT STAGE
* ; ARG POINTS TO X(1)

```

```

ADDI R7,ARG,ARG2 ; ARG POINTS TO X(L)
LDH R7,RC
SUBI 1,RC
*
* FIRST LOOP
*
* RPTB
* ADDF ARG0,ARG2,R0 ; R0=X(1)+X(L)
* SUBF ARG2++,ARG0++,R1 ; R1=X(1)-X(L)
* ADDF ARG2,ARG0,R2 ; R2=X(1)+X(L)
* SUBF ARG2,ARG0,R3 ; R3=X(1)-X(L)
* STF R2,ARG0-- ; Y(L)=R2 AND...
* STF R3,ARG2-- ; Y(L)=R3
* STF R0,ARG0++(1R0) ; X(L)=R0 AND...
* STF R1,ARG2++(1R0) ; X(L)=R1 AND ARG0,2 = ARG0,2 + 2*H/4
*
* ; IF THIS IS THE LAST STAGE, YOU ARE DONE
*
* CMPI LOGFFT,ARG
* BZD END
*
* MAIN INNER LOOP
*
* LDH 2,ARG1
* RPTB SINTAB,ARG1
* ADDF ARG5,ARG1
* LDH ARG1,ARG0
*
* INLOOP:
* LDH 2,ARG1
* ADDI INPUT,ARG0
* ADDI R7,ARG0,ARG2
* LDH R7,RC
* SUBI 1,RC
* LDF ARG4,R6
*
* ; INIT LOOP COUNTER FOR INNER LOOP
* ; INITIALIZE IA INDEX (ARG+1A)
* ; IA=IA+IE; ARG POINTS TO COSINE
*
* ; INCREMENT INNER LOOP COUNTER
* ; X(1),Y(1) POINTER
* ; X(L),Y(L) POINTER
*
* ; RC SHOULD BE ONE LESS THAN DESIRED #
* ; R6=SIN
*
* SECOND LOOP
*
* RPTB
* SUBF ARG2,ARG0,R2 ; R2=X(1)-X(L)
* SUBF ARG2,ARG0,R1 ; R1=X(1)-Y(L)
* MPYF R2,R6,R0 ; R0=R2*SIN AND...
* MPYF ARG2,ARG0,R3 ; R3=Y(1)+Y(L)
* MPYF R1,ARG4(1R1),R3 ; R3=R1*COS AND...
* STF R3,ARG0++ ; Y(L)=Y(1)+Y(L)
* SUBF R0,R3,R4 ; R4=R1*COS-R2*SIN
* MPYF R1,R6,R0 ; R0=R1*SIN AND...
* ADDF ARG2,ARG0,R3 ; R3=X(1)+X(L)
* MPYF R2,ARG4(1R1),R3 ; R3=R2*COS AND...
* STF R3,ARG0++(1R0) ; X(1)=X(1)+X(L) AND ARG0=ARG0+2*H/4
* ADDF R0,R3,R5 ; R5=R2*COS+R1*SIN
* SUBF R5,ARG2++(1R0) ; X(L)=R2*COS+R1*SIN, INCR ARG2 AND...
* STF R4,ARG2 ; Y(L)=R1*COS-R2*SIN
*
* CMPI R7,ARG1
* BNE INLOOP
*
* ; LOOP BACK TO THE INNER LOOP

```

```

*      LSH      1,AR7      ; INCREMENT LOOP COUNTER FOR NEXT TIME
*
*      LSH      1,AR5      ; IS=2*IE
*      LDI      R7,R0      ; NI=NI
*      LSH      -1,R7      ; NI=NI/2
*      BR       LOOP      ; NEXT FFT STAGE
*
*      ; STORE RESULT OUT USING BIT-REVERSED ADDRESSING
*      ;
*      ; END:
*      LDI      @FFTSIZ,RC ; RC=N
*      SUBI     1,RC       ; RC SHOULD BE ONE LESS THAN DESIRED #
*      LDI      @FFTSIZ,RO ; ; INQ=SIZE OF FFT=N
*      LDI      2,IR1
*      LDI      @INOUT,AR0
*      LDI      @OUTPUT,AR1
*
*      RPTB     BITRV
*      LDF      ++AR0(1),RO
*      LDF      ++AR0++(IR0),R1
*      STF      R0,++AR1(1)
*      STF      R1,++AR1++(IR1)
*
*      BR       SELF      ; BRANCH TO ITSELF AT THE END
*      SELF
*      .END

```

Appendix A2. fft_2—Radix-2 Complex FFT to Be Called as a C Function

```

* NAME:
*   fft_2 --- RADIX-2 COMPLEX FFT TO BE CALLED AS A C FUNCTION.
*
* SYNOPSIS:
*   INT fft_2(IN, N, DATA)
*   INT N      FFT SIZE: N=2**M
*   INT M      NUMBER OF STAGES = LOG2(N)
*   FLOAT DATA ARRAY WITH INPUT AND OUTPUT DATA
*
* DESCRIPTION:
*   GENERIC FUNCTION TO DO A RADIX-2 FFT COMPUTATION ON THE 320C30.
*   THE DATA ARRAY IS 2M-LONG, WITH REAL AND IMAGINARY VALUES ALTERNATING.
*   THE PROGRAM IS BASED ON THE FORTRAN PROGRAM IN THE BURRUS AND PARKS
*   BOOK, P. 111.
*
*   THE COMPUTATION IS DONE IN PLACE, AND THE ORIGINAL DATA IS DESTROYED.
*   BIT REVERSAL IS IMPLEMENTED AT THE END OF THE FUNCTION. IF THIS IS NOT
*   NECESSARY, THIS PART CAN BE COMMENTED OUT.
*
*   THE SINE/COSINE TABLE FOR THE TWIDDLE FACTORS IS EXPECTED TO BE SUPPLIED
*   DURING LINK TIME, AND IT SHOULD HAVE THE FOLLOWING FORMAT:
*
*   .GLOBAL _sine
*   .DATA
*   _sine .FLOAT VALUE1 = sin(0x299/N)
*   _sine .FLOAT VALUE = sin(1x299/N)
*   .....
*   _sine .FLOAT VALUE(5M/4) = sin((5M/4-1)x299/N)
*
*   THE VALUES VALUE1, VALUE2, ETC., ARE THE SAME WAVE VALUES. FOR AN
*   N-POINT FFT, THERE ARE N*M/4 VALUES FOR A FULL AND A QUARTER PERIOD OF
*   THE SINE WAVE. IN THIS WAY, A FULL SINE AND COSINE PERIOD ARE AVAILABLE
*   (SUPERIMPOSED).
*
* STACK STRUCTURE UPON THE CALL:
*   +-----+
*   -FP(4) : DATA
*   -FP(3) : N
*   -FP(2) : M
*   -FP(1) : RETURN ADDR
*   -FP(0) : OLD FP
*   +-----+
*
* REGISTERS USED: R0, R1, R2, R4, R5, R6, R7, ARO, AR1, AR2, AR4, AR5
*                AR6, AR7, IRO, IRI, RS, RE, RC
*
* AUTHOR: PANDOS E. PAPADIMITRAKIS
*        TEXAS INSTRUMENTS
*        OCTOBER 13, 1987
*
*****

```

```

*
* .set AR3
*
* .globl _fft_2
* .globl _sine
*
* .bss FFTSIZ,1
* .bss LOFFT,1
* .bss INPUT,1
*
* .TEXT
*
* .word _sine
*
* SINTAB
*
* INITIALIZE C FUNCTION
*
* _fft_2: PUSH FP
*        LODI SP,FP
*        PUSH RS
*        PUSH R6
*        PUSH R7
*        PUSH AR4
*        PUSH AR5
*        PUSH AR6
*        PUSH AR7
*
*        LODI *-FP(2),R0
*        STI R0,FFTSIZ
*        LODI *-FP(3),R0
*        STI R0,LOFFT
*        LODI *-FP(4),R0
*        STI R0,INPUT
*
*        INITIALIZE FFT ROUTINE
*
*        LODI FFTSIZ,IRI
*        LSH -2,IRI
*        LODI 0,AR6
*        FFTSIZ,IRO
*        LSH 1,IRO
*        LODI FFTSIZ,R7
*        LODI 1,AR7
*        LODI 1,AR5
*
*        OUTER LOOP
*
* LOOP: NOP
*        LODI *-AR6,1
*        AROD R7,ARO,AR2
*        LODI AR7,RC
*        SUBI 1,RC
*
*

```

ENTRY POINT FOR EXECUTION
 ADDRESS OF SINE TABLE

SAVE DEDICATED REGISTERS

MOVE ARGUMENTS TO LOCATIONS WATCHING
 THE MARKS IN THE PROGRAM

IRI=M/4, POINTER FOR SIN/COS TABLE
 AR6 HOLDS THE CURRENT STAGE NUMBER
 IRO=2*M (BECAUSE OF REAL/IMAG)
 R7=M/2
 INITIALIZE REPEAT COUNTER OF FIRST
 LOOP
 INITIALIZE IE INDEX (AR6+IE)

CURRENT FFT STAGE
 ARO POINTS TO X(1)
 AR2 POINTS TO X(L)
 AR7, RC
 RC SHOULD BE ONE LESS THAN DESIRED

Appendix A3. Complex, Radix-2 DIT FFT – R2DIT.ASM

```

*****
*
* COMPLEX, RADIX-2 DIT FFT : R2DIT.ASM
*
* -----
*
* GENERIC PROGRAM FOR A FAST LOOPED-CODE RADIX-2 DIT FFT COMPUTATION
* ON THE TRS320C30
*
* WRITTEN BY: RAIMUND MEYER, KARL SCHWARZ 19.07.89
* LEHRSTUHL FÜR ANWENDUNGSTECHNIK
* UNIVERSITÄT DUISBURG-ESSEN
* CAULOSTRASSE 7, D-6500 ESSEN, FRG
*
* THE (COMPLEX) DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE
* IN-PLACE, BUT THE RESULT IS MOVED TO ANOTHER MEMORY SECTION TO
* DEMONSTRATE THE BIT-REVERSED ADDRESSING.
*
* FOR THIS PROGRAM THE MINIMUM FFTLENGTH IS 32 POINTS BECAUSE OF THE
* SEPARATE STAGES.
*
* FIRST TWO PHASES ARE REALIZED AS A FOUR BUTTERFLY LOOP SINCE THE
* MULTIPLIES ARE TRIVIAL. THE MULTIPLIER IS ONLY USED FOR A LOAD IN
* PARALLEL WITH AN ADD OR SUB.
*
* *****
*
* EXAMPLE FOR A 1024-POINT FFT (EXCLUDING BIT REVERSAL):
*
* MEMORY SIZE:
* PROGRAM = 229 WORDS
* DATA (TWIDDLE FACTORS) = 512 WORDS
*
* CYCLES PER BUTTERFLY:
* STAGES 1 AND 2 = 4
* STAGES 3 TO 8 = 8
* STAGE 9 = 8.25
* STAGE 10 = 8.5
*
* AVERAGE CYCLES/BUTTERFLY = 7.275
* TOTAL BUTTERFLY CYCLES = 37248
* INITIALIZATION OVERHEAD = 2181 = 5.55 % OF TOTAL TIME
* TOTAL NUMBER OF INSTRUCTION CYCLES = 39429
* TOTAL TIME FOR A 1024 POINT FFT = 2.36 ms (EXCLUDING BIT REVERSAL)
*
* *****

```

```

*****
*
* THIS PROGRAM INCLUDES FOLLOWING FILES:
*
* -----
*
* THE FILE 'TWIDDLEFAC.ASM' CONSISTS OF TWIDDLE FACTORS
*
* THE TWIDDLE FACTORS ARE STORED IN BIT-REVERSED ORDER AND WITH A TABLE
* LENGTH OF N/2 (N = FFTLENGTH).
*
* EXAMPLE: SHOW FOR N=32, W(n) = COS(2*PI*n/N) - j*SIN(2*PI*n/N)
*
* ADDRESS COEFFICIENT
* 0 R(W(0)) = COS(2*PI*0/32) = 1
* 1 -I(W(0)) = SIN(2*PI*0/32) = 0
* 2 R(W(4)) = COS(2*PI*4/32) = 0.707
* 3 -I(W(4)) = SIN(2*PI*4/32) = 0.707
*
* 12 R(W(3)) = COS(2*PI*3/32) = 0.831
* 13 -I(W(3)) = SIN(2*PI*3/32) = 0.556
* 14 R(W(7)) = COS(2*PI*7/32) = 0.195
* 15 -I(W(7)) = SIN(2*PI*7/32) = 0.981
*
* WHEN GENERATED FOR A FFT LENGTH OF 1024, THE TABLE IS FOR ALL
* AVAILABLE FFT OF LESS OR EQUAL LENGTH.
*
* THE MISSING TWIDDLE FACTORS (W(1),W(5),....) ARE GENERATED BY USING
* THE SYMMETRY W(N/4+n) = -j*W(n). THIS CAN BE EASILY REALIZED BY
* CHANGING REAL- AND IMAGINARY PART OF THE TWIDDLE FACTORS AND BY
* NEGATING THE NEW REAL PART.
*
* TO CHANGE THE FFT LENGTH, ONLY THE PARAMETERS IN THE HEADER OF
* TWIDDLEFAC.ASM AND THE INPUT AND OUTPUT VECTOR LENGTHS NEED TO BE
* ALTERED.
*
* *****

```

$$AR + j AI \quad \xrightarrow{\quad} \quad AR' + j AI'$$

$$BR + j BI \quad \xrightarrow{\quad} \quad BR' + j BI'$$

```

*
* TR = BR * COS + BI * SIN
* TI = BR * SIN - BI * COS
* AR' = AR + TR
* AI' = AI - TI
* BR' = AR - TR
* BI' = AI + TI
*
* *****

```


[illegible]

[illegible]

[illegible]

[illegible]

[illegible]

Appendix A4. Complex, Radix-2 DIT FFT – R2DITB.ASM

```

*****
* THIS PROGRAM INCLUDES FOLLOWING FILES:
* -----
* THE FILE 'TUDITDOR.ASH' CONSISTS OF TUIDUBLE FACTORS
* THE TUIDUBLE FACTORS ARE STORED IN BIT REVERSED ORDER AND WITH A TABLE
* LENGTH OF N/2 (N = FTLENGTH).
*
* EXAMPLE: SHOW FOR N=32, W(16) = COS(2*PI*16/N) - j*SIN(2*PI*16/N)
*
* ADDRESS COEFFICIENT
* 0 RUM(0) = COS(2*PI*0/32) = 1
* 1 -JUM(0) = SIN(2*PI*0/32) = 0
* 2 RUM(4) = COS(2*PI*4/32) = 0.707
* 3 -JUM(4) = SIN(2*PI*4/32) = 0.707
* .
* .
* 12 RUM(31) = COS(2*PI*31/32) = 0.831
* 13 -JUM(31) = SIN(2*PI*31/32) = 0.356
* 14 RUM(7) = COS(2*PI*7/32) = 0.195
* 15 -JUM(7) = SIN(2*PI*7/32) = 0.981
*
* WHEN GENERATED FOR A FFT LENGTH OF 1024, THE TABLE IS FOR ALL
* AVAILABLE FFT OF LESS OR EQUAL LENGTH.
*
* THE MISSING TUIDUBLE FACTORS (UM(1),UM(2),....) ARE GENERATED BY USING
* THE SYMMETRY UM(N/4+n) = -jUM(n). THIS CAN BE EASILY REALIZED, BY
* CHANGING REAL- AND IMAGINARY PART OF THE TUIDUBLE FACTORS AND BY
* NEGATING THE REAL PART.
*
* TO CHANGE THE FFT LENGTH ONLY THE PARAMETERS IN THE HEADER OF
* TUDITDOR.ASH AND THE INPUT AND OUTPUT VECTOR LENGTHS NEED TO BE
* ALTERED.
*****
*
* AR + j AI ----- AR' + j AI'
*
*
*
*
* BR + j BI ----- ( COS - j SIN ) ----- BR' + j BI'
*
*
*
*
* TR = BR * COS + BI * SIN
* TI = BR * SIN - BI * COS
* AR' = AR + TR
* AI' = AI - TI
* BR' = AR - TR
* BI' = AI + TI
*****

```

```

*****
* APPENDIX A4
*
* COMPLEX, RADIX-2 DIT FFT : R2DITB.ASM
*
*
* GENETIC PROGRAM FOR A FAST LOADED-CODE RADIX-2 DIT FFT COMPUTATION
* ON THE TMS320C30
*
* WRITTEN BY: RAIMUND NEYER, KARL SCHWABZ 24.07.89
* LEHRSTUHL FUER MOCHTIDENTITECHNIK
* UNIVERSITAET ERLANGEN-NUERNBERG
* CAUERSTRASSE 7, D-91020 ERLANGEN, FRG
*
* THE (COMPLEX) DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE
* IN-PLACE, BUT THE RESULT IS MOVED TO ANOTHER MEMORY SECTION TO
* DEMONSTRATE THE BIT-REVERSED ADDRESSING.
*
* FOR THIS PROGRAM THE MINIMUM FFT LENGTH IS 32 POINTS BECAUSE OF
* THE SEPARATE STAGES.
*
* FIRST TWO PHASES ARE REALIZED AS A FOUR BUTTERFLY LOOP SINCE THE
* MULTIPLIES ARE TRIVIAL. THE MULTIPLIER IS ONLY USED FOR A LOAD IN
* PARALLEL WITH AN ADD OR SUB.
*****
*
* EXAMPLE FOR A 1024-POINT FFT (WITH BIT REVERSAL) :
*
* MEMORY SIZE :
* PROG = 231 WORDS
* DATA = 512 WORDS
*
* CYCLES FOR BUTTERFLY :
* STAGES 1 AND 2 = 4
* STAGES 3 TO 8 = 8
* STAGE 9 = 8.25
* STAGE 10 = 10.5 (DUE TO EXT. MEMORY WAIT)
*
* AVERAGE CYCLES/BUTTERFLY = 7.475
* TOTAL BUTTERFLY CYCLES = 38272
* INITIALIZATION OVERHEAD = 2185 = 5.4 % OF TOTAL TIME
* TOTAL NUMBER OF INSTRUCTION CYCLES = 40457
* TOTAL TIME FOR A 1024 POINT FFT = 2.42 ms (INCLUDING BIT
* REVERSAL)
*****

```

Op	Op1	Op2	Op3	Op4	Op5	Op6	Op7	Op8	Op9	Op10	Op11	Op12	Op13	Op14	Op15	Op16	Op17	Op18	Op19	Op20	Op21	Op22	Op23	Op24	Op25	Op26	Op27	Op28	Op29	Op30	Op31	Op32	Op33	Op34	Op35	Op36	Op37	Op38	Op39	Op40	Op41	Op42	Op43	Op44	Op45	Op46	Op47	Op48	Op49	Op50	Op51	Op52	Op53	Op54	Op55	Op56	Op57	Op58	Op59	Op60	Op61	Op62	Op63	Op64	Op65	Op66	Op67	Op68	Op69	Op70	Op71	Op72	Op73	Op74	Op75	Op76	Op77	Op78	Op79	Op80	Op81	Op82	Op83	Op84	Op85	Op86	Op87	Op88	Op89	Op90	Op91	Op92	Op93	Op94	Op95	Op96	Op97	Op98	Op99	Op100	Op101	Op102	Op103	Op104	Op105	Op106	Op107	Op108	Op109	Op110	Op111	Op112	Op113	Op114	Op115	Op116	Op117	Op118	Op119	Op120	Op121	Op122	Op123	Op124	Op125	Op126	Op127	Op128	Op129	Op130	Op131	Op132	Op133	Op134	Op135	Op136	Op137	Op138	Op139	Op140	Op141	Op142	Op143	Op144	Op145	Op146	Op147	Op148	Op149	Op150	Op151	Op152	Op153	Op154	Op155	Op156	Op157	Op158	Op159	Op160	Op161	Op162	Op163	Op164	Op165	Op166	Op167	Op168	Op169	Op170	Op171	Op172	Op173	Op174	Op175	Op176	Op177	Op178	Op179	Op180	Op181	Op182	Op183	Op184	Op185	Op186	Op187	Op188	Op189	Op190	Op191	Op192	Op193	Op194	Op195	Op196	Op197	Op198	Op199	Op200	Op201	Op202	Op203	Op204	Op205	Op206	Op207	Op208	Op209	Op210	Op211	Op212	Op213	Op214	Op215	Op216	Op217	Op218	Op219	Op220	Op221	Op222	Op223	Op224	Op225	Op226	Op227	Op228	Op229	Op230	Op231	Op232	Op233	Op234	Op235	Op236	Op237	Op238	Op239	Op240	Op241	Op242	Op243	Op244	Op245	Op246	Op247	Op248	Op249	Op250	Op251	Op252	Op253	Op254	Op255	Op256	Op257	Op258	Op259	Op260	Op261	Op262	Op263	Op264	Op265	Op266	Op267	Op268	Op269	Op270	Op271	Op272	Op273	Op274	Op275	Op276	Op277	Op278	Op279	Op280	Op281	Op282	Op283	Op284	Op285	Op286	Op287	Op288	Op289	Op290	Op291	Op292	Op293	Op294	Op295	Op296	Op297	Op298	Op299	Op300	Op301	Op302	Op303	Op304	Op305	Op306	Op307	Op308	Op309	Op310	Op311	Op312	Op313	Op314	Op315	Op316	Op317	Op318	Op319	Op320	Op321	Op322	Op323	Op324	Op325	Op326	Op327	Op328	Op329	Op330	Op331	Op332	Op333	Op334	Op335	Op336	Op337	Op338	Op339	Op340	Op341	Op342	Op343	Op344	Op345	Op346	Op347	Op348	Op349	Op350	Op351	Op352	Op353	Op354	Op355	Op356	Op357	Op358	Op359	Op360	Op361	Op362	Op363	Op364	Op365	Op366	Op367	Op368	Op369	Op370	Op371	Op372	Op373	Op374	Op375	Op376	Op377	Op378	Op379	Op380	Op381	Op382	Op383	Op384	Op385	Op386	Op387	Op388	Op389	Op390	Op391	Op392	Op393	Op394	Op395	Op396	Op397	Op398	Op399	Op400	Op401	Op402	Op403	Op404	Op405	Op406	Op407	Op408	Op409	Op410	Op411	Op412	Op413	Op414	Op415	Op416	Op417	Op418	Op419	Op420	Op421	Op422	Op423	Op424	Op425	Op426	Op427	Op428	Op429	Op430	Op431	Op432	Op433	Op434	Op435	Op436	Op437	Op438	Op439	Op440	Op441	Op442	Op443	Op444	Op445	Op446	Op447	Op448	Op449	Op450	Op451	Op452	Op453	Op454	Op455	Op456	Op457	Op458	Op459	Op460	Op461	Op462	Op463	Op464	Op465
----	-----	-----	-----	-----	-----	-----	-----	-----	-----	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------	-------

[illegible]


```

*      B1' = A1 + T1
*
*      RPTB
*      BFLV2
*      MPLY  *+AR1,R7,R5 ; R5 = B1 * COS , (AR' = R5)
*      STP   R1,R0,R2
*      ADDF  *+AR1,R6,R0 ; (R2 = T1 = R0 + R1)
*      MPLY  R2,*AR0,R3 ; R0 = BR * SIN , (R3 = A1 + T1)
*      SUBF  R2,*AR0++ ,R4 ; (R4 = A1 - T1 , B1' = R3)
*      STP   R3,*AR3++
*      SUBF  R0,R5,R3
*      MPLY  *+AR1++ ,R7,R0 ; TR = R3 * R5 - R0
*      SUBF  R3,*AR0,R2 ; R0 = BR * COS , R2 = AR - TR
*      MPLY  *+AR1++ ,R6,R1 ; R1 = B1 * SIN , (A1' = R4)
*      STP   R4,*AR2++
*      ADDF  *+AR0++ ,R3,R5 ; R5 = AR + TR , BR' = R2
*      STP   R2,*AR3++
*
*      CLEAR PIPELINE
*
*      ADDF  R1,R0,R2
*      ADDF  R2,*AR0,R3
*      STP   R5,*AR2++
*      CPTI  AR6,AR4
*      BRND  GROUPE
*      SUBF  R2,*AR0++*(IR1),R4 ; DO FOLLOWING 3 INSTRUCTIONS
*      STP   R3,*AR3++*(IR1) ; R4 = A1 - T1 , B1' = R3
*      LUF   *+AR7,R7 ; R7 = COS
*      STP   R4,*AR2++*(IR1) ; A1' = R4
*      NOP   *+AR1++*(IR1) ; BRANCH HERE
*
*      END OF THIS BUTTERFLY GROUP
*
*      CPTI  4,IR0
*      BRN   STUPE
*
*      SECOND TO LAST STAGE
*
*      LDI   @INPUT,AR0
*      LDI   AR0,AR2
*      ADDI   IR0,AR0,AR1 ; UPPER OUTPUT
*      LDI   AR1,AR3 ; LOWER OUTPUT
*      LDI   @SIMP2,AR7 ; POINTER TO TRIPLE FACTOR
*      LDI   5,IR0
*      LDI   @GARC,RC
*
*      FILL PIPELINE
*
*      1. BUTTERFLY: *0
*      ADDF  *+AR0,*AR1,R2 ; AR' = R2 = AR + BR
*      SUBF  *+AR1++,*AR0++ ,R3 ; BR' = R3 = AR - BR
*      ADDF  *+AR0,*AR1,R0 ; A1' = R0 = A1 + B1

```

```

*      SUBF  *+AR1++,*AR0++ ,R1 ; B1' = R1 = A1 - B1
*
*      2. BUTTERFLY: *0
*
*      ADDF  *+AR0,*AR1,R6 ; AR' = R6 = AR + BR
*      SUBF  *+AR1++,*AR0++ ,R7 ; BR' = R7 = AR - BR
*      ADDF  *+AR0,*AR1,R4 ; A1' = R4 = A1 + B1
*      SUBF  *+AR1++*(IR0),*AR0++*(IR0),R5 ; B1' = R5 = A1 - B1
*      STP   R2,*AR2++ ; (AR' = R2)
*      STP   R3,*AR3++ ; (BR' = R3)
*      STP   R0,*AR2++ ; (A1' = R0)
*      STP   R1,*AR3++ ; (B1' = R1)
*      STP   R6,*AR2++ ; AR' = R6
*      STP   R7,*AR3++ ; BR' = R7
*      STP   R4,*AR2++*(IR0) ; A1' = R4
*      STP   R5,*AR3++*(IR0) ; B1' = R5
*
*      3. BUTTERFLY: *0/4
*
*      ADDF  *+AR0++,*+AR1,R5 ; AR' = R5 = AR + B1
*      SUBF  *+AR1,*+AR0,R4 ; A1' = R4 = A1 - BR
*      ADDF  *+AR1++,*+AR0-- ,R6 ; B1' = R6 = A1 + BR
*      SUBF  *+AR1++,*+AR0++ ,R7 ; BR' = R7 = AR - B1
*
*      4. BUTTERFLY: *0/4
*
*      ADDF  *+AR1,*+AR0,R3 ; AR' = R3 = AR + B1
*      LUF   *+AR7,R1 ; R1 = 0 (FOR INNER LOOP)
*      LUF   *+AR1++ ,R0 ; R0 = BR (FOR INNER LOOP)
*      SUBF  *+AR1++*(IR0),*+AR0++ ,R2 ; BR' = R2 = AR - B1
*      STP   R5,*AR2++ ; (AR' = R5)
*      STP   R7,*AR3++ ; (BR' = R7)
*      STP   R6,*AR3++ ; (B1' = R6)
*
*      5. TO M. BUTTERFLY:
*
*      RPTB  BF2ND0
*
*      LUF   *+AR7++ ,R7 ; R7 = COS , (A1' = R4)
*      STP   R4,*AR2++
*      LUF   *+AR7++ ,R6 ; R6 = SIN , (BR' = R2)
*      STP   R2,*AR3++
*      MPLY  *+AR1,R6,R3 ; R5 = B1 * SIN , (AR' = R3)
*      STP   R2,*AR2++
*      ADDF  R1,R0,R2 ; (R2 = T1 = R0 + R1)
*      MPLY  *+AR1,R7,R0 ; R0 = BR * COS , (R3 = A1 + T1)
*      SUBF  R2,*AR0,R3 ; (R4 = A1 - T1 , B1' = R3)
*      STP   R3,*AR3++*(IR0),R4
*      ADDF  *+AR1++ ,R6,R0 ; R3 = TR = R0 + R5
*      MPLY  *+AR1++ ,R6,R0 ; R0 = BR * SIN , R2 = AR - TR
*      SUBF  *+AR1++ ,R7,R1 ; R1 = B1 * COS , (A1' = R4)
*      MPLY  R4,*AR2++*(IR0)
*      STP   R4,*AR2++*(IR0)

```



```

;;      R3, #480, R2
MOVFE  #481, *(IR1), R7, R1
STF    R4, #483, *(IR0)B
SUBFE  R3, #480, R3
;;      BR' = R3 = AR - TR, AR' = R2
STF    R2, #482, *(IR0)B
*
MOVFE  #481, R7, R5
STF    R3, #482, *(IR0)B
SUBFE  R1, R0, R2
MOVFE  #481, R6, R0
SUBFE  R2, #480, R3
;;      (R2 = T1 = R0 - R1)
STF    R3, #483, *(IR0)B
MOVFE  R0, R5, R3
SUBFE  R3, #480, R2
MOVFE  #481, R7, R0
;;      R3 = TR = R0 - R5
ANDFE  R3, #480, R2
MOVFE  #481, *(IR1), R6, R1
SUBFE  R3, #480, R3
;;      R0 = BR * SIN, (A1' = R3 = A1 - T1)
*      R5 = B1 * COS, (BR' = R3)
*      (R2 = T1 = R0 - R1)
*      R0 = BR * SIN, (A1' = R3 = A1 - T1)
*      (B1' = R4 = A1 + T1, A1' = R3)
*      R3 = TR = R0 - R5
*      R0 = BR * COS, AR' = R2 = AR + TR
*      R1 = B1 * SIN, BR' = R3 = AR - TR
BELEND
*
*      CLEAR PIPELINE
*
;;      R2, #482, *(IR0)B
STF    R4, #483, *(IR0)B
MOVFE  R1, R0, R2
SUBFE  R2, #480, R3
;;      R2 = T1 = R0 + R1
STF    R3, #482
MOVFE  R2, #480, R4
SUBFE  R3, #483, *(IR0)B
;;      A1' = R3 = A1 - T1, BR' = R3
*      B1' = R4 = A1 + T1, A1' = R3
*      B1' = R4
*      END OF FFT
*
END:
NOP
NOP
NOP
*
SELF  BR    SELF
.end
*

```

```

7.11432195745216e-001
7.02754744457225e-001
6.13588464915432e-003
9.99981175282601e-001

```

```
*****  
# APPENDIX A5  
#  
# TITLE: TWO-DIM.ASM  
#  
# TABLE WITH TWO-DIM FACTORS FOR A FFT UP TO A LENGTH OF 1024 COMPLEX  
# POINTS.  
#  
# FILE TO BE LINKED WITH THE SOURCE CODE : R2DIT.ASM OR R2DITL.ASM  
#  
# WRITTEN BY : RAUHMUND MEYER AND KARL SCHWARZ  
#              LEHRSTUHL FÜR MASCHINENGEOMETRIK  
#              UNIVERSITÄT DUISBURG-ESSEN  
#  
# LENGTH OF TWO-DIM FACTOR TABLE : 512 REAL VALUES (=1024 FFT)  
#  
*****  
#  
# global sine  
# global n  
# global ohalb  
# global invert  
# global macthel  
# global macthel  
#  
# .set 1024 ; FFT-LENGTH n  
# .set 512 ; n/2  
# .set 256 ; n/4  
# .set 128 ; n/8  
# .set 10 ; NUMBER OF STAGES = 14(n)  
#  
# ANOTHER EXAMPLE OF FFT-LENGTH n = 32:  
# ONLY THE FIRST 16 VALUES OF THE TABLE ARE NEEDED  
#  
# .set 2  
# .set 16  
# .set 8  
# .set 4  
# .set 5  
#  
# .delta  
#  
# sine  
# float 1.000000000000000000000000000000  
# float 0.000000000000000000000000000000  
# float 7.07106781186548e-001  
# float 7.07106781186548e-001  
# float 9.2387932311287e-001  
# float 3.8268343236599e-001  
# float 3.8268343236599e-001  
# float 9.2387932311287e-001  
# float 9.8078328640223e-001
```

Appendix B. Radix-4 Complex FFT

[illegible]

[illegible]

LSH 2,AR7 : INCREMENT REPEAT COUNTER FOR NEXT

Appendix B2. fft_4--Radix-4 Complex FFT to Be Called as a C Function

```

*
* APPENDIX B2
*
* NAME: fft_4 --- RADIX-4 COMPLEX FFT TO BE CALLED AS A C FUNCTION.
*
* SYNOPSIS:
*
*   int  fft_4(int N, DATA)
*   int  N      FFT SIZE: N=4*4*4*4
*   float data[] ARRAY WITH INPUT AND OUTPUT DATA
*
* DESCRIPTION:
*
*   GENERIC FUNCTION TO DO A RADIX-4 FFT COMPUTATION ON THE TMS320C30.
*   THE DATA ARRAY IS 24M-LONG, WITH REAL AND IMAGINARY VALUES ALTERNATING.
*   THE PROGRAM IS BASED ON THE FORTRAN PROGRAM IN THE BURRUS AND PARRIS BOOK, P. 117.
*
*   IN ORDER TO HAVE THE FINAL RESULT IN BIT-REVERSED ORDER, THE TWO
*   RIDGE BRANCHES OF THE RADIX-4 BUTTERFLY ARE INTERCHANGED DURING
*   STORAGE. NOTE THIS DIFFERENCE WHEN COMPARING WITH THE PROGRAM ON
*   P. 117. THE COMPUTATION IS DONE IN-PLACE, AND THE ORIGINAL DATA IS
*   DESTROYED. BIT REVERSAL IS IMPLEMENTED AT THE END OF THE FUNCTION.
*   IF THIS IS NOT NECESSARY, THIS PART CAN BE COERCED OUT. THE
*   SINE/COSINE TABLE FOR THE TWIDDLE FACTORS IS EXPECTED TO BE SUPPLIED
*   DURING LINK TIME, AND IT SHOULD HAVE THE FOLLOWING FORMAT:
*
*   .global  _sine
*   .data
*   _sine:   value1 = sin(0*2pi/N)
*   .float   value2 = sin(1*2pi/N)
*   .....
*   .float   value(SN/4) = sin((SN/4-1)*2pi/N)
*
*   THE VALUES value1, value2, ETC., ARE THE SINE WAVE VALUES. FOR AN
*   N-POINT FFT, THERE ARE SN/4 VALUES FOR A FULL AND A QUARTER PERIOD
*   OF THE SINE WAVE. IN THIS WAY, A FULL SINE AND COSINE PERIOD ARE
*   AVAILABLE (SUPERIMPOSED).
*
*   STACK STRUCTURE UPON THE CALL:
*
*   +-----+
*   -FP(4)  : DATA
*   -FP(3)  : N
*   -FP(2)  : N
*   -FP(1)  : RETURN ADDR
*   -FP(0)  : OLD FP
*   +-----+
*
*   REGISTERS USED: R0, R1, R2, R3, R4, R5, R6, R7, R8, R9, A01, A02, A03, A04,
*   A05, A06, A07, I00, I01, I02, I03, I04
*
*   AUTHOR: PAMOS E. PAPADIMITRAKIS
*   TEXAS INSTRUMENTS
*
*   OCTOBER 13, 1987
*
* *****

```

```

*
* FP      .SET      A03
*
* .GLOBAL _fft_4
* .GLOBAL _sine
*
* .BSS    FFTSIZ,1
* .BSS    LOFFFT,1
* .BSS    INPUT,1
*
* .TEXT
*
* .SYMTAB .word      _sine
*
* * INITIALIZE C FUNCTION
*
* _fft_4:  PUSH      FP
*          LOD      SP,FP
*          PUSH     R4
*          PUSH     R5
*          PUSH     R6
*          PUSH     R7
*          PUSH     A04
*          PUSH     A05
*          PUSH     A06
*          PUSH     A07
*
*          LOD      -FP(2),R0
*          STI      R0,0FFTSIZ
*          LOD      -FP(3),R0
*          STI      R0,LOFFFT
*          LOD      -FP(4),R0
*          STI      R0,INPUT
*
*          * MOVE ARGUMENTS TO LOCATIONS MATCHING
*          * THE NAMES IN THE PROGRAM
*
*          * FFT STAGE #
*          * REPEAT COUNTER
*          * IE INDEX FOR SINE/COSINE
*          * SECOND-LOOP COUNT
*          * JT COUNTER IN PROGRAM, P. 117
*          * I01 INDEX IN PROGRAM, P. 117
*
*          .BSS     STAGE,1
*          .BSS     RPTCNT,1
*          .BSS     IEINDX,1
*          .BSS     LPOUNT,1
*          .BSS     JT,1
*          .BSS     I01,1
*
*          LOD      0FFTSIZ,R0
*          LOD      0FFTSIZ,I00
*          LOD      0FFTSIZ,I01
*          LOD      0,A07
*          STI      A07,ISTAGE
*
*          * ISTAGE HOLDS THE CURRENT STAGE
*          * NUMBER
*          * I00=24M (BECAUSE OF REAL/IMAG)
*          * I01=4M/A, POINTER FOR SINE/COS TABLE
*          * INITIALIZE REPEAT COUNTER OF FIRST
*          * LOOP
*
*          LSH      1,I00
*          LSH      -2,I01
*          LOD      1,A07
*          STI      A07,0RPTCNT

```

[illegible]

• LOOP BACK TO THE INNER LOOP

POP A07
POP A06
POP A05
POP A04
POPF R7
POPF R6
POP R5
POP R4
POP FP
RETS

Appendix C.Radix-2 Real FFT

[illegible]

```

*
LDI  AR5,AR1
ADDI 1,AR1
LDI  AR1,AR3
LDI  R2,AR3
LDI  AR3,AR2
SUBI 2,AR2
ADDI R3,AR2,AR4
*
LDF  *AR5++(IR1),R0
AUIF *AR5(IR1),R0,R1
SUBF R0,*AR5(IR1),R0
STF  R1,*AR5(IR1)
NEGF R0
NEGF *AR5(IR1),R1
STF  R0,AR5
STF  R1,AR5
*
* INCREMENT LOOP
*
LDI  OFFSET,IR1
LSH  -2,IR1
LDI  R4,RC
SUBI 2,RC
*
BLK3
RPTB
RPTF *AR3,*AR0(IR1),R0
RPTF *AR4,*AR0,R1
RPTF *AR4,*AR0(IR1),R1
AUIF R0,R1,R2
RPTF *AR3,*AR0++(IR0),R0
SUBF *AR2,R0,R1
SUBF *AR2,R0,R1
AUIF R1,AR3++
STF  R1,AR3++
AUIF R1,R2,R1
STF  R1,AR4--
SUBF R2,AR0,R1
STF  R1,AR0++
BLK3
*
SUBI INPUT,AR5
ADDI R4,AR5
CPI  OFFSET,AR5
RLTO INLOP
ADDI INPUT,AR5
NOP
*
AUI 1,R5
CPI  BLOFFT,R5
BLE LOOP
NOP
NOP

```

```

NOP
NOP
*
END BR
END
;BRANCH TO ITSELF AT THE END
;END

```

Appendix C2. fft_rl – Radix-2 Real FFT to Be Called as a C Function

```

*
* APPENDIX C2
*
* NAME:
*   fft_rl --- RADIX-2 REAL FFT TO BE CALLED AS A C FUNCTION.
*
* SYNOPSIS:
*   int fft_rl(int n, data);
*   int n      FFT SIZE: M/2+M
*   int m      NUMBER OF STAGES = LOG2(N)
*   float *data  ARRAY WITH INPUT AND OUTPUT DATA
*
* DESCRIPTION:
*   GENERIC FUNCTION TO DO A RADIX-2 FFT COMPUTATION ON THE TMS320C30.
*   THE DATA ARRAY IS M-LONG, WITH ONLY REAL DATA. THE OUTPUT IS STORED
*   IN THE SAME LOCATIONS WITH REAL AND IMAGINARY POINTS R AND I AS
*   FOLLOWS: R(0), R(1), ..., R(M/2), I(M/2+1), ..., I(1)
*
*   THE PROGRAM IS BASED ON THE FORTRAN PROGRAM IN THE PAPER BY SOROKEN
*   ET AL., JUNE 1987 ISSUE OF TRANS. ON ASSP. THE COMPUTATION IS DONE
*   IN-PLACE, AND THE ORIGINAL DATA IS DESTROYED. BIT REVERSAL IS
*   IMPLEMENTED AT THE BEGINNING OF THE FUNCTION. IF THIS IS NOT
*   NECESSARY, THIS PART CAN BE COMMENTED OUT.
*
*   THE SINE/COSINE TABLE FOR THE INITIAL FACTORS IS EXPECTED TO BE
*   SUPPLIED DURING LINK TIME, AND IT SHOULD HAVE THE FOLLOWING FORMAT:
*
*   .global  _sine
*   .data
*   _sine:
*   .float value1 = sin((0+2*pi)/N)
*   .float value2 = sin((1+2*pi)/N)
*   .....
*   .float value(N/2) = cos((N/4)*2*pi/N)
*
*   THE VALUES value1 TO value(N/4) ARE THE FIRST QUARTER OF THE SINE
*   PERIOD AND value(N/4+1) TO value(N/2) ARE THE FIRST QUARTER OF THE
*   COSINE PERIOD.
*
*   STACK STRUCTURE UPON THE CALL:
*   +-----+
*   -FP(4)  : DATA
*   -FP(3)  : n
*   -FP(2)  : N
*   -FP(1)  : RETURN ADDR
*   -FP(0)  : OLD FP
*   +-----+
*
*   REGISTERS USED: R0, R1, R2, R3, R4, R5, AR0, AR1, AR2, AR4, AR5, IRO,
*   IRI, RS, RE, RC
*
*   AUTHOR: PHILIP E. PAPATHOMASIS
*   TEXAS INSTRUMENTS
*
*   OCTOBER 13, 1987
*
* *****

```

```

*
* .SET      AR0
*
* .GLOBAL   _fft_rl
* .GLOBAL   _sine
*
* .BSS      _fft_rl
* .BSS      _sine
* .BSS      _fft_rl
* .BSS      _sine
* .BSS      _fft_rl
* .BSS      _sine
*
* .TEXT
*
* .SIMTAB   .word  _sine
*
* * INITIALIZE C FUNCTION
*
* _fft_rl:
*   PUSH    FP
*   LODI    SP, FP
*   PUSH    R4
*   PUSH    R5
*   PUSH    AR4
*   PUSH    AR5
*
*   LODI    +FP(2), R0
*   STI     R0, _fft_rl
*   LODI    +FP(3), R0
*   STI     R0, _sine
*   LODI    +FP(4), R0
*   STI     R0, _fft_rl
*
* * DO THE BIT REVERSING AT THE BEGINNING
*
*   LODI    _fft_rl, RC
*   SUBI    1, RC
*   LODI    _fft_rl, IRO
*   LSH     -1, IRO
*   LODI    @INPUT, AR0
*   LODI    @INPUT, AR1
*
*   RPTB    BITRV
*   CPTI    AR1, AR0
*   BGE     CONT
*   LDF     #AR0, R0
*   LDF     #AR1, R1
*   STF     R0, AR0
*   STF     R1, AR1
*   NOP     #AR0++
*   NOP     #AR1++(IRO)B
*   CONT
*   BITRV
*
* * LENGTH-TWO BUTTERFLIES
*
*   LODI    @INPUT, AR0
*   LODI    IRO, RC
*   SUBI    1, RC

```


[illegible]

Appendix C3. Generic Program to Do a Radix-2 Real Inverse FFT Computation on the TMS320C30

```

*
* APPENDIX C3
*
* GENERIC PROGRAM TO DO A RADIX-2 REAL INVERSE FFT COMPUTATION ON THE
* TMS320C30.
*
* THE (REAL) DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE
* IN-PLACE. THE BIT REVERSAL IS DONE AT THE BEGINNING OF THE PROGRAM. THE
* INPUT DATA ARE STORED IN THE FOLLOWING ORDER:
*
* RE(0), RE(1),..., RE(N/2-1), ..., IN(1)
*
* THE TWIDDLE FACTORS ARE SUPPLIED IN A TABLE PUT IN A DATA SECTION. THIS
* DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC NATURE OF THE
* PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FFT N AND LOG2(N) ARE
* DEFINED IN A .GLOBL DIRECTIVE AND SPECIFIED DURING LINKING. THE LENGTH OF
* THE TABLE IS N/4 * N/4 = N/2.
*
* AUTHOR: RAMOS PARACHALIS
* TELAS INSTRUMENTS
*
* .GLOBL IFFT
* .GLOBL N
* .GLOBL M
* .GLOBL SINE
*
* .BSS TMP,1024
*
* .TEXT
*
* INITIALIZE
*
* .WORD IFFT
*
* .SPACE 100
*
* FFTSIZ .WORD N
* LOGFFT .WORD M
* SINTAB .WORD SINE
* INPUT .WORD TMP
*
* IFFT: LUP FFTSIZ
*
* MAIN LOOP (FFT STAGES)
*
* LDI 1,R0
* LDI 3,R5
* LDI 0,FFTSIZ,R0
* LSH -1,R0
* LDI 0,FFTSIZ,R4
* LSH -2,R4
*
* INNER LOOP
*
* ; ENTRY POINT FOR EXECUTION
* ; FFT SIZE
* ; LOG2(N)
* ; ADDRESS OF SINE TABLE
* ; MEMORY WITH INPUT DATA
*
* ; STARTING LOCATION OF THE PROGRAM
* ; RESERVE 100 WORDS FOR VECTORS, ETC.
*
* ; COMMAND TO LOAD DATA PAGE POINTER
*
* ; IRO=INDEX FOR E
* ; R5 HOLDS THE CURRENT STAGE NUMBER
* ; R0=N/1/2+2
* ; R4=N/1/4+4
*
* ; IR1=SEPARATION BETWEEN SIN/COS TABLS
* ; REPEAT NM-1 TIMES
*
* ; IR1=SEPARATION BETWEEN SIN/COS TABLS
* ; REPEAT NM-1 TIMES
*
* ; R1=T*(11)-(112)
* ; R0=T*4COS
* ; X(11)=X(11)+X(12)
* ; R2=72+X(13)+X(14)
* ; R4=72+SIN
* ; X(12)=X(14)-X(13)
* ; R2=72+SIN
* ; R4=72+COS
* ; X(13)=T*4COS-72+SIN
* ; R0=T*5SIN
* ; X(14)=T*5SIN+72+COS
*
* ; LOOP BACK TO THE INNER LOOP
*
* ; ARO POINTS TO SIN/COS TABLE

```

[illegible]

Appendix D. Discrete Hartley Transform

```

* APPENDILI D1
*
* GENETIC PROGRAMM TO DO A PAULI-2 HARTLEY TRANSFORM ON THE TRESOCCO.
*
* THE PROGRAM IS TAKEN FROM THE PAPER BY SØRENSEN ET AL., OCT 1968 ISSUE
* OF THE TRANSACTIONS ON ASSP".
*
* THE IDEAL DATA RESIDE IN INTERNAL MEMORY, THE COMPIATION IS DONE
* IN-PACE. THE BIT-REVERSAL IS DONE AT THE BEGINNING OF THE PROGRAM.
*
* THE TRIVIAL FACTORS ARE SUPPLIED IN A TABLE PUT IN A .DATA SECTION. THIS
* DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC NATURE OF THE
* PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FIT N AND LOGZ(N) ARE
* DEFINED IN A .GLOBAL DIRECTIVE AND SPECIFIED DURING LINKING. THE LENGTH OF
* THE TABLE IS M/A * N/A = N/2.
*
* AUTHOR: PHILIP PAPATHOMASIS
*          TEXAS INSTRUMENTS
*
*          DECEMBER 14, 1980
*
* .GLOBAL FIT           ; ENTRY POINT FOR EXECUTION
* .GLOBAL N             ; FIT SIZE
* .GLOBAL M             ; LOGZ(N)
* .GLOBAL SINE          ; ADDRESS OF SINE TABLE
*
* .RSS IMP, LOGZ        ; MEMORY WITH INPUT DATA
*
* .TEXT
*
* INITIALIZE
*
* .WORD FIT             ; STARTING LOCATION OF THE PROGRAM
*
* .SPACE 100            ; RESERVE 100 WORDS FOR VECTORS, ETC.
*
* FRTSIZ               ;
* .WORD N               ;
* .WORD N               ;
* .WORD N               ;
* .WORD SINE            ;
* .WORD IMP             ;
*
* FRTSIZ               ;
* LBP                  ; COMMAND TO LOAD DATA PAGE POINTER
*
* DO THE BIT REVERSING AT THE BEGINNING
*
*          LDI           @FRTSIZ,PC
*          SUB# 1,PC      ; PC-M
*          LDI           @FRTSIZ,IRO
*          LSH -1,IRO     ; PC SHOULD BE ONE LESS THAN DESIRED #
*          LDI           @IMP,IRO
*          LDI           @IMP,ARI
*
*          RPTB          BITRV
*          CPT#          ARI,ARO
*          BBE          CNT
*          JDE          #ARO,RD
*          IF ARX(ARI)   ; XCHANGE LOCATIONS ONLY
*                   ; IF ARX(ARI)

```


Appendix E. Discrete Cosine Transform

Appendix E1. A Fast Cosine Transform

[illegible]


```

* INCLUDES LAST BUTTERFLIES AND FIRST STAGE OF BIT REVERSE ADDITIONS.
*
LUI 4,IR1 ; INITIALIZE INDEX REGISTER.
ANDI 1,AR3 ; SET UP DATA POINTERS.
LSH -1,AR5
ANDI 3,AR4
ANDI 2,IR0,RC ; INITIALIZE REPEAT COUNTER.
ANDI 4,AR1 ; CALCULATE (2/IR1)*MODS(PI/4).
PPF3 4AR7,4AR7,IR4 ; (I.E.-> (SORT(2)/IR1) THIS VALUE IS CALLED, S, BELOW.)
RPTB END_2ND_LOOP ; TWO BUTTERFLIES ARE CALCULATED PER LOOP.
*
SUBF3 4AR2,4AR1,RO ; SUBTRACT 1ST BUTTERFLY DATA.
SUBF3 4AR4,4AR3,RI ; SUBTRACT 2ND BUTTERFLY DATA.
PPF3 RO,IR4,RO ; MULTIPLY 1ST SUBTRACTION RESULT BY S.
AUF3 4AR3+((IR1),4AR4+((IR1),R3 ; BY S. AND 2ND BUTTERFLY DATA.
PPF3 RI,IR4,RI ; MULTIPLY 2ND SUBTRACTION RESULT BY S.
AUF3 4AR1+((IR1),4AR2+((IR1),R2 ; BY S. AND 1ST BUTTERFLY DATA.
PPF3 R3,4AR7,R3 ; MULTIPLY 2ND ADDITION RESULT BY 7071.
STF RO,4AR2(IR1) ; 7071. SAVE 1ST SUBTRACTION IN LOWER 1/2 OF 1ST BUTTERFLY.
PPF3 R2,4AR7,R2 ; MULTIPLY 1ST ADDITION RESULT BY 7071.
STF RI,4AR4(IR1) ; 7071. SAVE 2ND SUBTRACTION IN LOWER 1/2 OF 2ND BUTTERFLY.
AUF3 R2,RI,R3 ; ADD 2ND SUBTRACTION MULTIPLY TO 2ND ADDITION MULTIPLY.
STF R2,4AR1(IR1) ; SAVE 1ST ADDITION MULTIPLY IN UPPER 1/2 OF BUTTERFLY.
*
END_2ND_LOOP:
*
STF R3,4AR3(IR1) ; SAVE 2ND ADDITION MULTIPLY IN UPPER 1/2 OF UPPER BUTTERFLY.
*
END OF FINAL BUTTERFLY STAGE LOOP.
*
BIT REVERSE ADDITION LOOP SERIES.
*
THIS LOOP SERIES DOES ALL OF THE BIT REVERSE ADDITIONS AT THE END OF FIRST COSTUME TRANSFORM.
*
LUI 2,IR0 ; INITIALIZE INDEX REGISTERS AND DATA.
LUI 4AR6,AR1 ; POINTERS FOR FINAL ADDITION SERIES.
ANDI 4,AR1
LUI 4AR1,AR2
LUI 8,IR1
*
LAST_OUTSIDE_LOOP:
*
LUI 4AR2,AR4 ; UPDATE POINTERS AND COUNTERS.
LSH -1,AR5

```

```

LUI 4AR5,RC ; SET UP REPEAT COUNTER.
AUF3 4AR2+((IR0),B,4AR4+((IR0),B,RO ; DATA POINTER UPDATE.
LUI 4AR1,IR4 ; USE INITIAL 4AR1 VALUE AS INNER LOOP CONTROL.
*
SUBI 1,RC ; CONTINUE UPDATING POINTERS.
NOP 4AR4+((IR0),B
LUI 4AR2,AR3 ; CONTINUE UPDATING POINTERS.
*
RPTB END_INSIDE ; TWO ADDITIONS ARE DONE IN EACH LOOP.
*
LAST_INSIDE_LOOP:
*
AUF3 4AR1,4AR2+((IR1),2,RO ; ADD FIRST TWO DATA.
AUF3 4AR3,4AR4+((IR1),2,RI ; ADD SECOND TWO DATA.
STF RO,4AR1+((IR1),2 ; SAVE FIRST ADDITION.
*
END_INSIDE:
*
STF RI,4AR3+((IR1),2 ; SAVE SECOND ADDITION.
*
END OF INSIDE LOOP FOR LAST LOOP SERIES.
*
AUF3 4AR1+((IR0),B,4AR2+((IR0),B,RO ; UPDATE DATA POINTERS.
AUF3 4AR3+((IR0),B,4AR4+((IR0),B,RO
AUF3 4AR5+((IR0),B,4AR4+((IR0),B,RO
CP1 R4,AR4 ; IS THIS LOOP COMPLETE?
BNE LAST_INSIDE_LOOP ; DELAYED BRANCH, IF NOT.
LUI 4AR5,RC ; SET UP REPEAT COUNTER.
SUBI 1,RC
OR 0100H,ST ; SET REPEAT MODE.
*
BRANCH DELAYED TO LAST_INSIDE_LOOP.
*
RPTB LAST_BLOCK ; SINCE THERE ARE AN ODD NUMBER OF ADDITIONS, THE FINAL ONES ARE DONE NOW.
*
LAST_BLOCK:
*
STF RO,4AR1+((IR1),2 ; SAVE ADDITION.
*
END OF LAST REPEAT BLOCK.
*
LSH 1,IR0 ; MULTIPLY IR0 BY 2.
ANDI 4AR1,IR4 ; UPDATE INNER LOOP CONTROL REGISTER.
CP1 1,AR5 ; ARE CALCULATIONS COMPLETE?
BGT LAST_OUTSIDE_LOOP ; DELAYED BRANCH, IF NOT.
LUI 4AR2,AR2 ; UPDATE DATA POINTERS.
LUI 4AR1,IR4
LSH 1,IR1 ; MULTIPLY IR1 BY 2.
*
BRANCH DELAYED TO LAST_OUTSIDE_LOOP.

```

```

*   END OF LAST LOOP SERIES.
*
*   MULTIPLY COEFFICIENT ZERO BY .5. IF NOT ZERO.
*
*       LDF    *A06,R0      ; SET ZERO FLAG IF *A06 = 0.
*       BEQ0   DONT_STORE   ; IF COEFFICIENT IS ZERO, DON'T DO
*                               ; THIS.
*       LSH    24,A06        ; USE INTEGER PART FOR FLOAT DIVIDE
*                               ; BY 2.
*       SUB13   A05,*A06,A01
*       NOP
*
*   DELAYED BRANCH FROM HERE IF VALUE IS NOT TO BE STORED.
*
*       ST1    A01,*A06      ; STORE, IF EXPONENT WASN'T -128.
*   DONT_STORE:
*       RETS

```

```

APPENDIX E2
A FAST COSINE TRANSFORM (INVERSE TRANSFORM)

BASED ON THE ALGORITHM OUTLINED BY BREYER OF LEE IN HIS ARTICLE, FET - A
FAST COSINE TRANSFORM, PUBLISHED IN THE PROCEEDINGS OF THE IEEE INTER-
NATIONAL CONFERENCE ON ACoustICS, SPEECH, AND SIGNAL PROCESSING, SAN
DIEGO, CA, 19-21 MARCH 1984, P. 264-37/-4 VOL. 2., (CH1954-5/84-0000-0299),
LEE'S ALGORITHM HAS BEEN MODIFIED TO ALLOW NATURAL ORDER TIME DOMAIN
COEFFICIENTS.

THE FREQUENCY DOMAIN COEFFICIENTS ARE IN BIT REVERSE ORDER. THIS IS AN IN
PLACE CALCULATION.

AUTHOR: PAUL WILHELM

.global TEXT
; INVERSE FAST COSINE TRANSFORM ENTRY
TEXT
.global H
; POINT.
.global COEFF
; LENGTH OF ARRAY TO BE TRANSFORMED.
.global COS_TAB
; TABLE OF COSINE COEFFICIENTS.
; TABLE OF ARRAY DATA TO BE
; TRANSFORMED.

.text
.word H
; COEFF
.word COS_TAB
; COS_TAB

FETS:
LDI #COSIZE,ARO
LDI #COSIZE,BK
; LOAD ARRAY SIZE.
; LOAD BLOCK SIZE FOR CIRCULAR ADDRESSING
LDI @DATA,A#6
LDI @COS,AR7
; LOAD POINTER TO DATA TABLE.
ADUI ARO,AR7
SUBI 2,AR7
; POINT TO LAST COSINE VALUE IN TABLE.
LDI ARO,IRO
LSH -2,IRO
LDI ARO,IRI
LDI ARA,ARI
ADUI IRO,ARI

START OF BIT REVERSED ADDITION LOOP SERIES.
OUTSIDE:
ADUI IRO,ARI
; TOP OF OUTSIDE LOOP FOR BIT REVERSED ADCTIONS.
LDI ARI,AR2
LDI IRO,RC
SUBI 2,RC
; UPDATE DATA POINTERS AND REPEAT COUNTING.

```

[illegible]

[illegible]

Appendix E3. FCT Cosine Tables File

```
*
* APPENDIX E3
*
* FCT COSINE TABLES FILE
*
* TO BE LINKED WITH FCT SOURCE CODE FOR 32 POINT FCT.
*
* COEFFICIENTS ARE  $1/(2 * \cos(N\pi/2M))$ , WHERE N IS A NUMBER FROM 1 to
* M-1. M IS THE ORDER OF THE TRANSFORM.
*
* FOR A 32 POINT FCT, N IS IN THE FOLLOWING ORDER:
*     1, 15, 3, 13, 5, 11, 7, 9,
*     2, 14, 6, 10,
*     4, 12,
*     8
*
* THE LAST VALUE IN THE TABLE IS 2/M.
*
*
*     .global  COS_TAB
*     .global  M
*
* M     .set    16
*
*     .data
*
* COS_TAB
*     .float   0.5024193
*     .float   5.1011487
*     .float   0.5224986
*     .float   1.7224471
*     .float   0.5669440
*     .float   1.0606777
*     .float   0.6468218
*     .float   0.7881546
*     .float   0.5097956
*     .float   2.5629154
*     .float   0.6013449
*     .float   0.8999762
*     .float   0.5411961
*     .float   1.3065630
*     .float   0.7071068
*     .float   0.1250000
*     .end
```

Appendix E4. Data File

```
*
* APPENDIX E4
*
* DATA FILE
*
      .global COEFF
*
      .data
*
COEFF
      .float 137.0
      .float 249.0
      .float 105.0
      .float 217.0
      .float 73.0
      .float 185.0
      .float 41.0
      .float 153.0
      .float 9.0
      .float 121.0
      .float 233.0
      .float 89.0
      .float 201.0
      .float 57.0
      .float 169.0
      .float 25.0
      .end
```

Appendix F. Test Vectors, 64-Point Sine Table, Link Command File

Appendix F1. Example of a 64-Point Vector to Test the FFT Routines

```
*
* APPENDIX F1
*
* EXAMPLE OF A 64-POINT VECTOR TO TEST THE FFT ROUTINES
*
X =
0.2113
0.0624
0.7599
0.0087
0.8096
0.8474
0.4534
0.8075
0.4832
0.6135
0.2749
0.8807
0.6538
0.4899
0.7741
0.9426
0.9933
0.8360
0.7469
0.0378
0.4237
0.2613
0.2403
0.3405
0.1167
0.6250
0.3510
0.9550
0.4943
0.0365
0.2260
0.8159
0.2284
0.8553
0.6421
0.7075
0.2408
0.6907
0.1062
0.2640
3.7034
0.4021
0.6553
0.9700
0.0380
0.0988
0.2560

*
* 64-POINT FFT CORRESPONDING TO VECTOR X
*
Y =
30.3774
1.7780 - 2.5584i
-1.0376 - 2.3999i
-1.0123 + 2.4889i
0.6594 + 2.3459i
-1.5228 - 0.7527i
-3.8171 - 0.2050i
-2.7096 + 1.2841i
2.1622 - 1.6831i
0.2879 + 1.8671i
-1.5479 + 1.6298i
-0.6366 - 0.1176i
2.2902 + 1.5549i
-2.4837 - 0.5942i
-1.7338 + 0.0738i
-0.2180 - 0.4726i
-0.2104 + 0.4897i
-1.7473 - 1.0213i
0.1230 - 2.3915i
-0.6415 - 1.1144i
-2.7719 - 0.4802i
-0.0063 - 0.3885i
-0.7163 + 1.5682i
0.3218 - 1.3316i
-0.7823 + 1.0607i
-0.2553 + 2.8270i
-1.0813 - 2.7861i
3.4649 + 1.9485i
3.0352 + 1.3853i
3.2099 + 2.5844i
-1.9511 - 0.7714i
1.8755 + 0.2867i
```

-1.5474
1.8795 - 0.2867i
-1.9511 + 0.7714i
3.2099 - 2.3544i
3.0352 - 1.3855i
3.4869 - 1.9485i
-1.0813 + 2.7811i
-0.2553 - 2.8276i
-0.7823 - 1.0607i
0.3218 + 1.3316i
-0.7163 - 1.5882i
-0.0063 + 0.3885i
-2.7719 + 0.4802i
-0.6415 + 1.1144i
0.1223 + 2.3915i
-1.7873 + 1.0213i
-0.2104 - 0.4897i
-0.2180 + 0.4728i
-1.7338 - 0.0738i
-2.4837 + 0.5942i
2.2902 - 1.5549i
-0.6366 + 0.1176i
-1.5479 - 1.6298i
0.2879 - 1.8671i
2.1622 + 1.6853i
-2.7096 - 1.2841i
-3.8171 + 0.2050i
-1.5228 + 0.7527i
0.6594 - 2.3639i
-1.0123 - 2.4889i
-1.0376 + 2.3999i
1.7780 + 2.5584i

Appendix F2. File to Be Linked with the Source Code for a 64-Point, Radix-4 FFT.

```
*
* APPENDIX F2
*
* FILE TO BE LINKED WITH THE SOURCE CODE FOR A 64-POINT, RADIX-4 FFT.
*
*
*      .global SINE
*      .global N
*
*      .set 64
*      .set 6
*
*      .data
*
* SINE
*      .float 0.000000
*      .float 0.098017
*      .float 0.195090
*      .float 0.290285
*      .float 0.382683
*      .float 0.471397
*      .float 0.555570
*      .float 0.634393
*      .float 0.707107
*      .float 0.773010
*      .float 0.831470
*      .float 0.881921
*      .float 0.923880
*      .float 0.956940
*      .float 0.980785
*      .float 0.995185
*
* COSINE
*      .float 1.000000
*      .float 0.995185
*      .float 0.980785
*      .float 0.956940
*      .float 0.923880
*      .float 0.881921
*      .float 0.831470
*      .float 0.773010
*      .float 0.707107
*      .float 0.634393
*      .float 0.555570
*      .float 0.471397
*      .float 0.382683
*      .float 0.290285
*      .float 0.195090
*      .float 0.098017
*      .float 0.000000
*      .float -0.098017
*      .float -0.195090
*      .float -0.290285
*      .float -0.382683
*      .float -0.471397
*      .float -0.555570
*      .float -0.634393
*      .float -0.707107
*      .float -0.773010
*      .float -0.831470
*      .float -0.881921
*      .float -0.923880
*      .float -0.956940
*      .float -0.980785
*      .float -0.995185
*
*      .float -0.555570
*      .float -0.634393
*      .float -0.707107
*      .float -0.773010
*      .float -0.831470
*      .float -0.881921
*      .float -0.923880
*      .float -0.956940
*      .float -0.980785
*      .float -1.000000
*      .float -0.995185
*      .float -0.980785
*      .float -0.956940
*      .float -0.923880
*      .float -0.881921
*      .float -0.831470
*      .float -0.773010
*      .float -0.707107
*      .float -0.634393
*      .float -0.555570
*      .float -0.471397
*      .float -0.382683
*      .float -0.290285
*      .float -0.195090
*      .float -0.098017
*      .float 0.000000
*      .float 0.098017
*      .float 0.195090
*      .float 0.290285
*      .float 0.382683
*      .float 0.471397
*      .float 0.555570
*      .float 0.634393
*      .float 0.707107
*      .float 0.773010
*      .float 0.831470
*      .float 0.881921
*      .float 0.923880
*      .float 0.956940
*      .float 0.980785
*      .float 0.995185
```

Appendix F3. Link Command File

```
*
*  APPENDIX F3
*
*
*  LINK COMMAND FILE
*
*  DO NOT TYPE IN THESE FIRST SEVEN LINES
-o 12opt64.out
12fopt.obj
sin64.obj

SECTIONS
{
    .text : {}
    .data : {}
    IN 809800h : { 12fopt.obj(IN) }
    .bss 809C00h: {}
}
```